



# An Energy-Efficient MAC Protocol for Wireless Sensor Networks

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# Outline

- Introduction
- Protocol description and discussion
  - Basic parameter design
  - Phase switch schedule
  - Time-slot assignment
- Simulations
- Conclusions



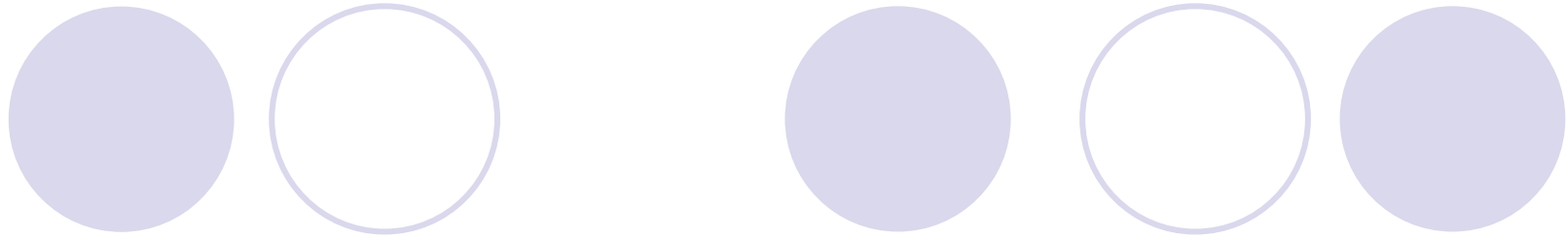
# Introduction

- Most energy-efficient MAC protocols assume perfect network time synchronization
  - Clock drifts exist
- The duration of power on/off is seldom discussed
- Traffic arrival rate for different nodes at different time is fluctuating

# Goals of This Paper



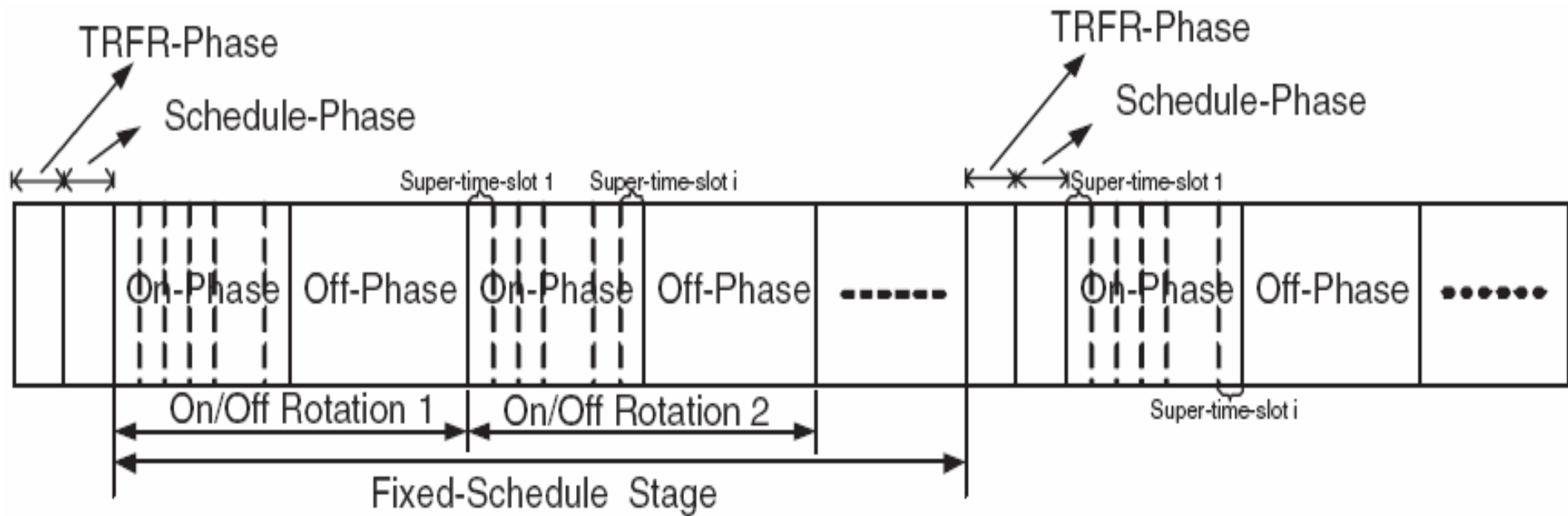
- The authors propose an energy-efficient MAC protocol
  - Remove tight dependency on time synchronization
  - Apply free-running method and fuzzy logic rescheduling scheme, phase-switching schedules are set up and **clock drifts among nodes are compensated.**
  - a trafficstrength- and network-density-based model to **determine essential algorithm parameters**



## ASCEMAC

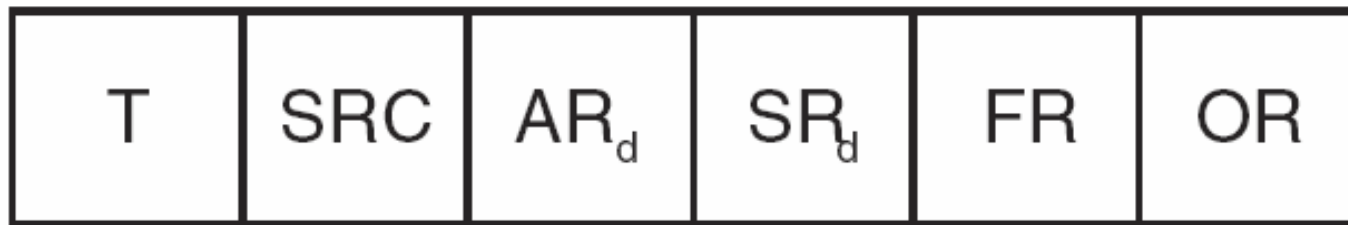
Asynchronous Energy-Efficient MAC Protocol

# ASCEMAC Protocol Model



1. **TRFR-Phase:** cluster members transmit message to cluster head
2. **Schedule-Phase:** cluster heads to broadcast phase-switching schedules
3. **On-Phase:** all nodes to power on their radios to make communication
4. **Off-phase:** all nodes to power off their radios

# TRFR Message Format



T	packet type
SRC	source address
AR <sub>d</sub>	data arrival rate
SR <sub>d</sub>	data service rate
FR	transmission failure rate
OR	buffer overflowing rate

Information for a cluster head to determine the scheduling

# Parameter Design

Combine

$$T_f = \min \left\{ (2W_{max} - T_n), \min_i \left( \frac{k_i}{\lambda_i} - T_n \right) \right\}$$
$$T_n = \frac{T_d T_f \sum_{i=1}^N \lambda_i}{1 - T_d \sum_{i=1}^N \lambda_i} \longrightarrow \phi = \sum_{i=1}^N \lambda_i$$

$T_n$ : on phase duration

$W_{max}$ : maximum waiting time

$T_f$ : off-phase duration

$\lambda_i$ : traffic arrival rate for node  $i$

$T_d$ : slot time

$K_i$ : buffer size



# Parameter Design

$$1) \text{ when } 2W_{max} < \min_i \left( \frac{k_i}{\lambda_i} \right) \quad 2) \text{ when } 2W_{max} \geq \min_i \left( \frac{k_i}{\lambda_i} \right)$$

$$\begin{cases} T_f = 2W_{max}(1 - T_d\phi) \\ T_n = 2W_{max}T_d\phi \end{cases} \quad \begin{cases} T_f = \min_i \left( \frac{k_i}{\lambda_i} \right) (1 - T_d\phi) \\ T_n = \phi T_d \min_i \left( \frac{k_i}{\lambda_i} \right) \end{cases}$$

- The goal of adjusting parameters
  - Extend the power off time to save more energy
  - Adjust the data packet waiting time to an acceptable value

# Schedule Message Format

T	SRC	$D_{off}$	$D_{on}$
$SRC_1$	$DEST_1$	$D_{df1}$	$D_{s1}$
$SRC_2$	$DEST_2$	$D_{df2}$	$D_{s2}$
-----			
$SRC_i$	$DEST_i$	$D_{dfi}$	T

- SRC source address
- $D_{off}$  off phase duration
- $D_{on}$  on phase duration
- $D_{si}$  supper time slot duration for node i
- $D_{dfi}$  supper time slot starts defer time
- $SRC_1$  source address for ith supper time slot
- $DEST_i$  destination address for ith supper time slot



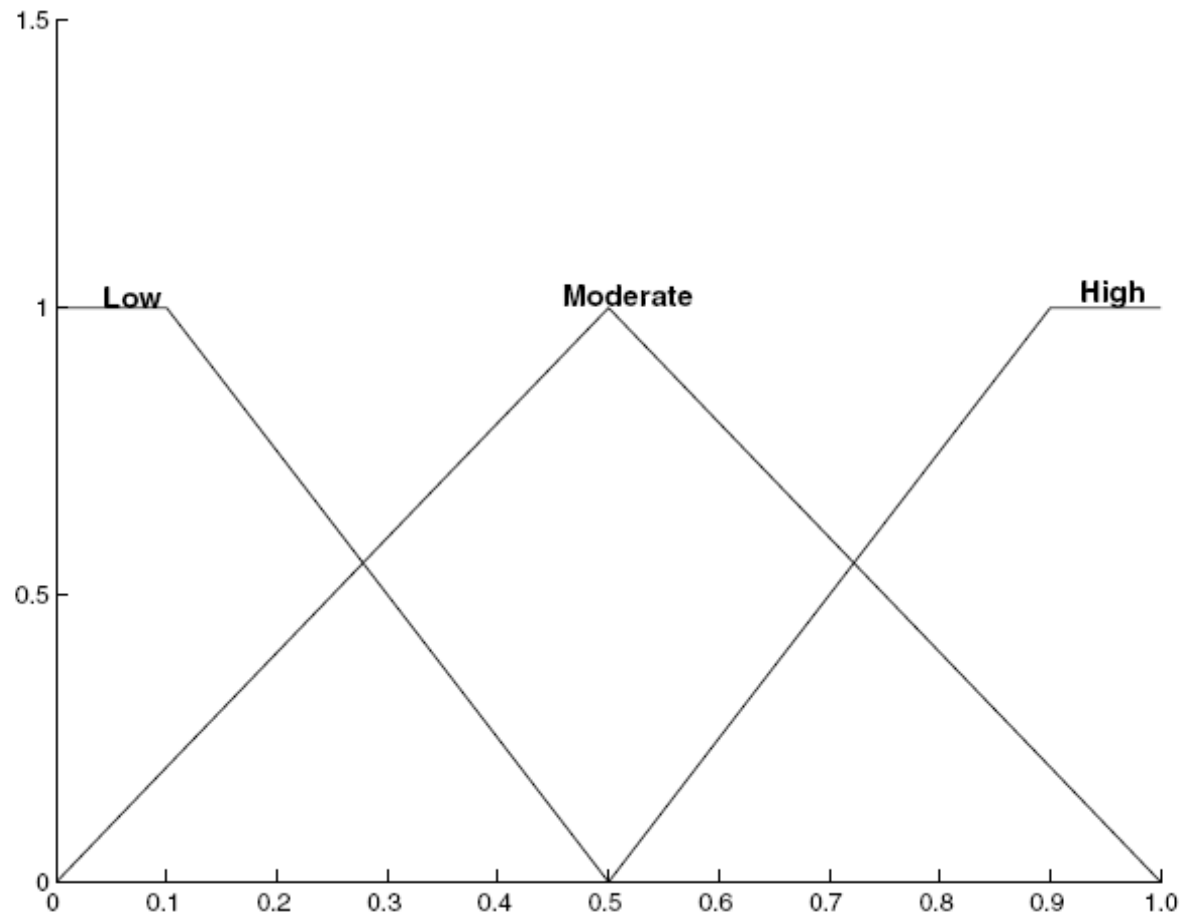
# Interval of Schedule Broadcast

$$T_i = \xi_i \times T_{i-1}$$

- $\xi_i$  is the interval adjusting function
- Three parameters are used to design  $\xi_i$ 
  - **Ante1**: ratio of nodes with overflowed buffer
  - **Ante2**: ratio of nodes with high failing transmission rate
  - **Ante3**: ratio of nodes experiencing unsuccessful transmission

# Interval of Schedule Broadcast

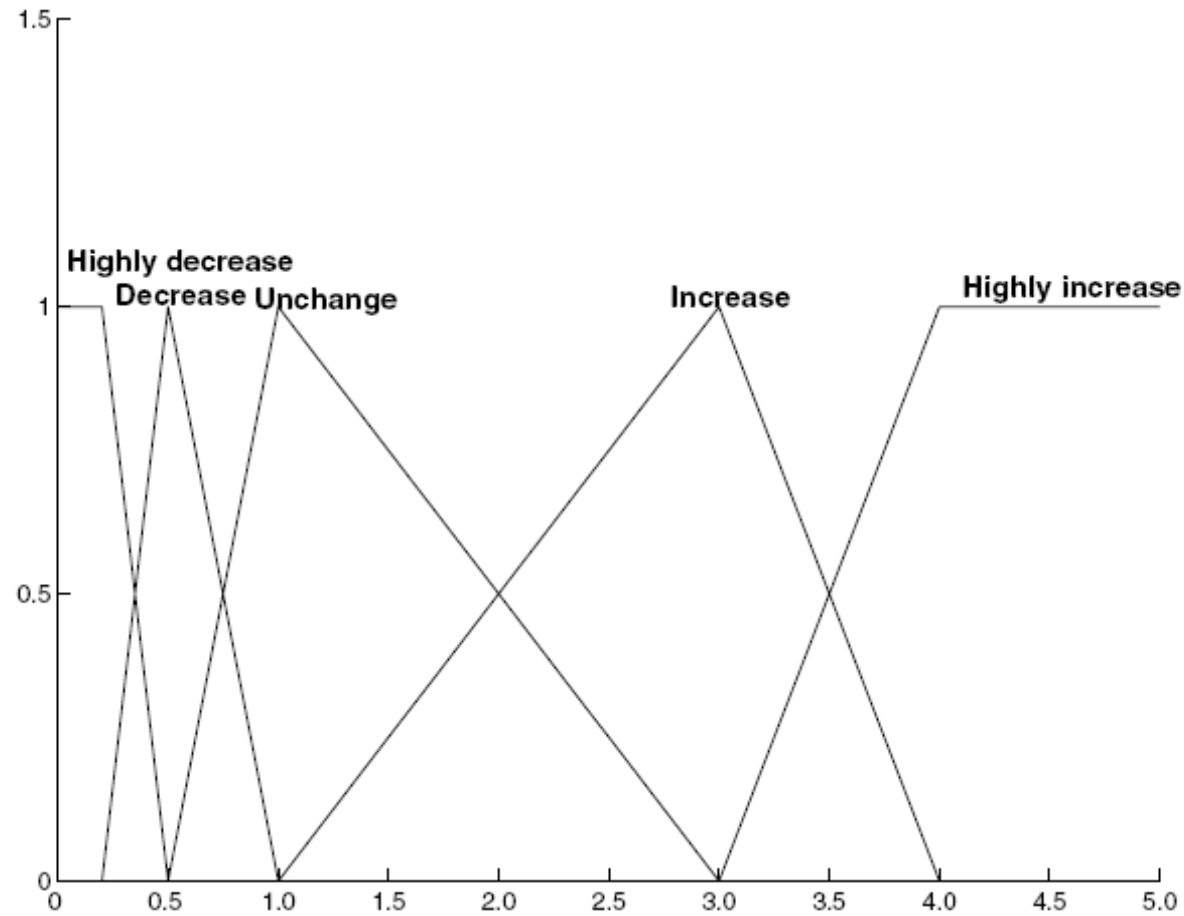
- The linguistic variables used to represent those parameters are divided into three levels: *Low*, *Moderate*, *High*



# The rules for adjusting the interval of schedule broadcast

<i>Rule</i>	<i>Ante1</i>	<i>Ante2</i>	<i>Ante3</i>	<i>Consequent</i>
1	<i>Low</i>	<i>Low</i>	<i>Low</i>	<i>HighlyIncrease</i>
2	<i>Low</i>	<i>Low</i>	<i>Moderate</i>	<i>Increase</i>
3	<i>Low</i>	<i>Moderate</i>	<i>Moderate</i>	<i>Decrease</i>
4	<i>Low</i>	<i>Moderate</i>	<i>High</i>	<i>Decrease</i>
5	<i>Moderate</i>	<i>Low</i>	<i>Moderate</i>	<i>Increase</i>
6	<i>Moderate</i>	<i>Low</i>	<i>High</i>	<i>Unchange</i>
7	<i>Moderate</i>	<i>Moderate</i>	<i>Moderate</i>	<i>Decrease</i>
8	<i>Moderate</i>	<i>Moderate</i>	<i>High</i>	<i>HighlyDecrease</i>
9	<i>Low</i>	<i>High</i>	<i>High</i>	<i>Decrease</i>
10	<i>Moderate</i>	<i>High</i>	<i>High</i>	<i>HighlyDecrease</i>
11	<i>High</i>	<i>Low</i>	<i>Moderate</i>	<i>Increase</i>
12	<i>High</i>	<i>Low</i>	<i>High</i>	<i>Unchange</i>
13	<i>High</i>	<i>Moderate</i>	<i>Moderate</i>	<i>Decrease</i>
14	<i>High</i>	<i>Moderate</i>	<i>High</i>	<i>Decrease</i>
15	<i>High</i>	<i>High</i>	<i>High</i>	<i>HighlyDecrease</i>

# Interval of Schedule Broadcast



# Using Fuzzy Function

- The authors mention that  $\xi_i$  is defuzzified by

$$\xi(x_1, x_2, x_3) = \frac{\sum_{l=1}^{15} \bar{\xi}^l \mu_{F_1^l}(x_1) \mu_{F_2^l}(x_2) \mu_{F_3^l}(x_3)}{\sum_{l=1}^{15} \mu_{F_1^l}(x_1) \mu_{F_2^l}(x_2) \mu_{F_3^l}(x_3)}$$

E. H. Mandani, "Application of fuzzy logic to approximate reasoning using linguistic systems", *IEEE Trans. On system, Man, and Cybernetics*, vol. 26, no. 12, pp. 1182-1191, 1977.

# Time-slot assignment



$\Delta t_1$ : time difference between nodes

$T_{s,min}$ : the least time needed to detect the synchronization information

1) If  $T_{s,min} < \Delta t_1 < T_d$

$\frac{n-1}{n} \%$  Successful transmission rate

2) If  $T_d < \Delta t_1 < 2T_d$

$\frac{n-2}{n} \%$  Successful transmission rate

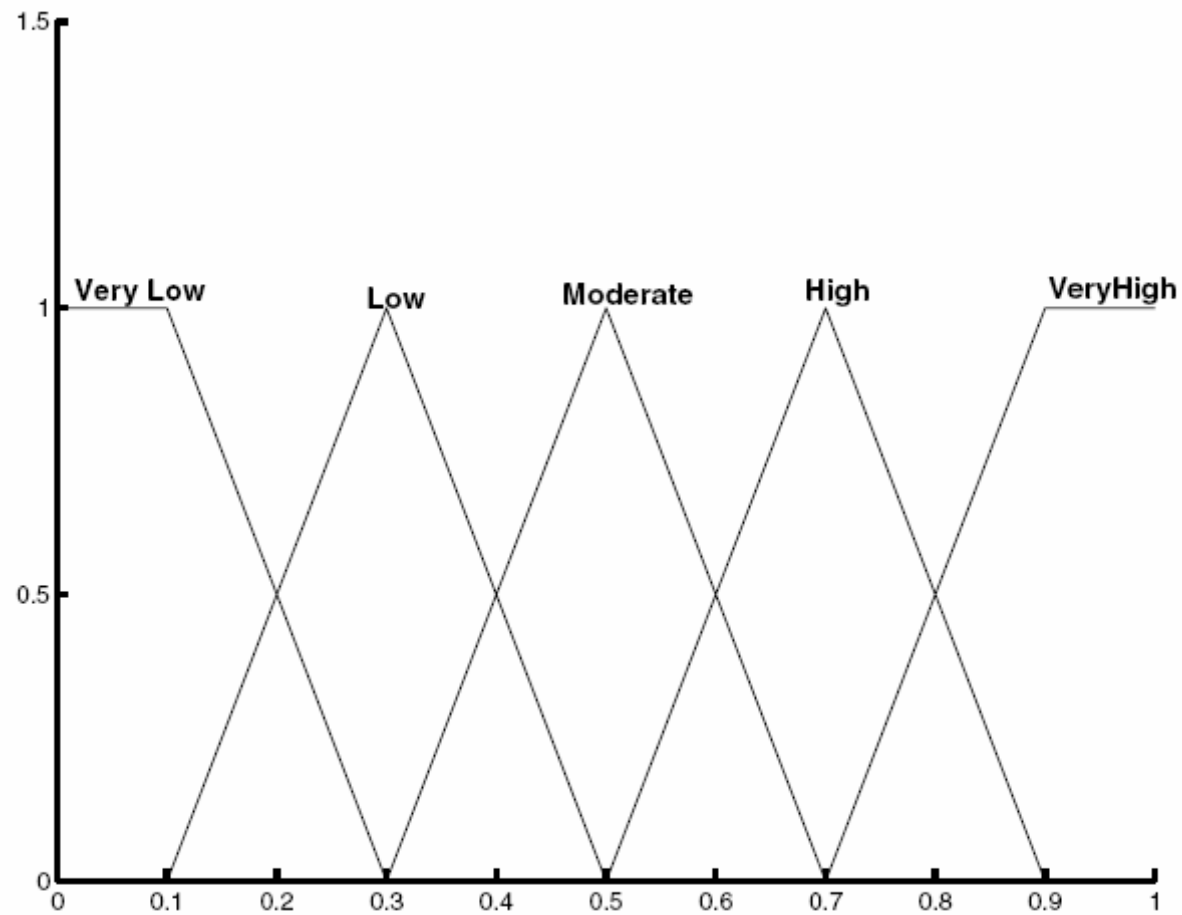


# Rules for Time Slot Allocation

- Ante1: traffic arrival rate
- Ante2: unsuccessful transmission rate

<i>Rule</i>	<i>Antecedent1</i>	<i>Antecedent2</i>	<i>Consequent</i>
1	<i>Low</i>	<i>Low</i>	<i>Moderate</i>
2	<i>Low</i>	<i>Moderate</i>	<i>High</i>
3	<i>Low</i>	<i>High</i>	<i>VeryHigh</i>
4	<i>Moderate</i>	<i>Low</i>	<i>Low</i>
5	<i>Moderate</i>	<i>Moderate</i>	<i>Moderate</i>
6	<i>Moderate</i>	<i>High</i>	<i>High</i>
7	<i>High</i>	<i>Low</i>	<i>VeryLow</i>
8	<i>High</i>	<i>Moderate</i>	<i>Low</i>
9	<i>High</i>	<i>High</i>	<i>Moderate</i>

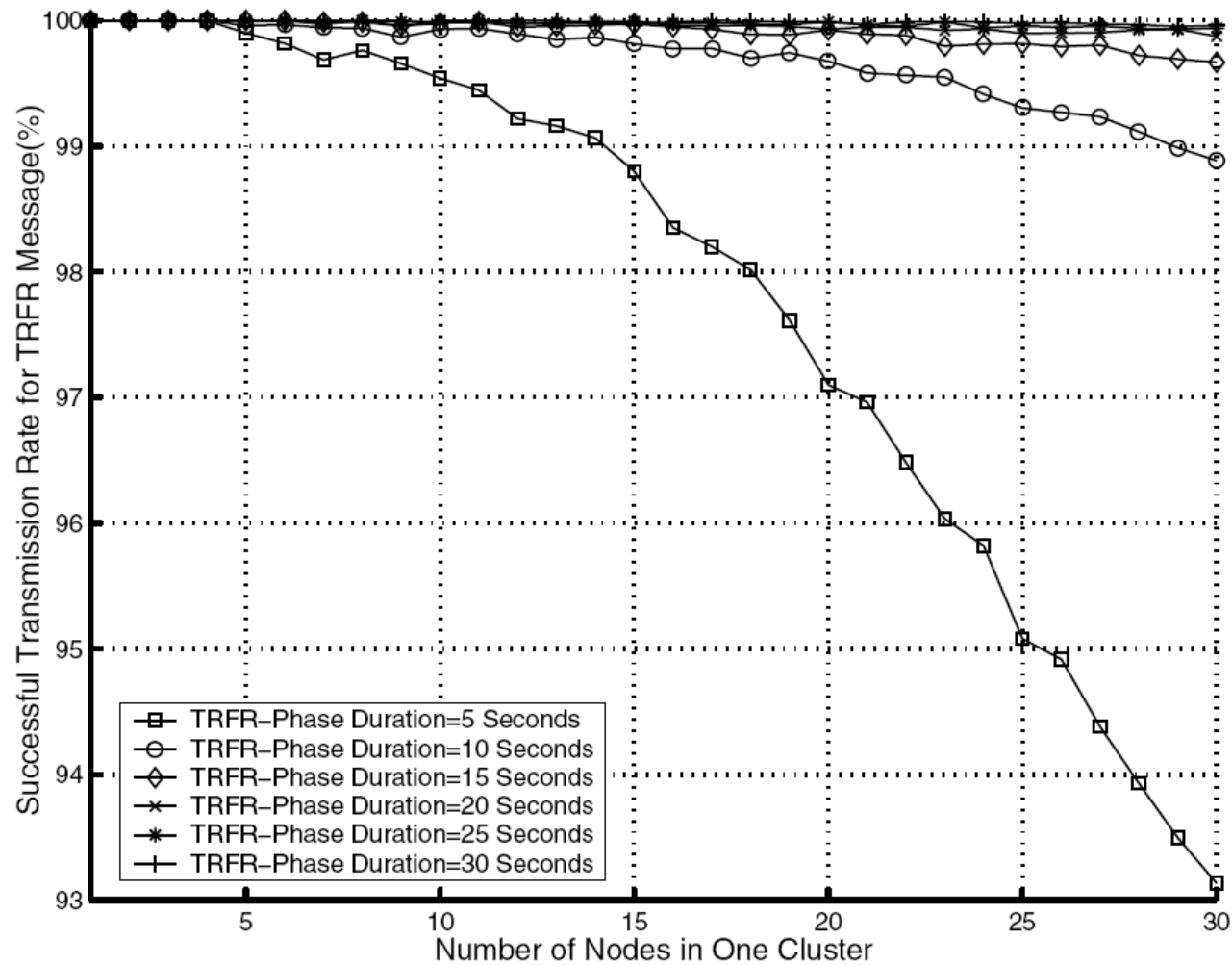
# The MF of time-slot assignment



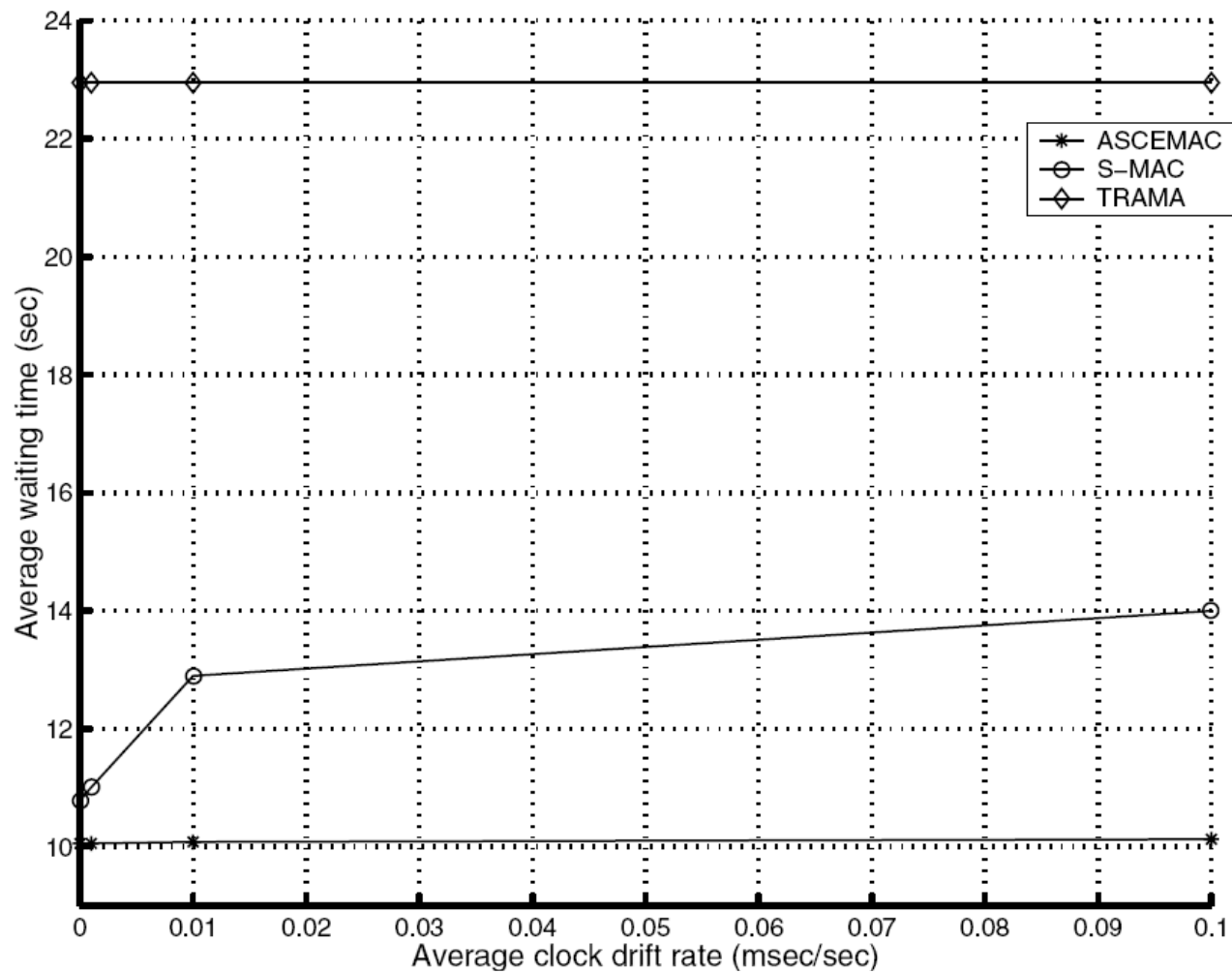
# Simulation Environment

Simulator	OPNET
Area	100m x 100m
Radio range	30m
Symbol rate	40 Ksps
Data frame length	1024 bits
Clock drift range	1 to 100us

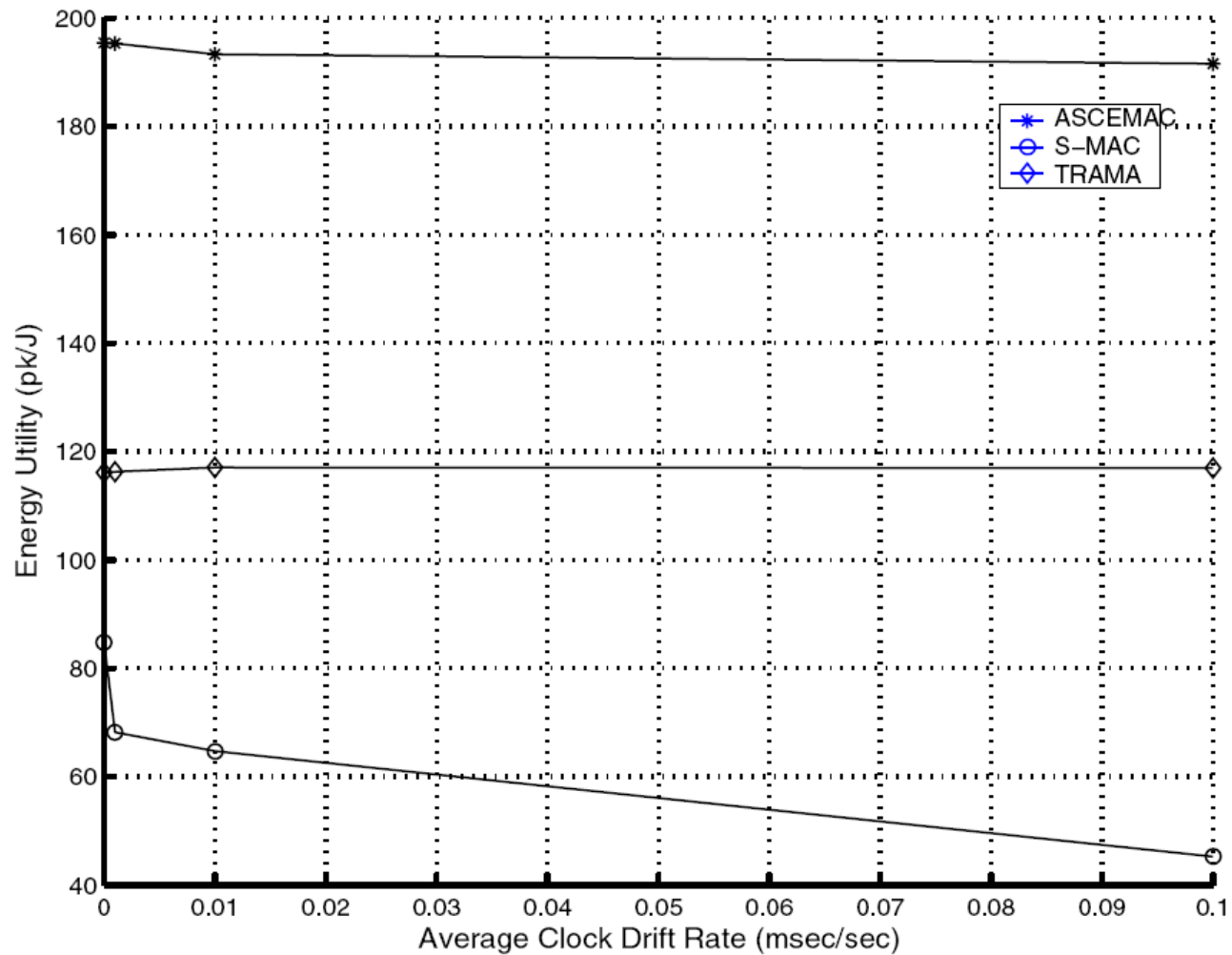
# Successful Transmission Rate



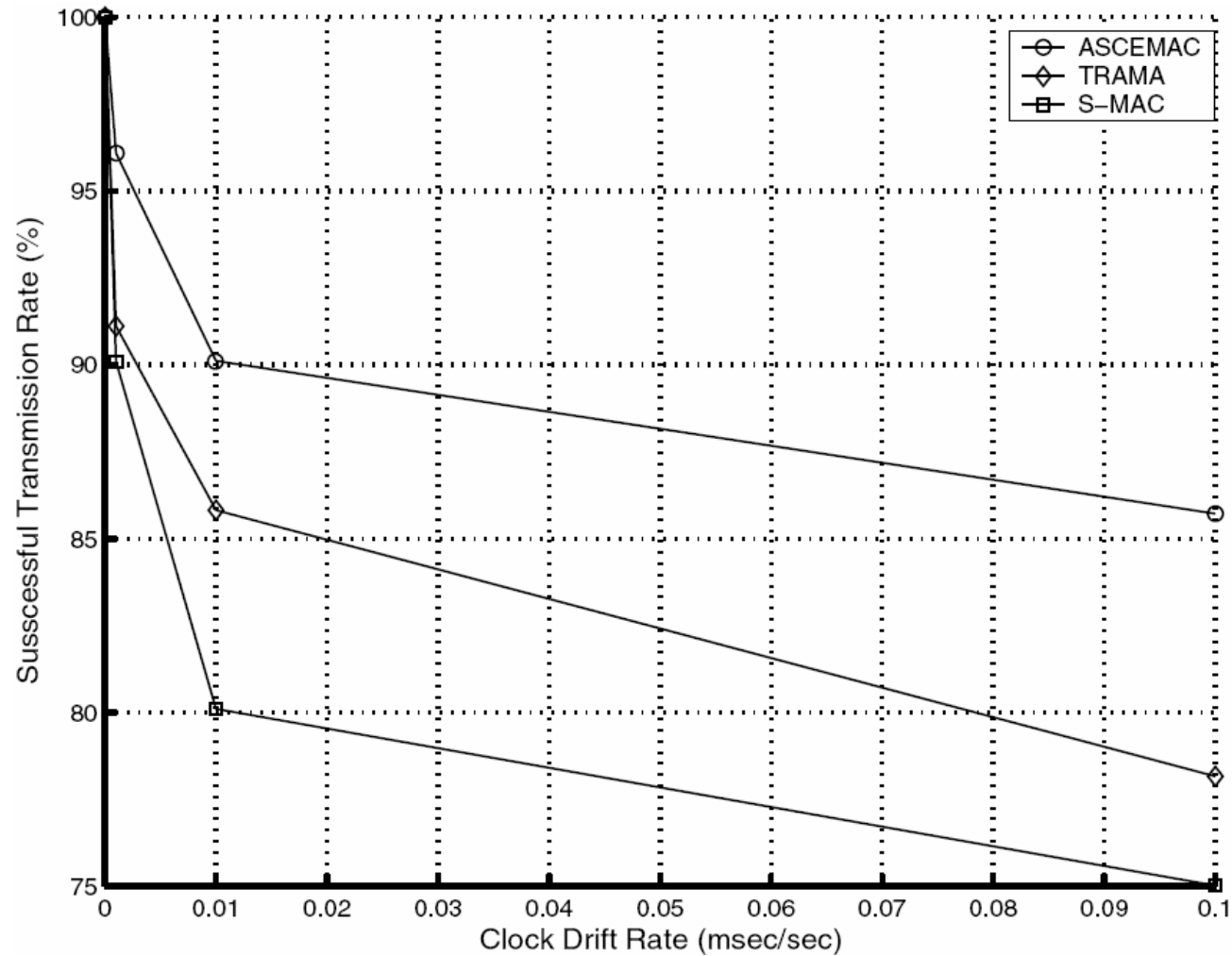
# Average Waiting Time



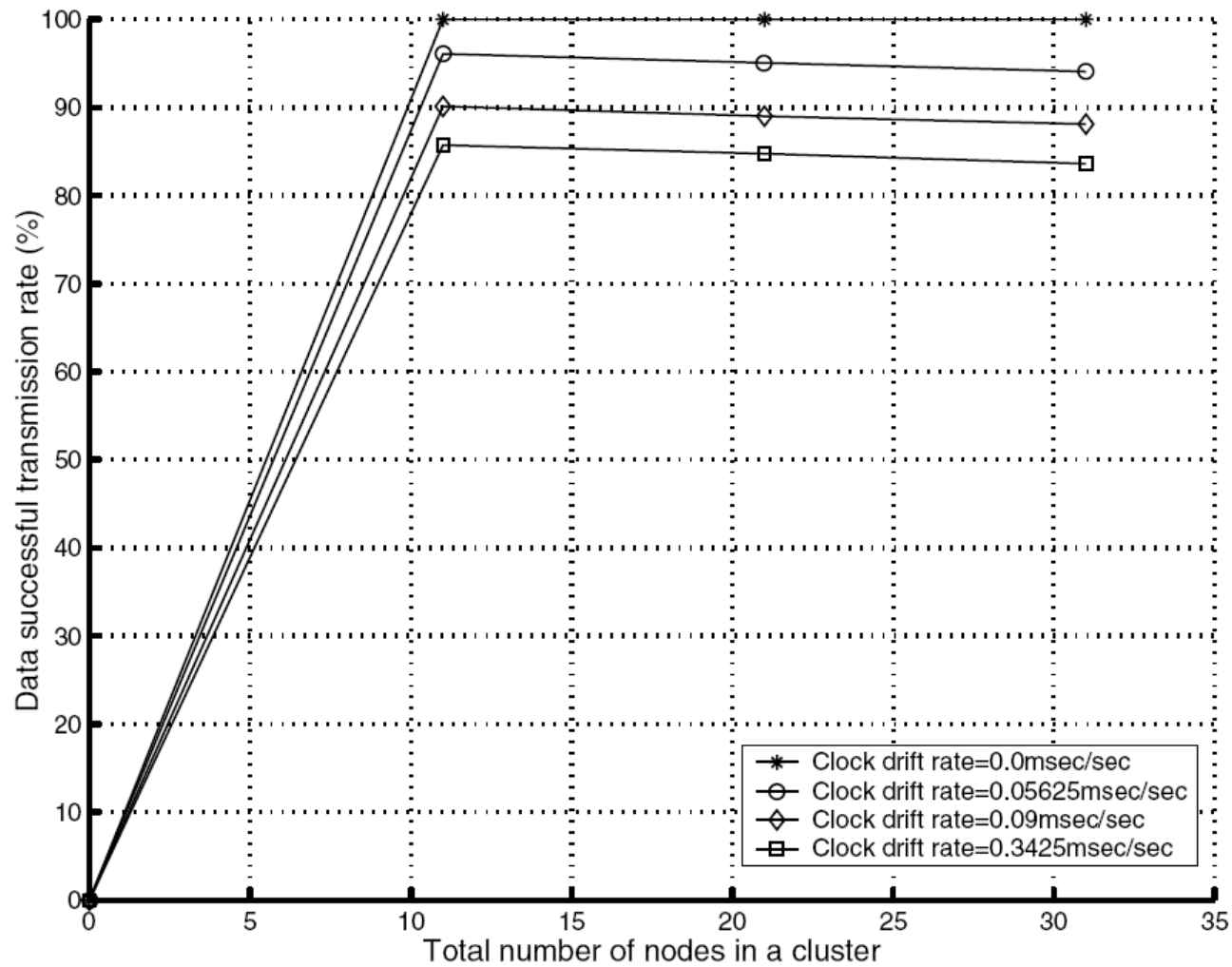
# Average Energy Utility



# Successful Transmission Rate

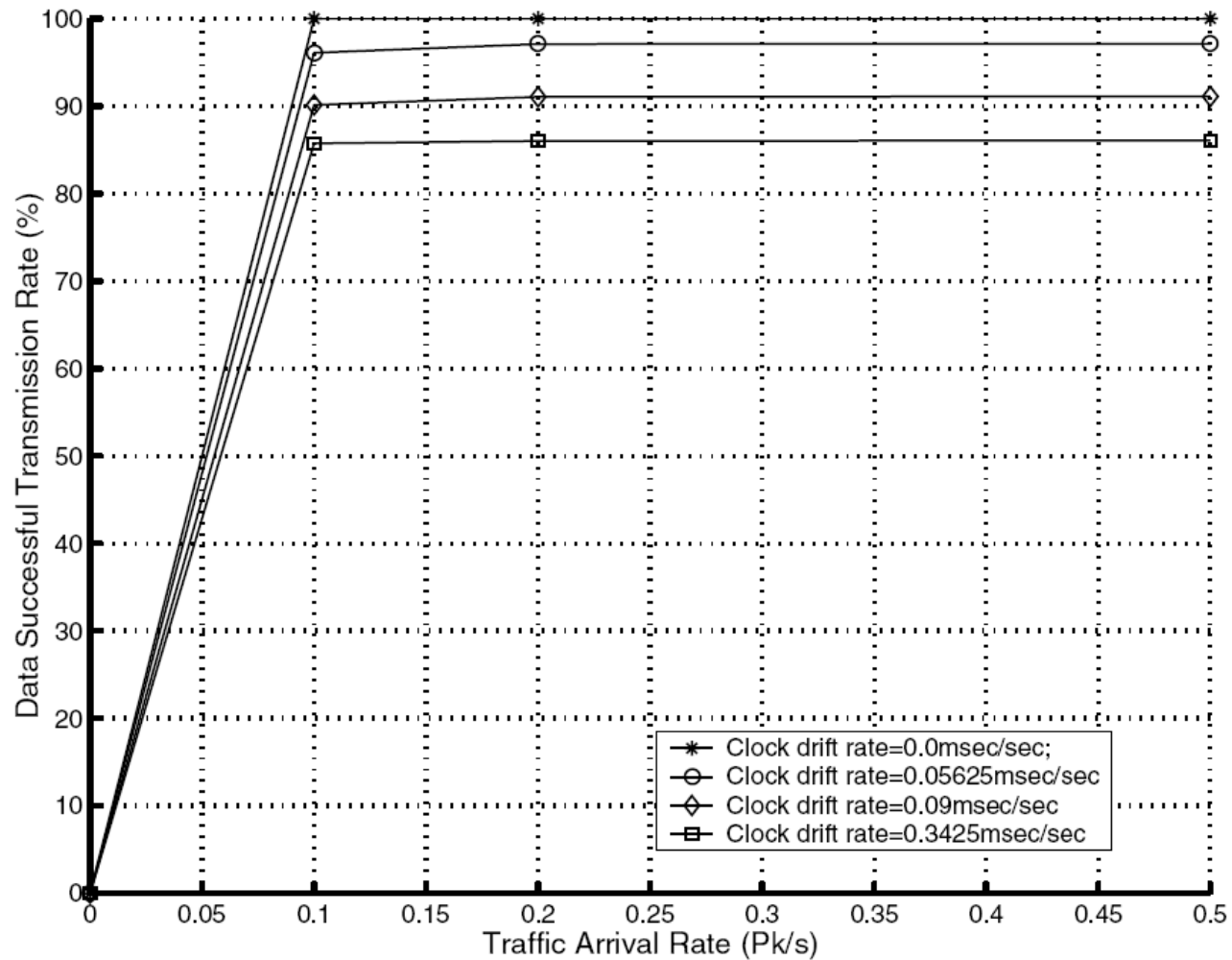


# Node Density Adaptation





# Traffic Strength Adaptation



# Conclusions



- Exploiting a rescheduling method, instead of time synchronization, to handle mismatch caused by clock drifts, as well as taking advantage of fuzzy logical theory, which has distinctive capabilities for coping with uncertainty
- ASCEMAC acquires the optimal values of essential algorithm parameters
  - Ensure average successful transmission rate
  - Decrease the data packet average waiting time
  - Reduce energy consumption



THANK YOU