

Location Awareness in Unstructured Peer-to-Peer Systems

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KuYou Lin

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Outline

- Introduction
 - Related Work
 - Location-Aware Topology Matching
 - Three Main Operations
 - Simulation & Performance Evaluation
 - Conclusion
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1. Introduction

□ LTM

(Location-aware Topology matching)

- Unstructured P2P
- Solve **Mismatching** problem
- Distributed , Doesn't require global view

□ Main idea

- Build an efficient overlay by **disconnecting** slow connection
 - Choosing **closer peer** as neighbor
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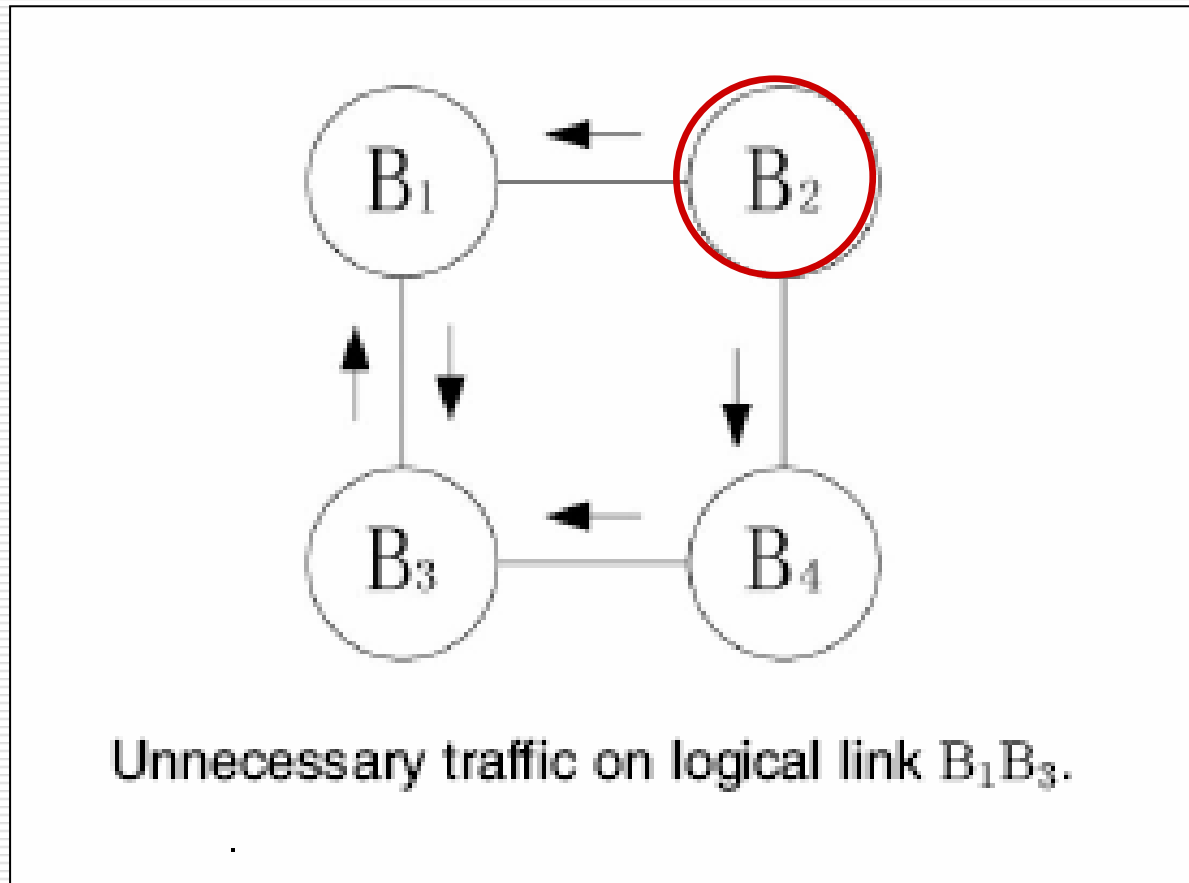
2.Related Work

- Other ways to reduce traffic cost in unstructure-P2P system
 - Forwarding-based
 - Only forward to subset of neighbors
 - Cache-based
 - Remember index of files/peers used before
 - Overlay topology optimization
 - Logical topology / physical topology
 - LTM
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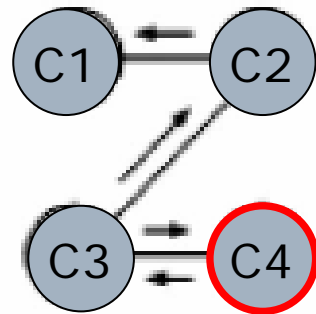
3. Location-Aware Topology Matching

- Problems at unstructured P2P system
 - Unnecessary Message Duplications in **Overlay** Connections
 - Topology Mismatch (logical/physical)
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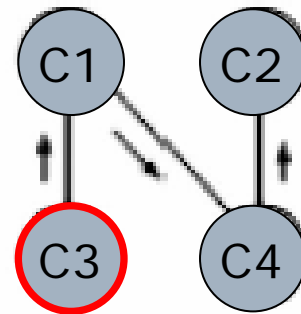
Unnecessary Message Duplications in Overlay Connections



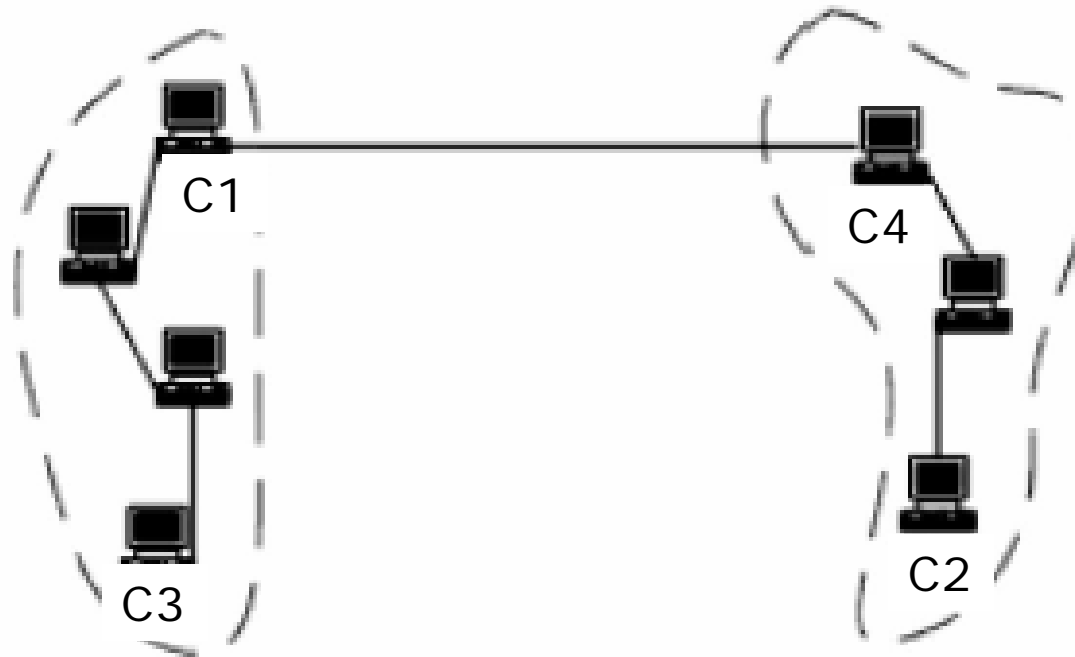
Topology Mismatch



(a)



(b)



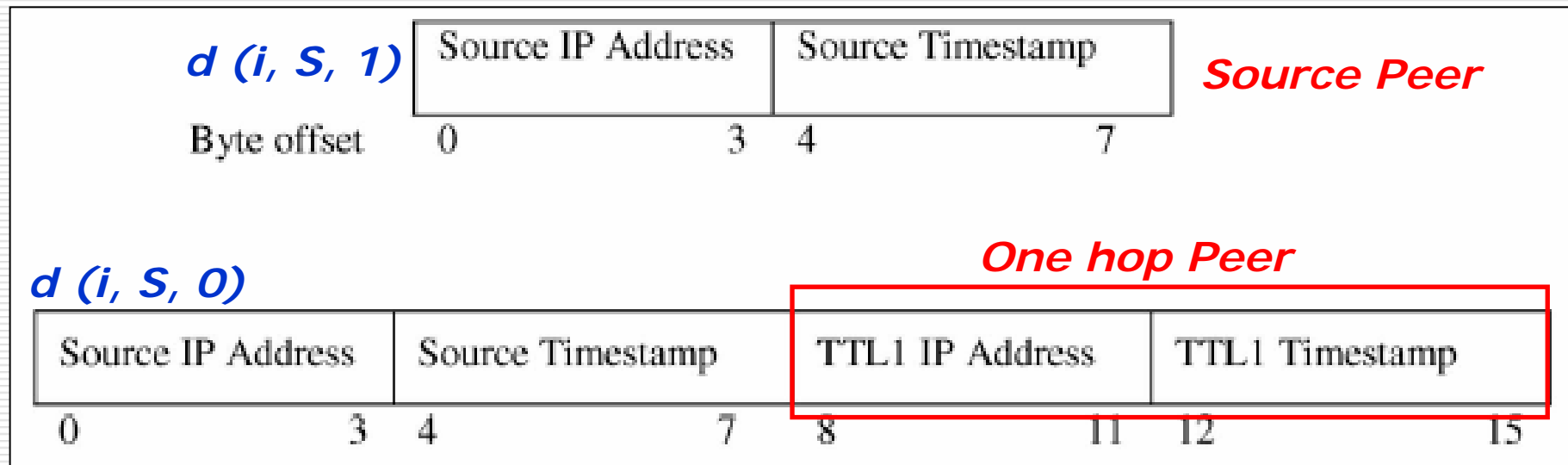
(c)

4. Three Main Operations

- TTL2-Detector Flooding
 - Flooding 2 hops
 - Slow Connection Cutting
 - Source Peer Probing
-

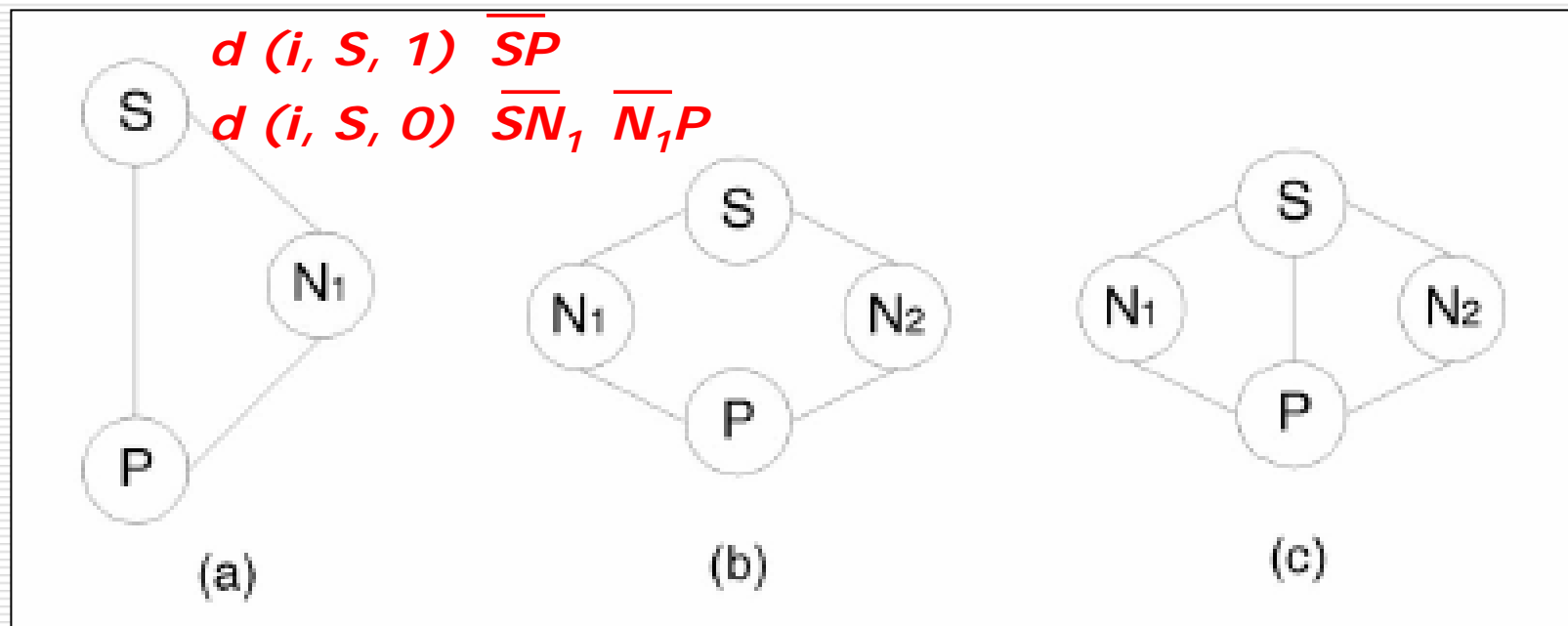
TTL2-Detector Flooding

- Each peer floods a TTL2-Detector periodically
- $d(i, S, v)$
 - i : Message ID
 - S : Initial Source
 - v : TTL value



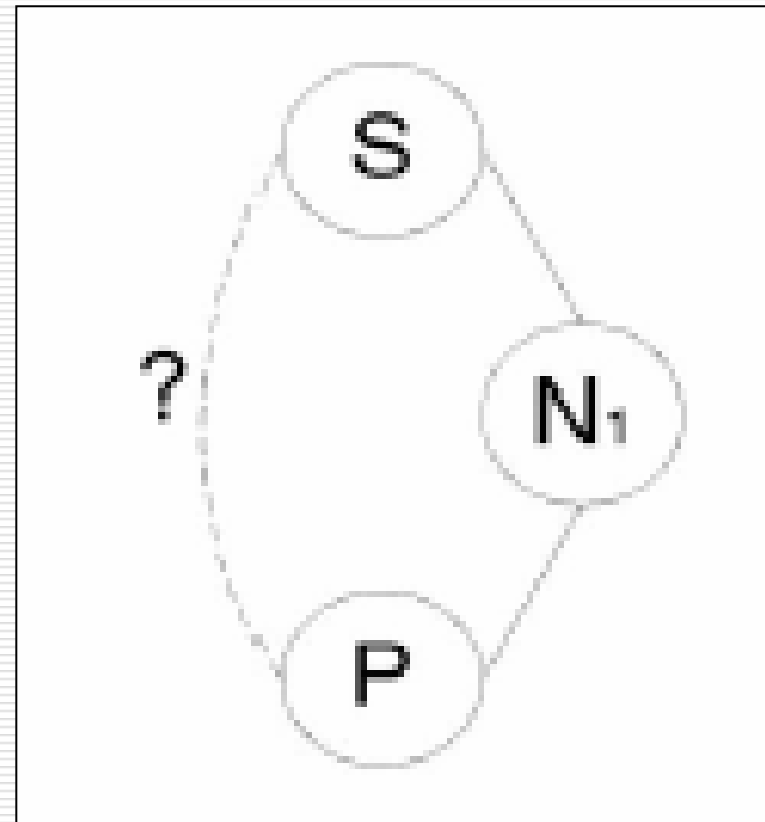
Slow Connection Cutting

- 3 kinds of cases
- Will-Cut List / Cut List

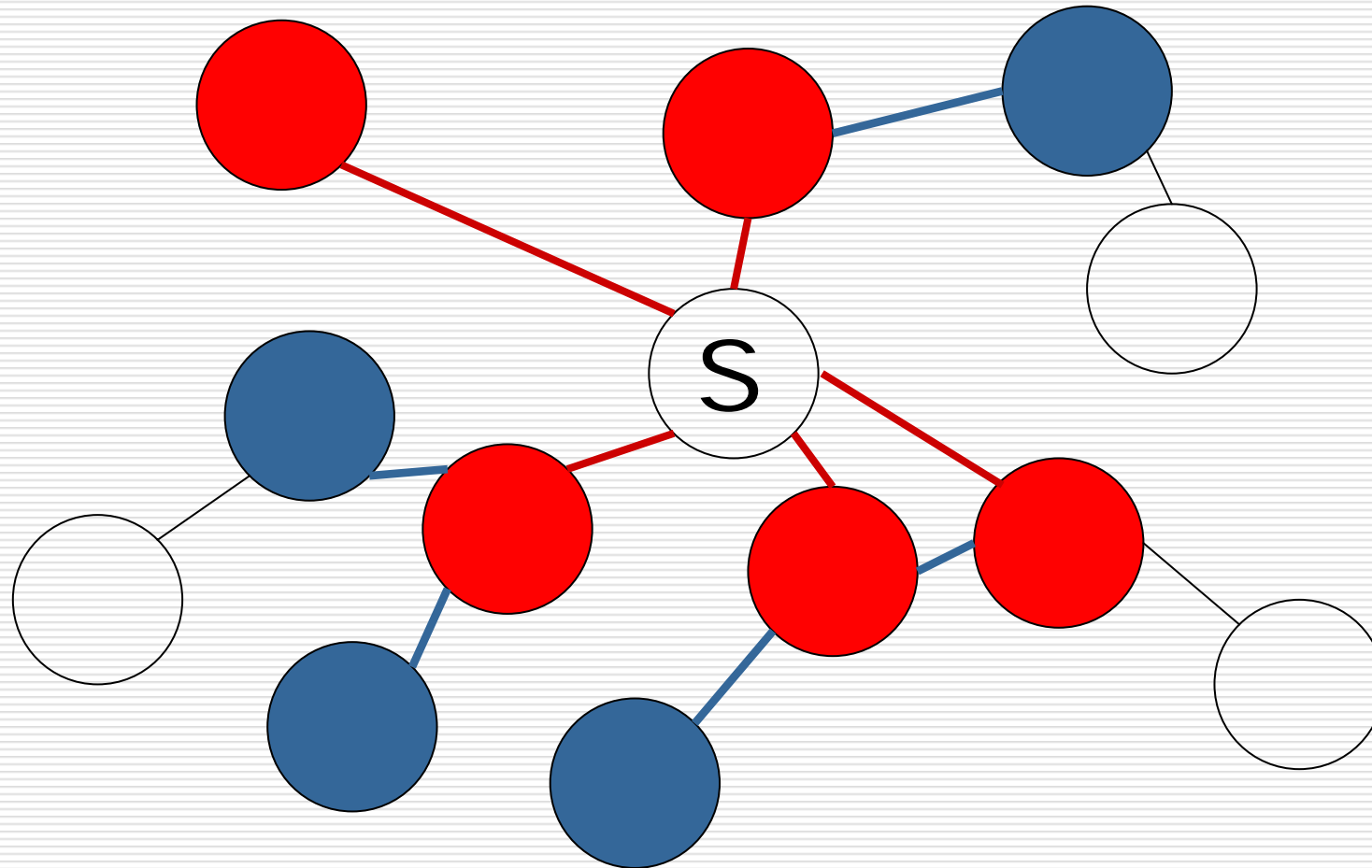


Source Peer Probing

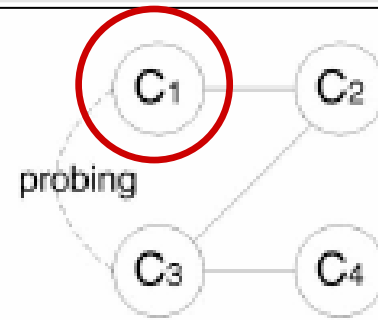
- Peer P receives only one $d(i, s, 0)$ during a certain time
- $P \in (N^2(S) - N(S))$



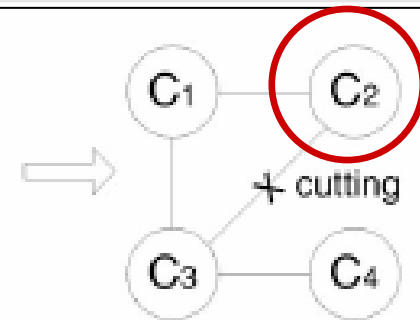
Definition : $N(S)$ $N^2(S)$



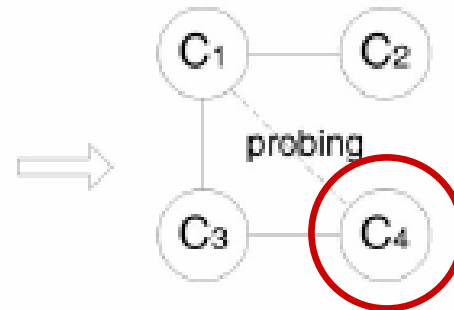
Operation Example



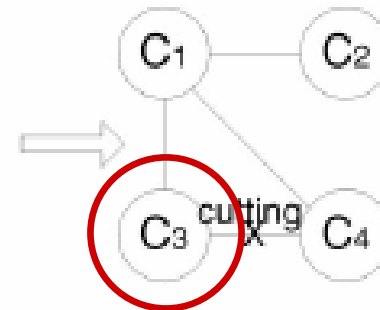
(a)



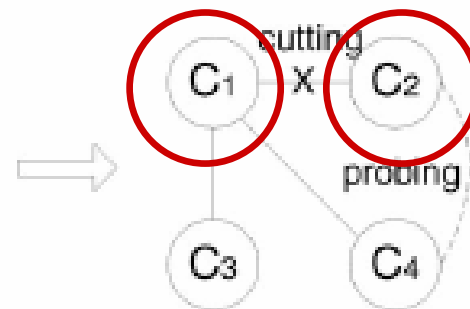
(b)



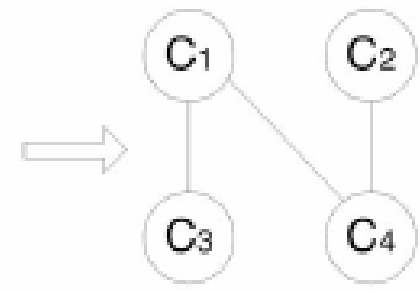
(c)



(d)



(e)



(f)

5. Simulation & Performance Evaluation

□ Performance Metrics

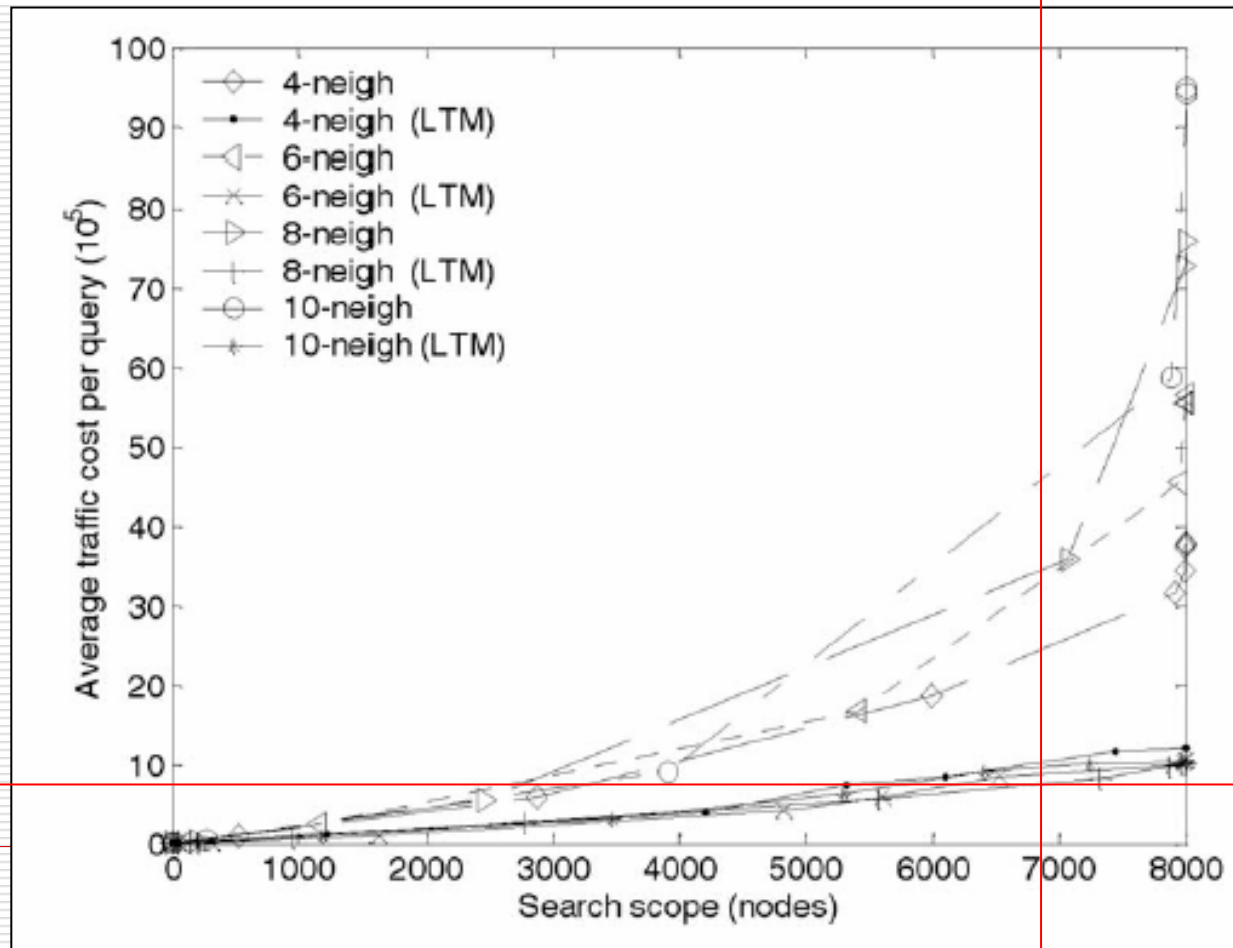
- Average traffic cost vs. search scope
 - $T_c = \text{Message} * \text{number of Links}$
 - Search scope = number of peers reached
 - Average neighbor distance
 - Query response time
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Simulation Environment

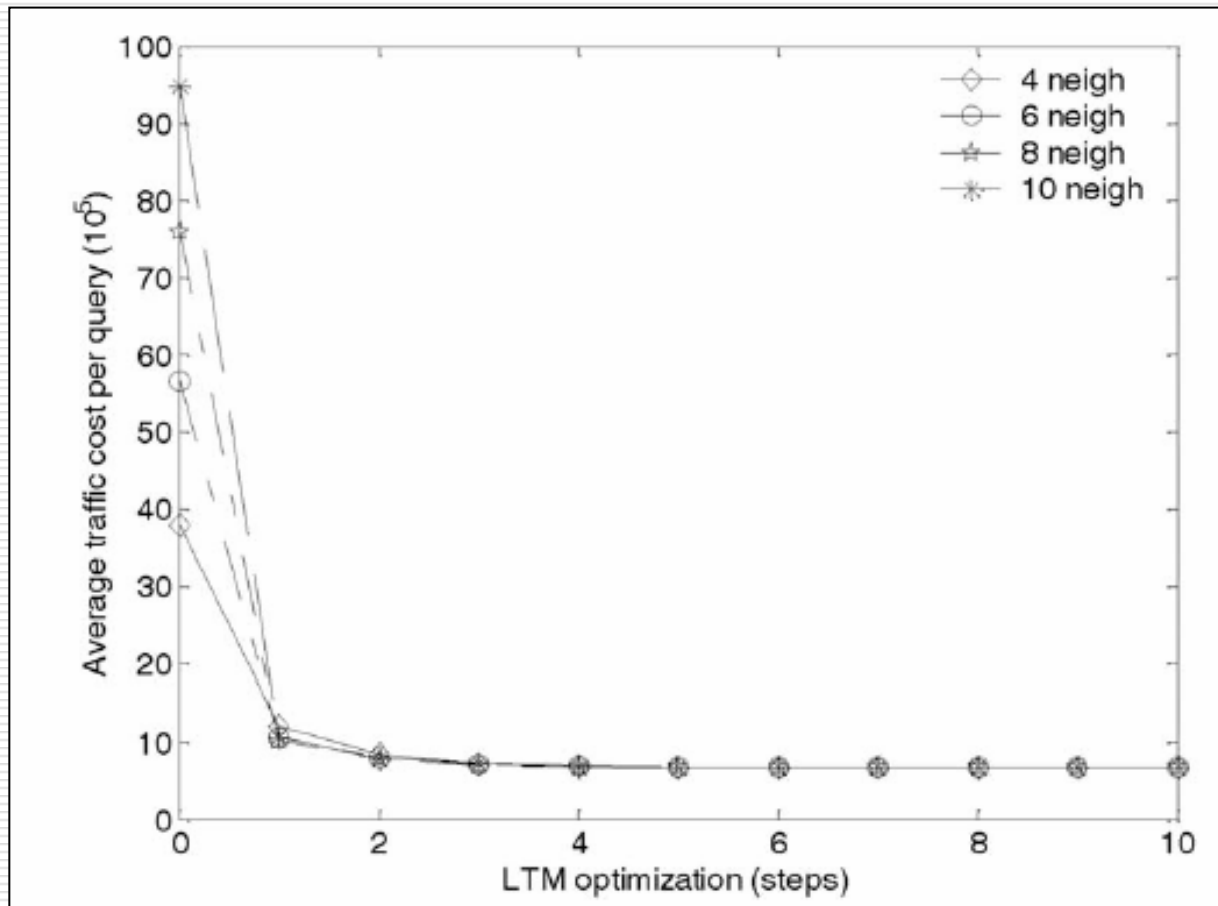
- Overlay (logical) :
2000, 3000, 5000, 8000 nodes
 - Physical :
22000 Internet-like nodes
 - Neighbors :
4, 6, 8, 10 neighbors
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LTM in Static Environment

- Traffic cost vs. Search scope (Only 1 step LTM)

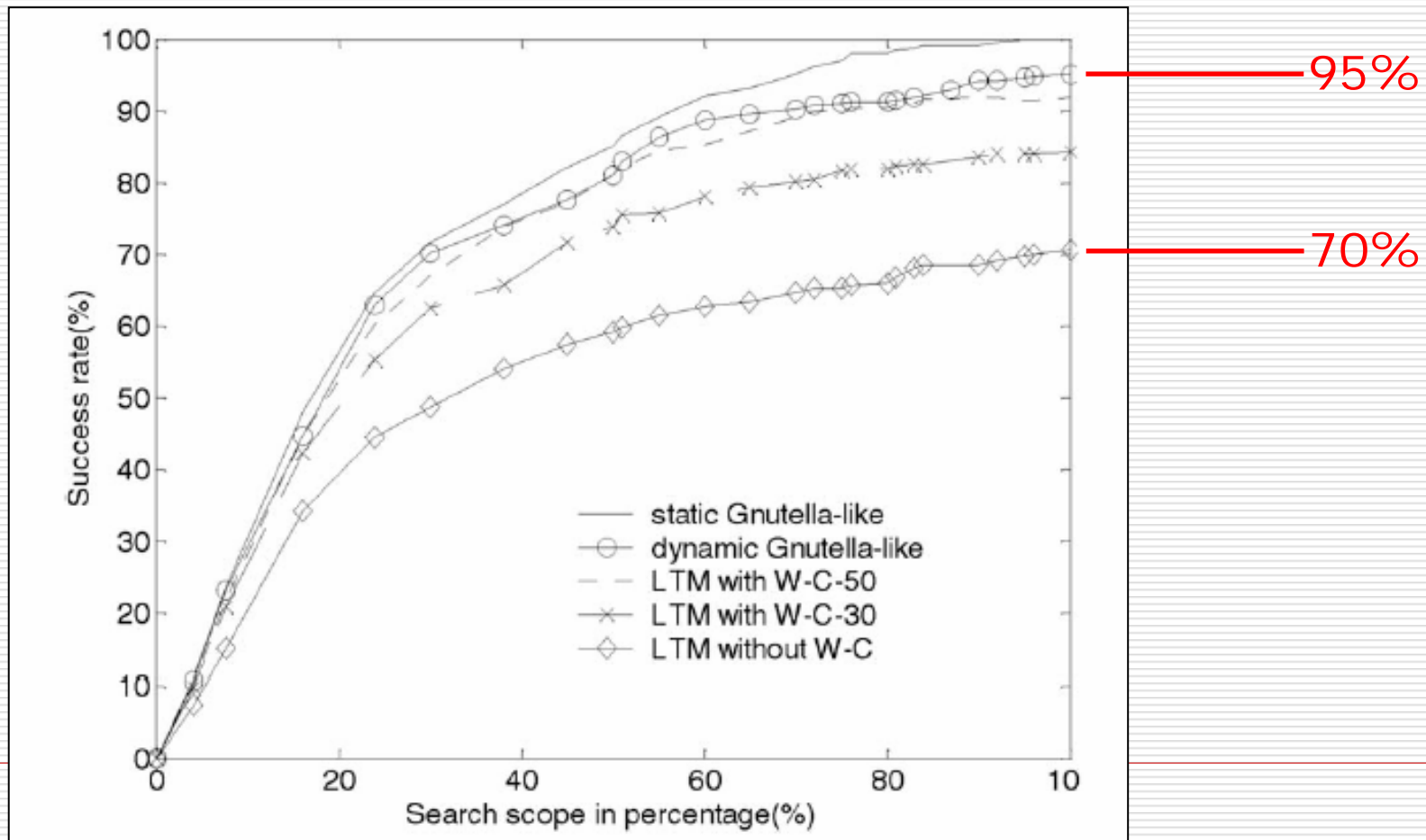


Traffic cost vs. optimization step



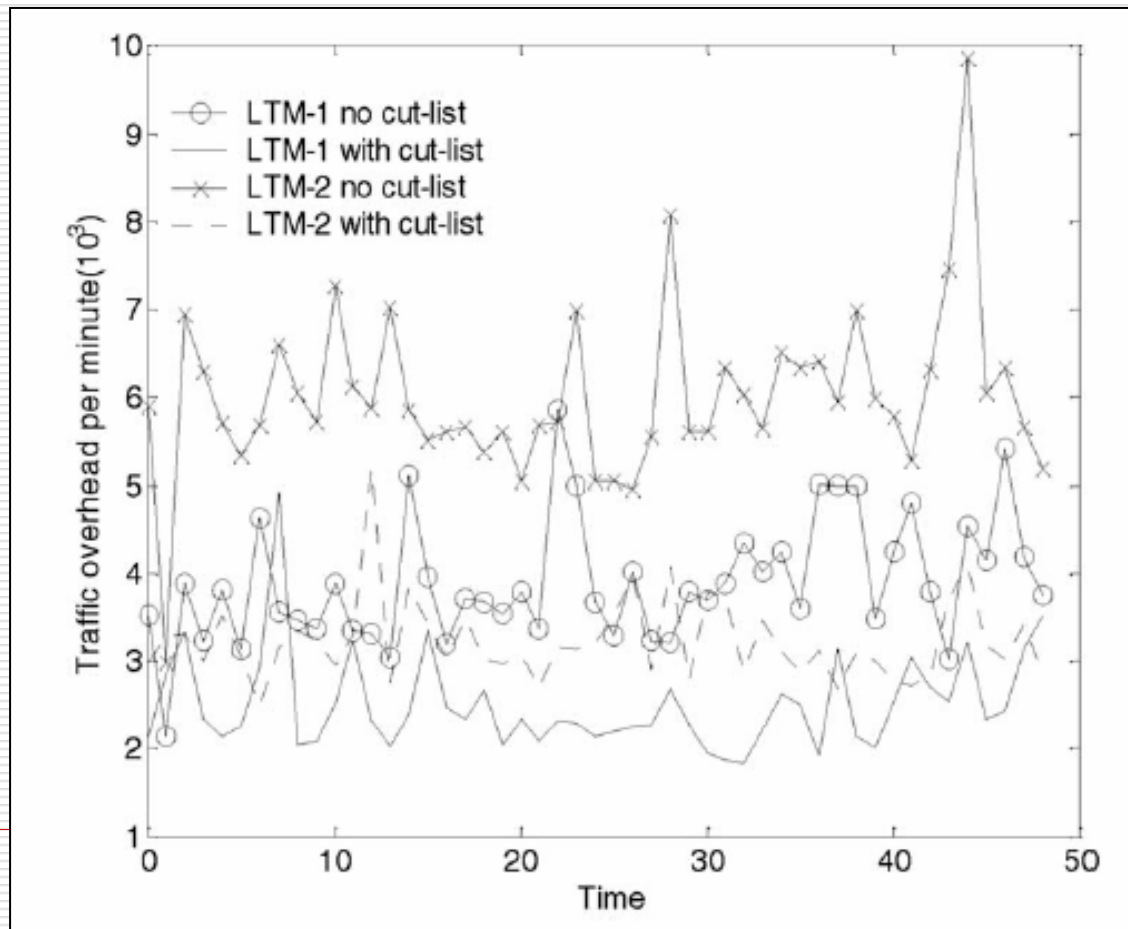
LTM in Dynamic Environment

- Effectiveness of Will-Cut List (W-c-50 → 50 sec)

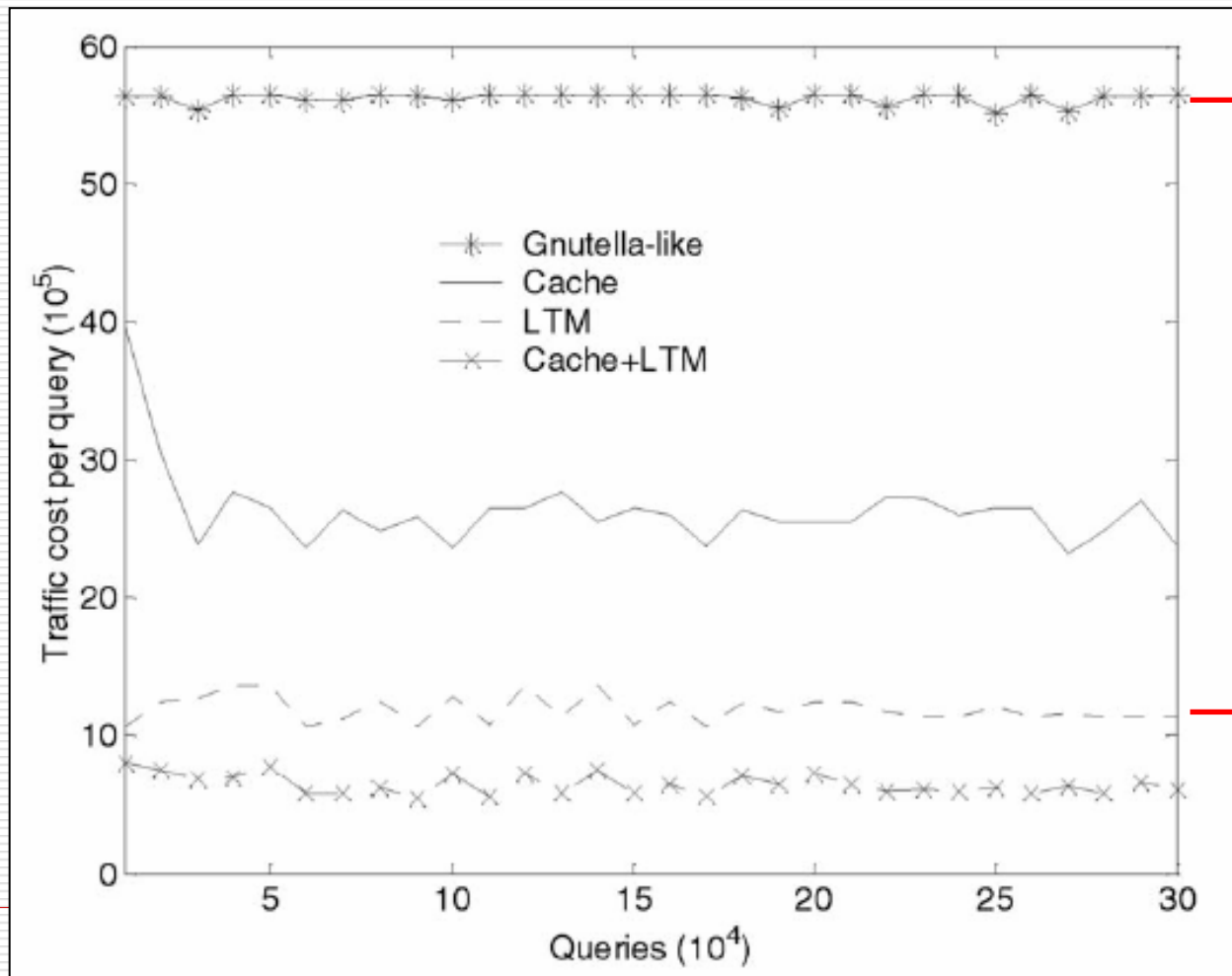


Effectiveness of cut list

- LTM-k means **k-times LTM** / per minute



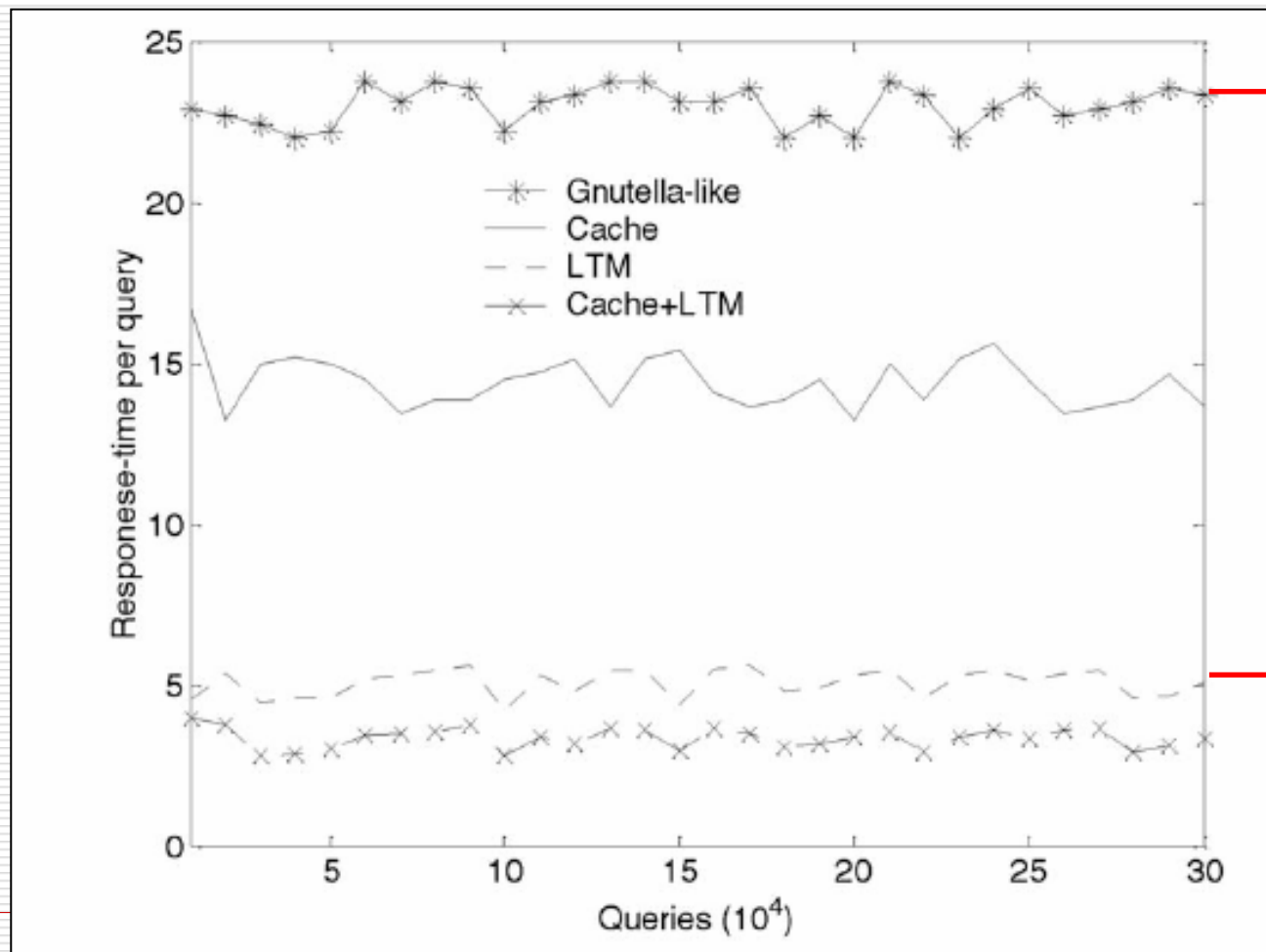
Traffic cost of 4 schemes



57

13

Average response time of 4 schemes



24

6

6. Conclusion

- Using LTM in unstructured P2P system can reduce **75% traffic cost** and **65% query response time**
 - **Will-cut List** and **cut List** can improve the performance of LTM
 - LTM is completely **distributed** and **scalable**
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Reference

- IEEE TRANSACTIONS ON PARALLEL AND DISTRIBUTED SYSTEMS, VOL. 16, NO. 2, FEBRUARY 2005
 - *Yunhao Liu, Member, IEEE, Li Xiao, Member, IEEE, Xiaomei Liu, Lionel M. Ni, Fellow, IEEE, and Xiaodong Zhang, Senior Member, IEEE*
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