



A Reliable Multicast Protocol for Wireless Mobile Multihop Ad Hoc Networks

Selected from CCNC 2004

Present by Lin Yu-Chen

9/17/2004



Outline

- Introduction
- Problem Definition and Motivation
- Proposed Solution
- Simulations and Performance Evaluation
- Conclusions and Discussions



Introduction

- In a typical ad hoc network , hosts work in groups to carry out a given task .
- Some multicast applications , such as audio/video conferencing , can tolerate packet error and/or loss .
- Other applications , such as one-to-many file transfer and military applications do not .



Introduction

- The characteristics of wireless mobile multihop ad hoc networks make reliable multicasting extremely challenging .
 - Lack of network infrastructure
 - Dynamic network topology
 - Scarce bandwidth and variable link capacity
 - High error rates



Problem Definition and Motivation

- A reliable multicasting guarantees the eventual delivery of all the multicast data to all the multicast group members .
- In recent year , several reliable multicast protocol have been proposed , but they still have some problems .



Problem Definition and Motivation

- It is suitable to multicast over a dense multicast group , but inadequate for sparse group.
- It is inoperable in nonclustered ad hoc networks.
- It restricts that the underlying unicast routing protocol is reliable



Problem Definition and Motivation

- Observing that a certain class of application will be almost infeasible on ad hoc networks without reliable multicasting .
- We need to have a reliable multicast protocol that requires minimal support from underlying network protocols and does not depend on an underlying clustering protocol .



Proposed Solution

- ReMHoc is a distributed receiver-initiated NACK-based reliable multicast protocol .
 - REQUEST
 - REPAIR
 - HeartBeat(HB)



Proposed Solution (REQUEST)

- When a receiver detects a missing DATA packet , it should request its retransmission by multicasting a negative acknowledgement (REQUEST)
- Request implosion
 - A great number of REQUEST are generated to request the retransmission of the same DATA packet
- Request timers
 - Whenever a receiver detects a missing DATA packet , it sets a request timer for a random interval.



Proposed Solution (REPAIR)

- If the member receiving a REQUEST has a cached copy of the requested DATA packet , it can respond by multicasting the cached copy (REPAIR)
- Repair implosion
- Repair timers

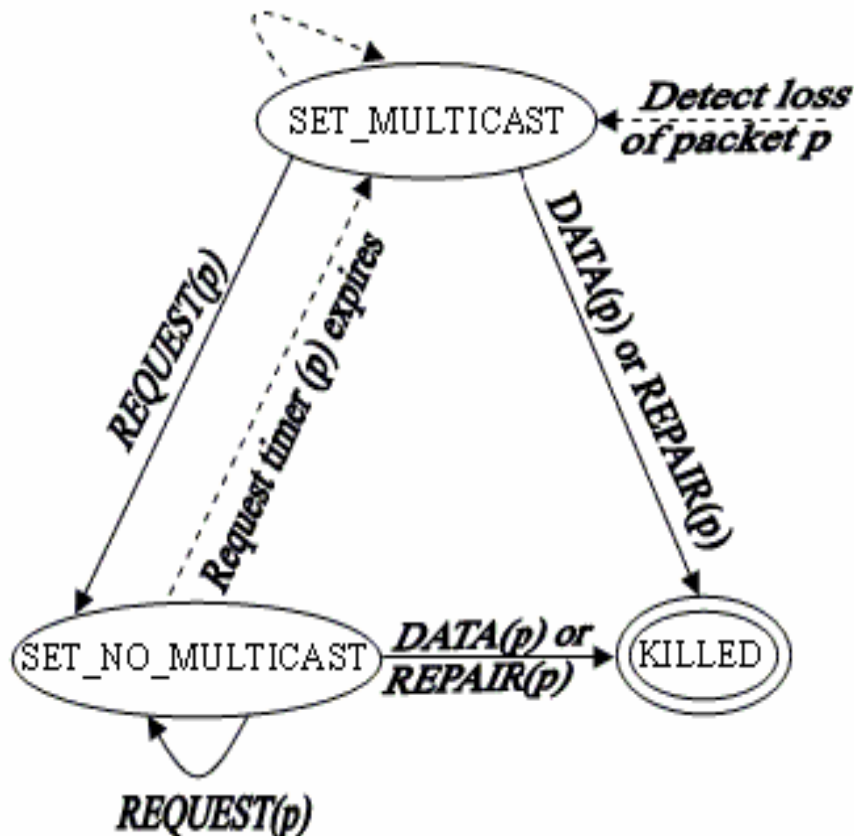


Proposed Solution (HB)

- As long as a receiver has not received that END packet , it sets a heartbeat timer. When the timer expires , a receiver multicasts a heartbeat (HB)
- Any group member which receives an HB packet , tries to find in its cache a copy of any DATA packet whose sequence number is higher than that indicated in HB packet.
- This mechanism ensures that all receivers will receive the last data packet and that receivers which do not receive any data packets for a “long” period can keep pace with other receivers.

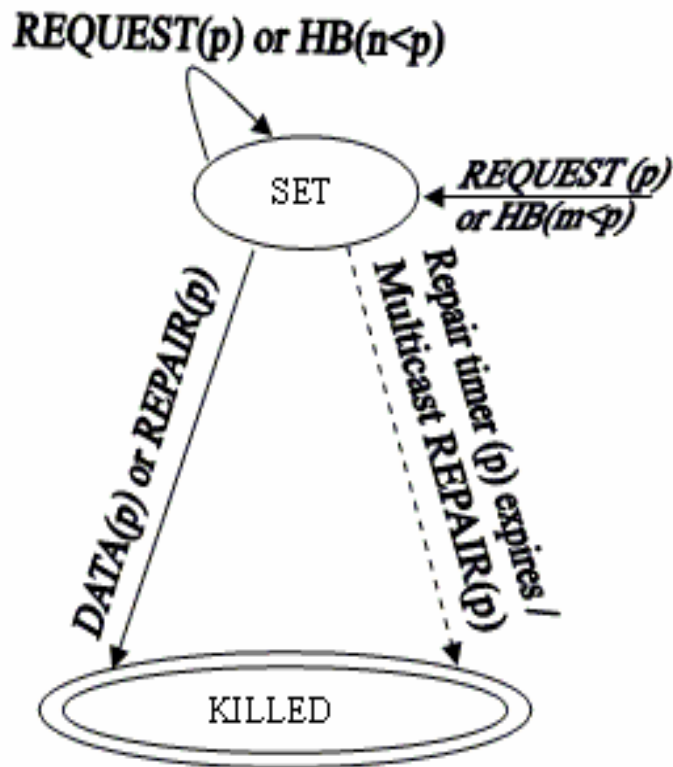
State Transition Diagram (REQUEST)

Request timer (p) expires / Multicast REQUEST(p)



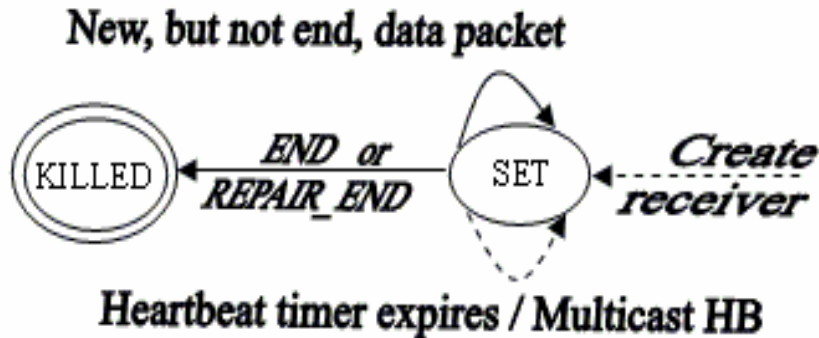
- The request timer interval is made dependent on the hop count between the receiver and source.
- It is desirable to make the request timer interval for a missing packet gets longer if the receiver has already sent REQUESTs for it before.

State Transition Diagram (REPAIR)



- It is desirable to make the one who is likely to time out first be the one that is closest to the requestor.

State Transition Diagram (HB)



- A receiver may not receive new data packets for a long period because the rate of arrival of data packets is low.
- The ReMHoc make the heartbeat timer interval a multiple of the packet interarrival time.



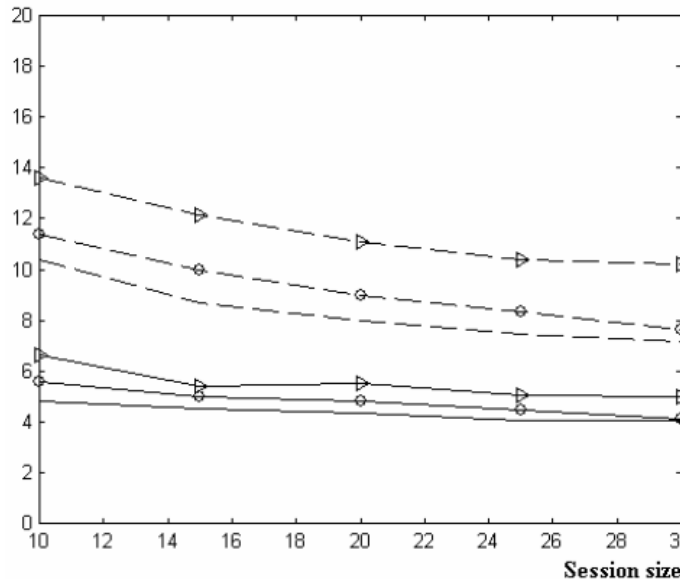
Simulations and Performance Evaluation

- Performance evaluation criteria can be stated as follows
 - Percentage of REQUESTs and HBs
 - Percentage of REPAIRs
 - Average recovery latency
 - Average end-to end delay
 - Overhead percentage

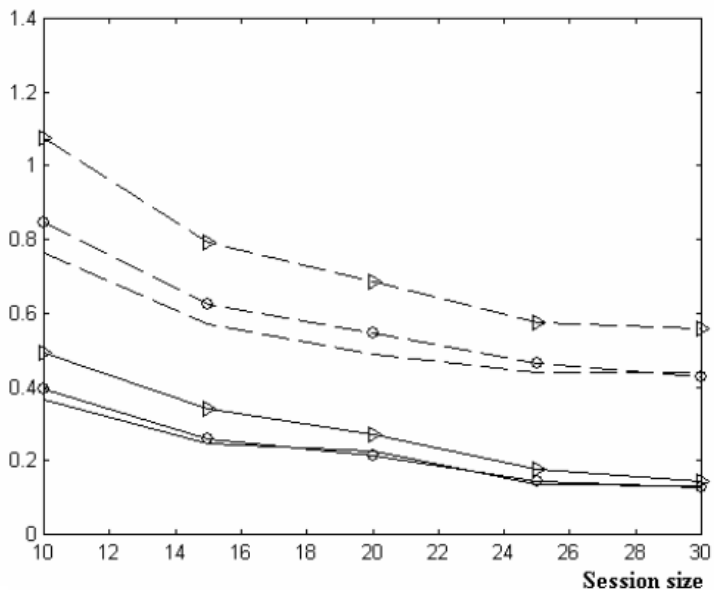
Simulations and Performance Evaluation

■ Effect of mobility and session size

Average Recovery Latency (sec.)



Average end-to-end delay (sec.)

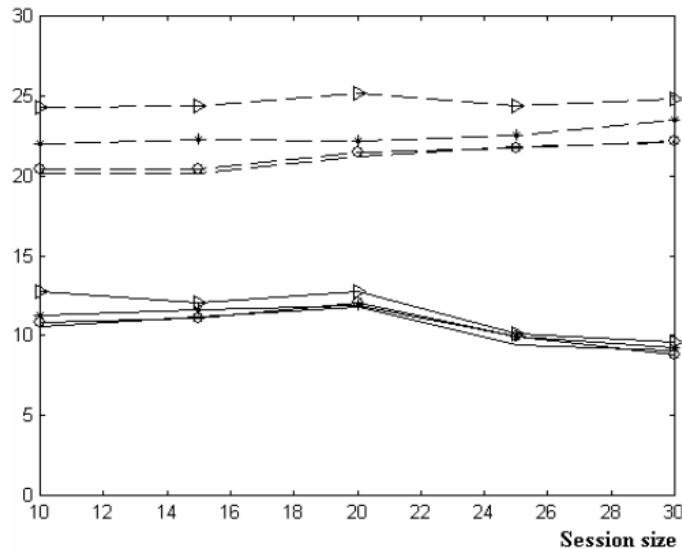


- 0.3 - 0.7 m/sec., Distributed Loss Recovery
- - - 11 - 25 m/sec., Distributed Loss Recovery
- 0.3 - 0.7 m/sec., Semi-distributed Loss Recovery
- -○- - 11 - 25 m/sec., Semi-distributed Loss Recovery
- ▷— 0.3 - 0.7 m/sec., Centralized Loss Recovery
- -▷- - 11 - 25 m/sec., Centralized Loss Recovery

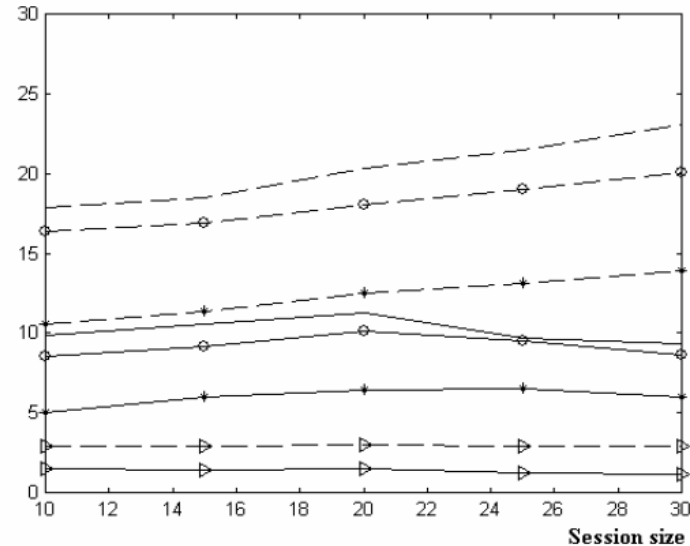
Simulations and Performance Evaluation

■ Effect of the receiver's cache size

Percentage of REQUESTs and HBs



Percentage of REPAIRs



- 0.3 - 0.7 m/sec., Cache_size = 2000 packets
- - - 11 - 25 m/sec., Cache_size = 2000 packets
- 0.3 - 0.7 m/sec., Cache_size = 50 packets
- -○- - 11 - 25 m/sec., Cache_size = 50 packets
- *— 0.3 - 0.7 m/sec., Cache_size = 25 packets
- -*-- 11 - 25 m/sec., Cache_size = 25 packets
- △— 0.3 - 0.7 m/sec., Cache_size = zero



Conclusions and Discussions

- ReMHoc is scalable with the number of multicast group members (session size)
- The centralized loss recovery has higher overhead percentage , higher average recovery latency and higher average end-to-end delay than distributed loss recovery.
- Can this method extend to multi-channel ad hoc networks?