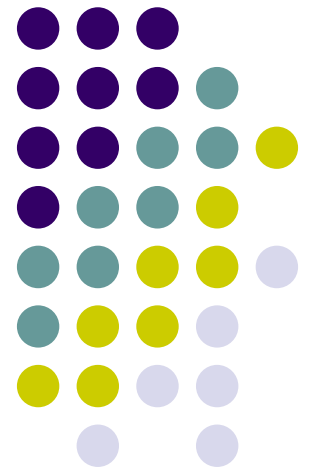


Efficient and Scalable Query Routing for Unstructured Peer-to-Peer Networks

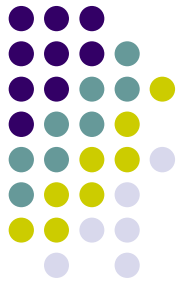
Presented by 曾胤燁

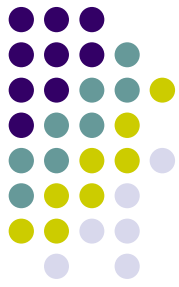
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Outline

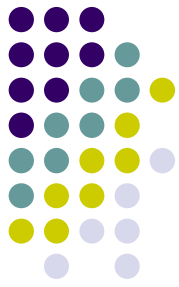
- Introduction
- Design
- Experiment
- Conclusion





Introduction

- Design an efficient query routing mechanism for unstructured peer-to-peer networks.
- To build routing tables at nodes.
 - Use a novel data structure EDBF
- Scalable Query Routing (SQR) mechanism
 - Maximum probability



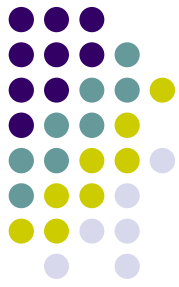
EDBF

- Exponential Decay Bloom Filter
- EDBF is an extension of the traditional Bloom filter to encode the probabilistic forwarding table.

Bloom filter

- Given a query, EDBF simply return indicator $\theta(x)$, the number of 1's in the filter.
 $\theta(x) / k \rightarrow$ probability

EDBF

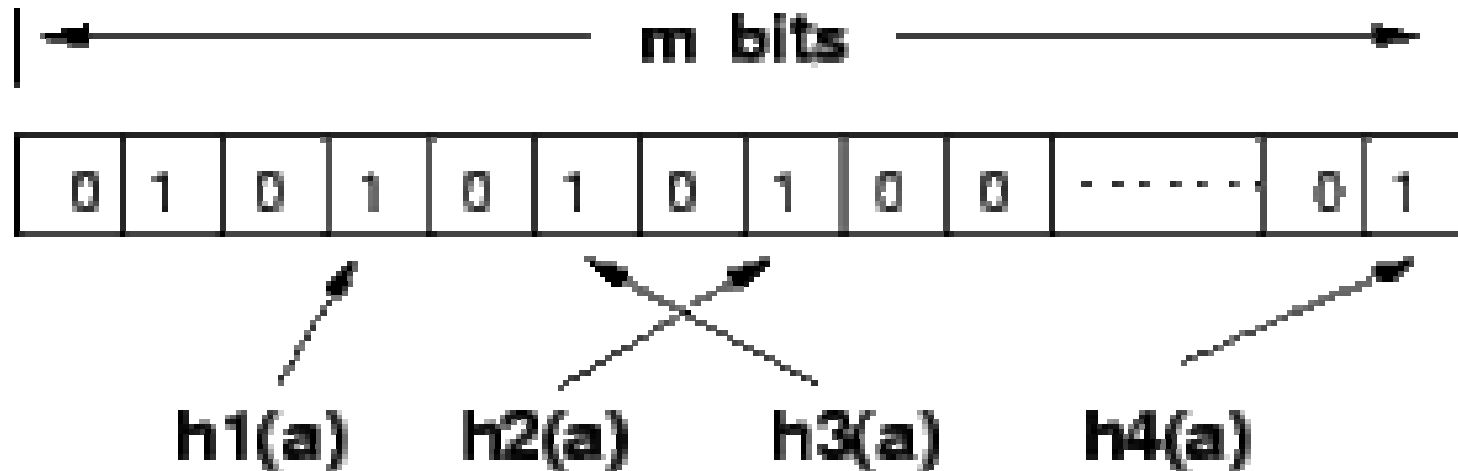


Bloom filter

- A data structure for approximately answering set membership questions.
- An array of bits A initialized to all 0.
- Fixed number (k) of hash functions

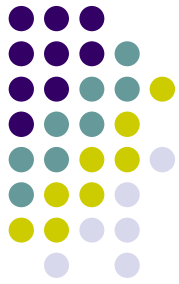


- When a inserts, $h_i(a)$, $i = 1, 2, \dots, k$ are set to 1.



- A query for y returns a yes if and only if all $h_i(y)$, $i = 1, 2, \dots, k$ are 1.

EXAMPLE



X Y Z

$H_1(X)$

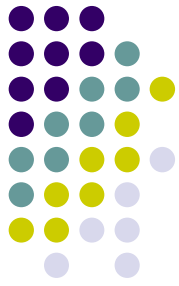
$H_2(X)$

$H_3(X)$

$H_4(X)$

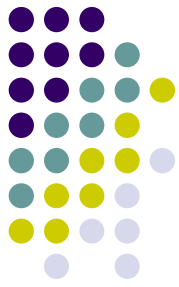


Q (Y): $H_1(Y)$ $H_2(Y)$ $H_3(Y)$ $H_4(Y)$ -> YES Q (W): $H_1(W)$ $H_2(W)$ $H_3(W)$ $H_4(W)$ -> NO



優點

- 可將多筆資料紀錄於同一array
- 利用多個hash function可減少碰撞的機會



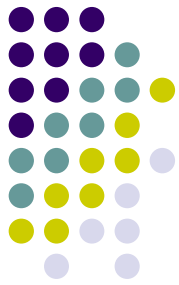
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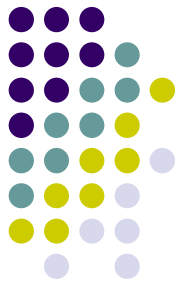
EDBF



1 0 1 1 1 0 0 0 1 0 1 0 1 1 1 1 0 1 0 1

Q (Y): $H_1(Y)$ $H_2(Y)$ $H_3(Y)$ $H_4(Y)$ -> 4 Q (W): $H_1(W)$ $H_2(W)$ $H_3(W)$ $H_4(W)$ -> 2

EDBF



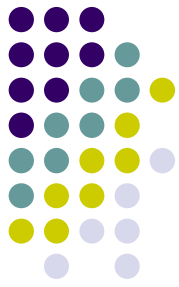
- Exponential Decay Bloom Filter
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Bloom filter

- Given a query, EDBF simply return indicator $\theta(x)$, the number of 1's in the filter.

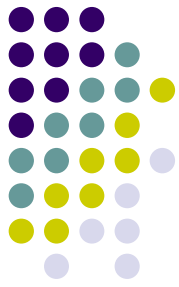
$\theta(x) / k \rightarrow$ probability

EDBF

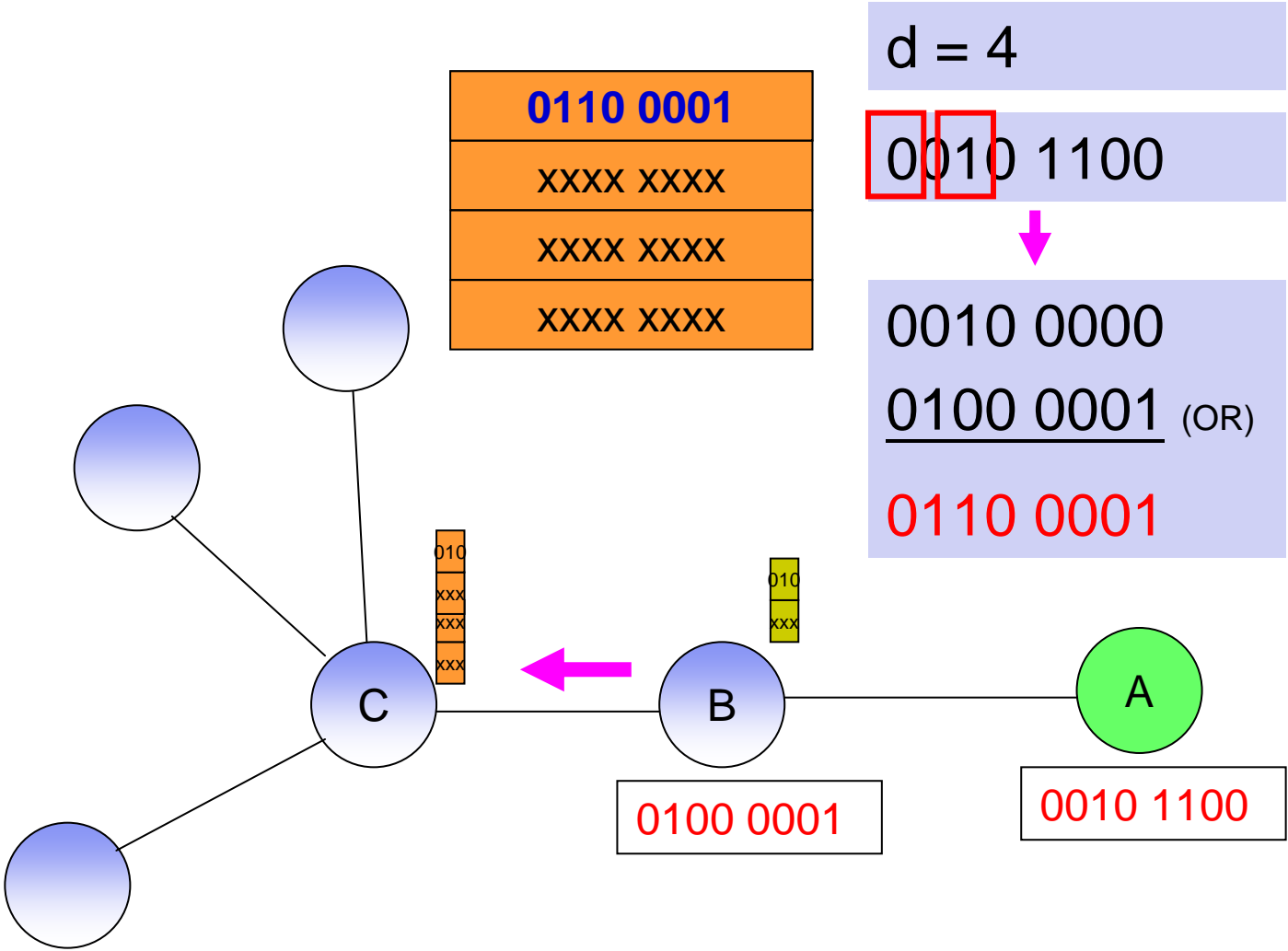


- The number of bits decays exponentially with its distance from the node where the object x is stored. → **EDBF**

Maintenance of routing tables



- A node's (with degree j) routing table consists of j EDBF arrays.
- Maintain consistency by periodic update
- Information regarding content x will decay with distance



0110 0001
XXXX XXXX
XXXX XXXX
XXXX XXXX

d = 4

0010 1100



0010 0000
0100 0001 (OR)
0110 0001

0100 0001

0010 1100

0010 1100
XXXX XXXX

010
xxx
xxx
xxx

010
xxx



Algorithms for creating updates

```
Create Local EDBF (given local content  $X$ ):  
  // Populate local EDBF  $A$ .
```

1. $\forall x \in X$
2. Set bits $A[h_1(x)], \dots, A[h_k(x)]$ to 1;

```
Create Update (for neighbor  $j$ ):
```

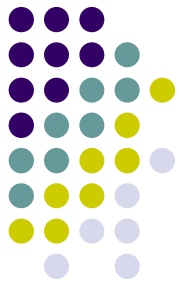
```
  // Copy all the bits from the local EDBF  $A$  into  
  // the update  $U_j$ .
```

1. $U_j \leftarrow A$;
 // Decay the information received from all neighbors
 // other than j by a factor of d , and add the
 // surviving bits to U_j .

2. $\forall i \in neighbor_list, i \neq j$
3. $\forall r \in \{1, \dots, m\}$
4. $if(A_i[r] = 1)$
5. $with\ probability\ 1/d, U_j[r] \leftarrow 1;$
6. Return U_j ;

Query Forwarding

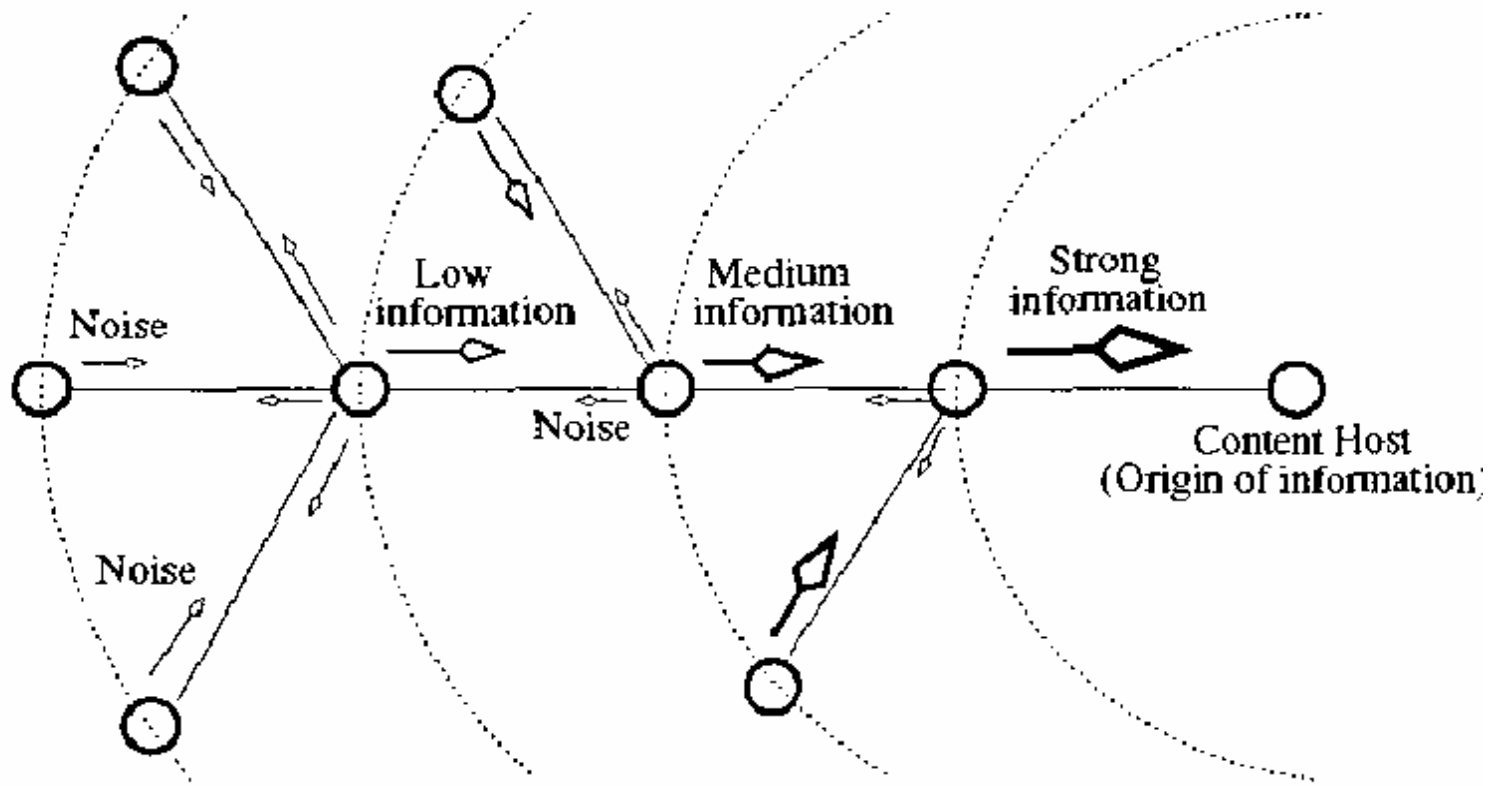
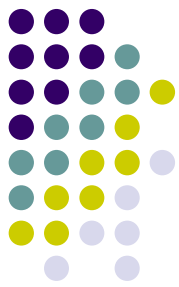
-Scalable Query Routing (SQR) mechanism

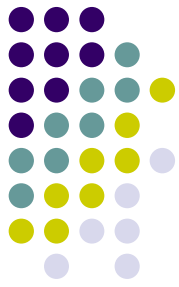


- The query is forwarded to the neighbor with the highest indicator value $\theta(x)$.

```
Forward Query (given query  $Y$ ):  
    // Forward previously seen queries to neighbor  $i$ ,  
    // chosen randomly from neighbor_list.  
1. if ( Seen Query( $Y$ ) )  
2.     Deliver Query( $Y, i$ );  
3. else  
    //Forward previously unseen queries to the neighbor  
    // with the maximum information about this query.  
4.      $\Theta = \text{Lookup}(Y)$ ;  
5.     Pick  $i$  such that  $\theta_i = \max(\Theta)$ ;  
6.     Deliver Query( $Y, i$ );
```

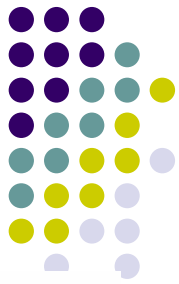
```
Lookup (given query  $Y$ ):  
1.  $\forall i \in \text{neighbor\_list}$   
2.      $\forall q \in \{1, \dots, k\}$   
3.          $\theta_i += A_i[h_q(y)]$ ;  
4. Return  $\Theta$ ;           /* $\Theta = \{\theta_i\}$ */
```

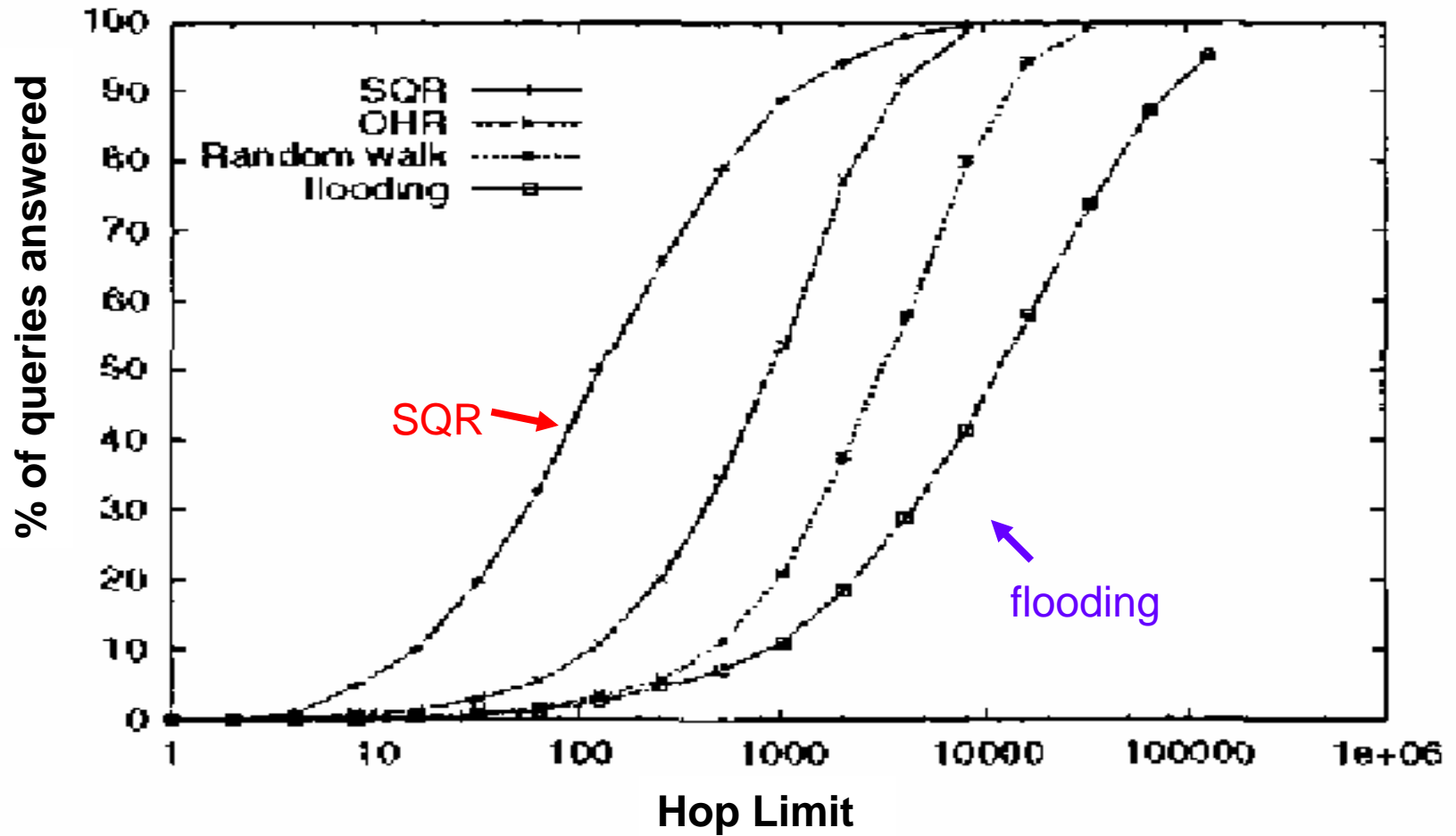


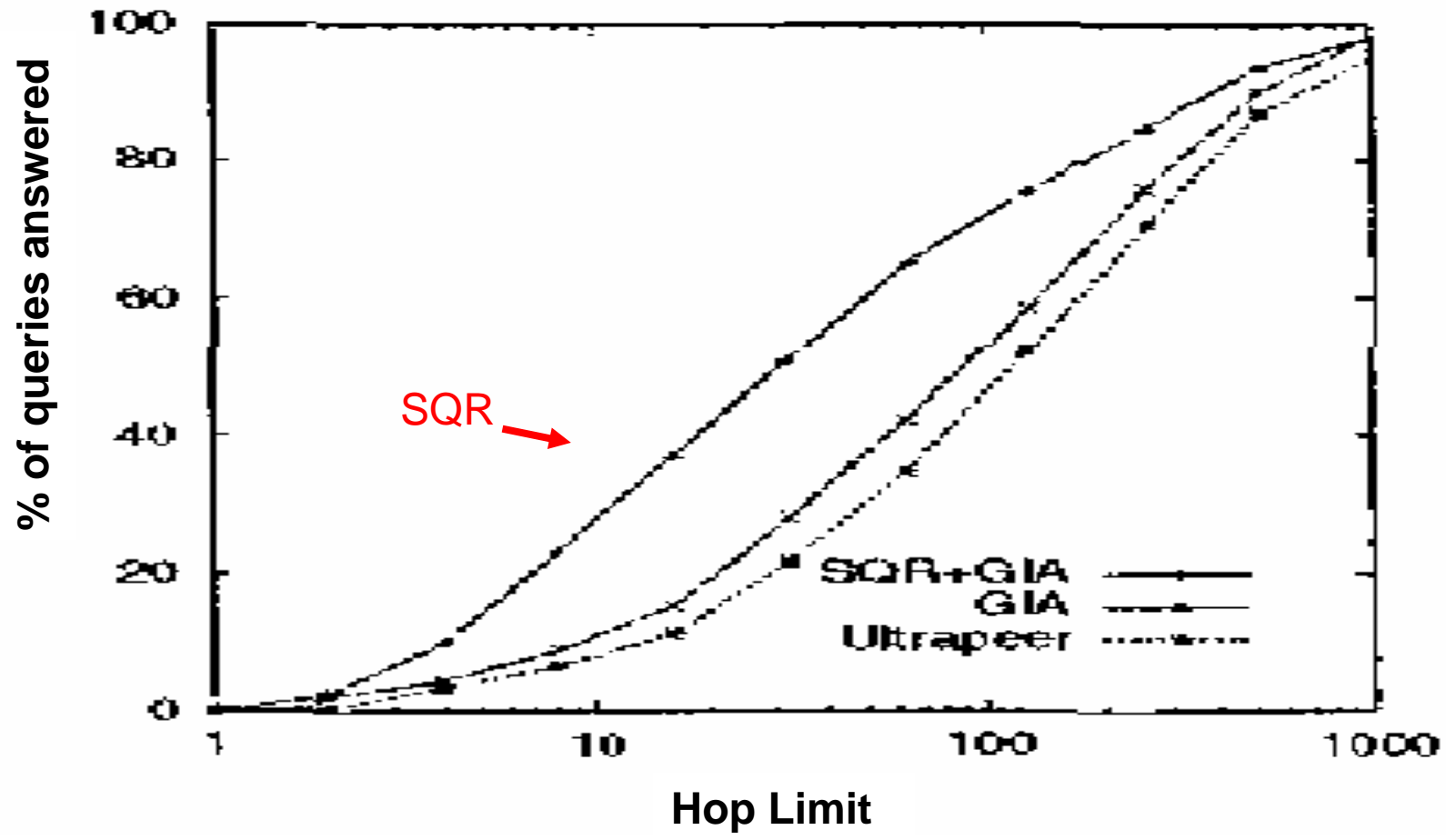
優點

- 使用**BF**，可使每一**node**所儲存的資料量不至於太大
- 利用**information decay**，可使在每一**peer**中的資料重要性被區分出來，距離較遠的資訊比較不會干擾到距離較近較重要的資訊-避免**Noise**

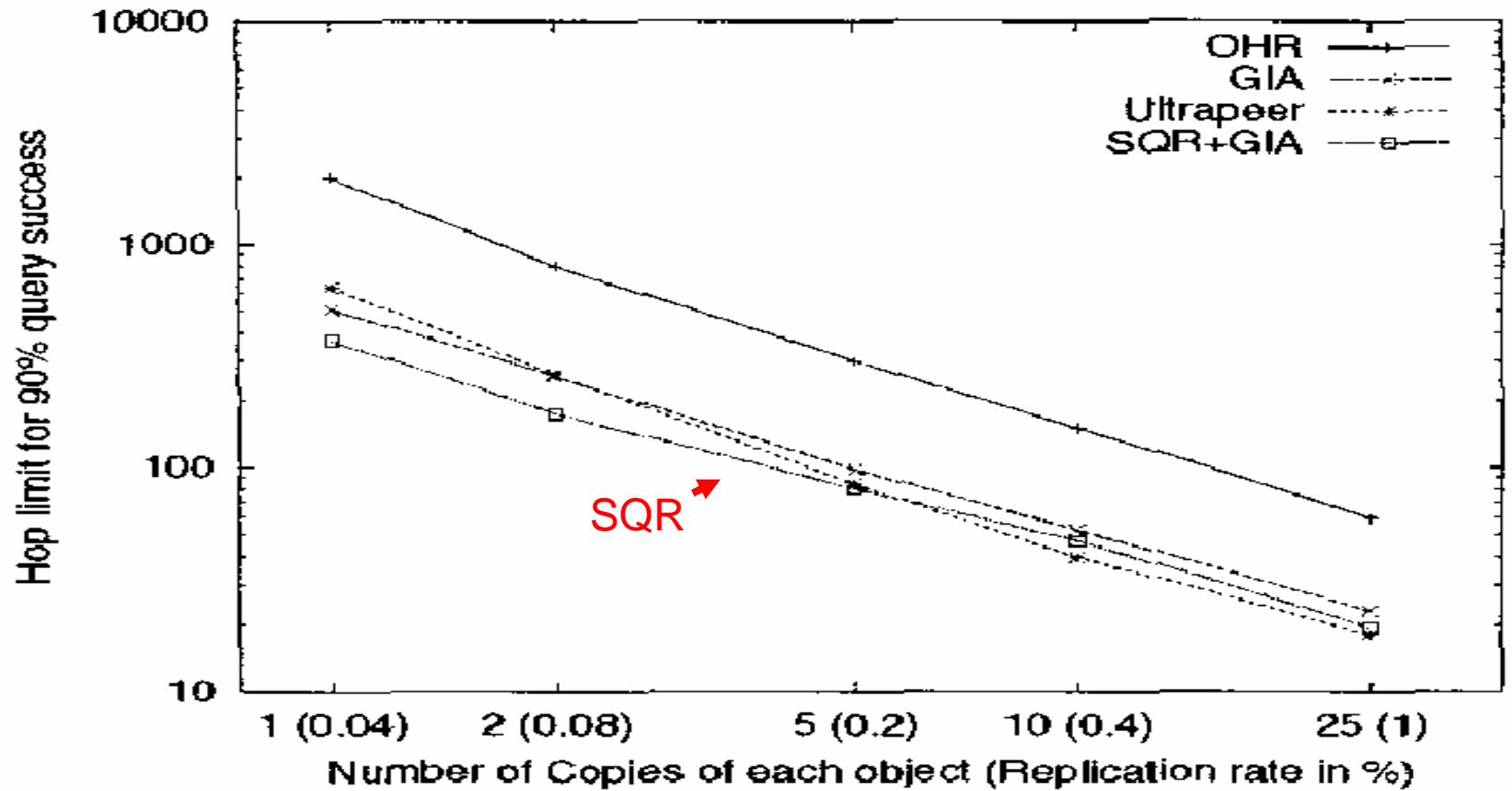
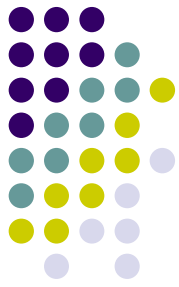


Query response rate





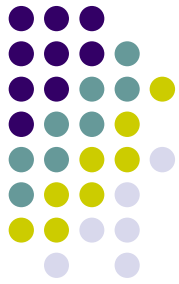
Impact of Replication





Conclusions

- In this paper, we have explored the approach of spreading probabilistic information about the location of hosted content in its neighborhood.
- Simulation based comparison with various query routing mechanisms establish the performance advantages of SQR



Thank you😊