

Distributed Navigation Algorithms for Sensor Networks

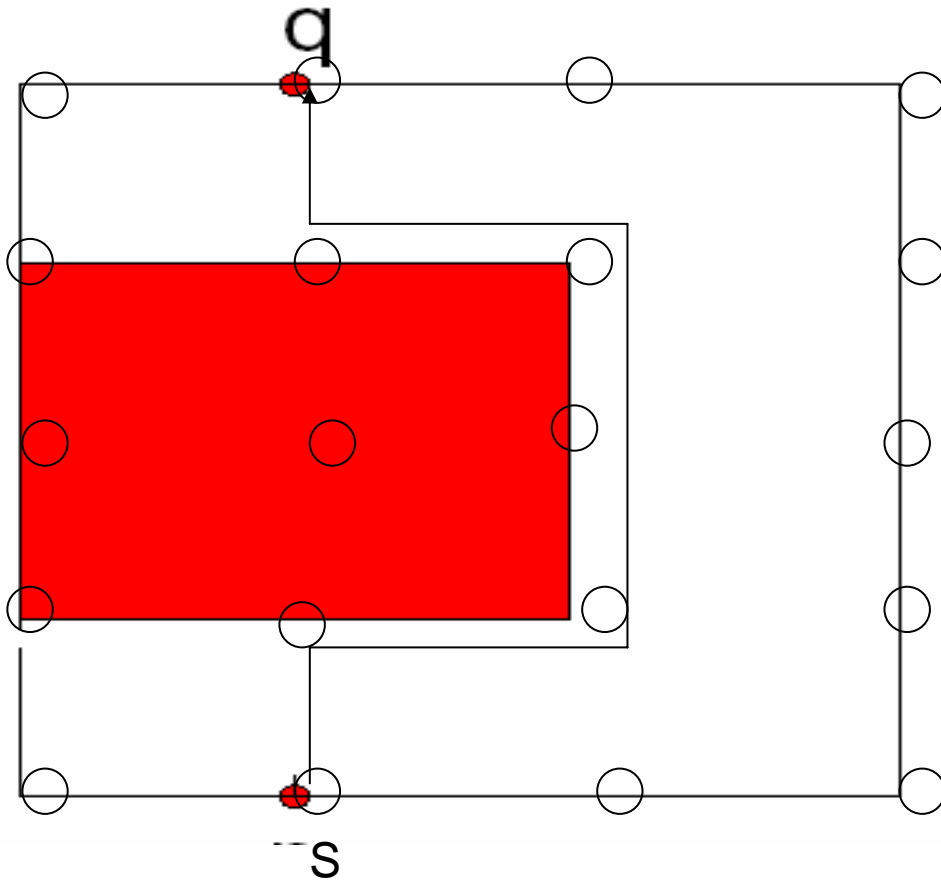
INFOCOM 2006

Speaker: Cheng-Han Wu

Outline

- Introduction
- The main advantage of this paper
- The main idea of this paper
- Simulation
- Conclusion

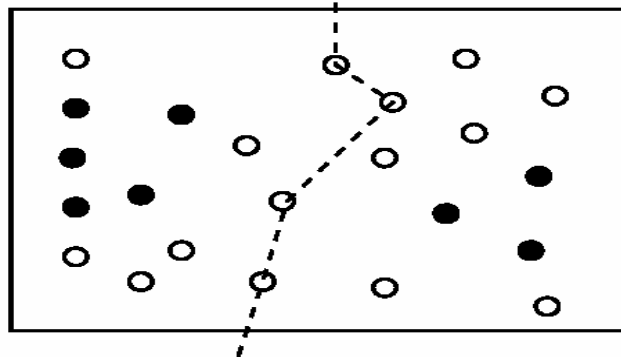
How to navigate ?



q: destination
s: start point

The Disadvantage of the past Method

- This scheme does not scale well due to a very high communication cost.
- In a network of n nodes, the number of total packet transmissions required for the past method to find the safe path is $O(n)$.



The Advantage of this Paper

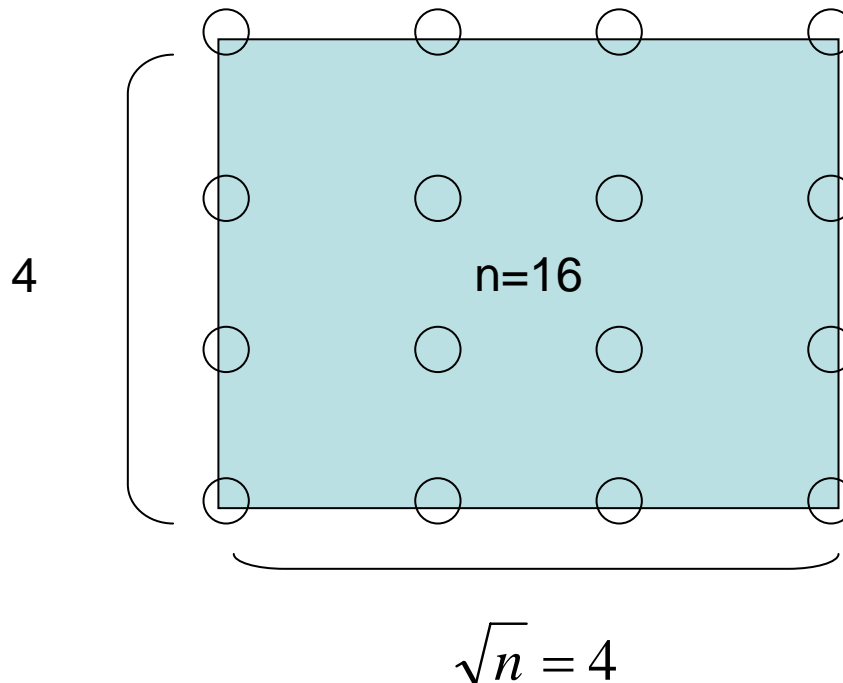
- The communication cost of discovering the shortest path is $O(n^{1/2+e})$, for any e such that $0 < e < 1/2$.

Assumptions

- Each sensor knows its geographic location.
- The operational environment is assumed to have no large holes in the coverage by sensors.

Related Works

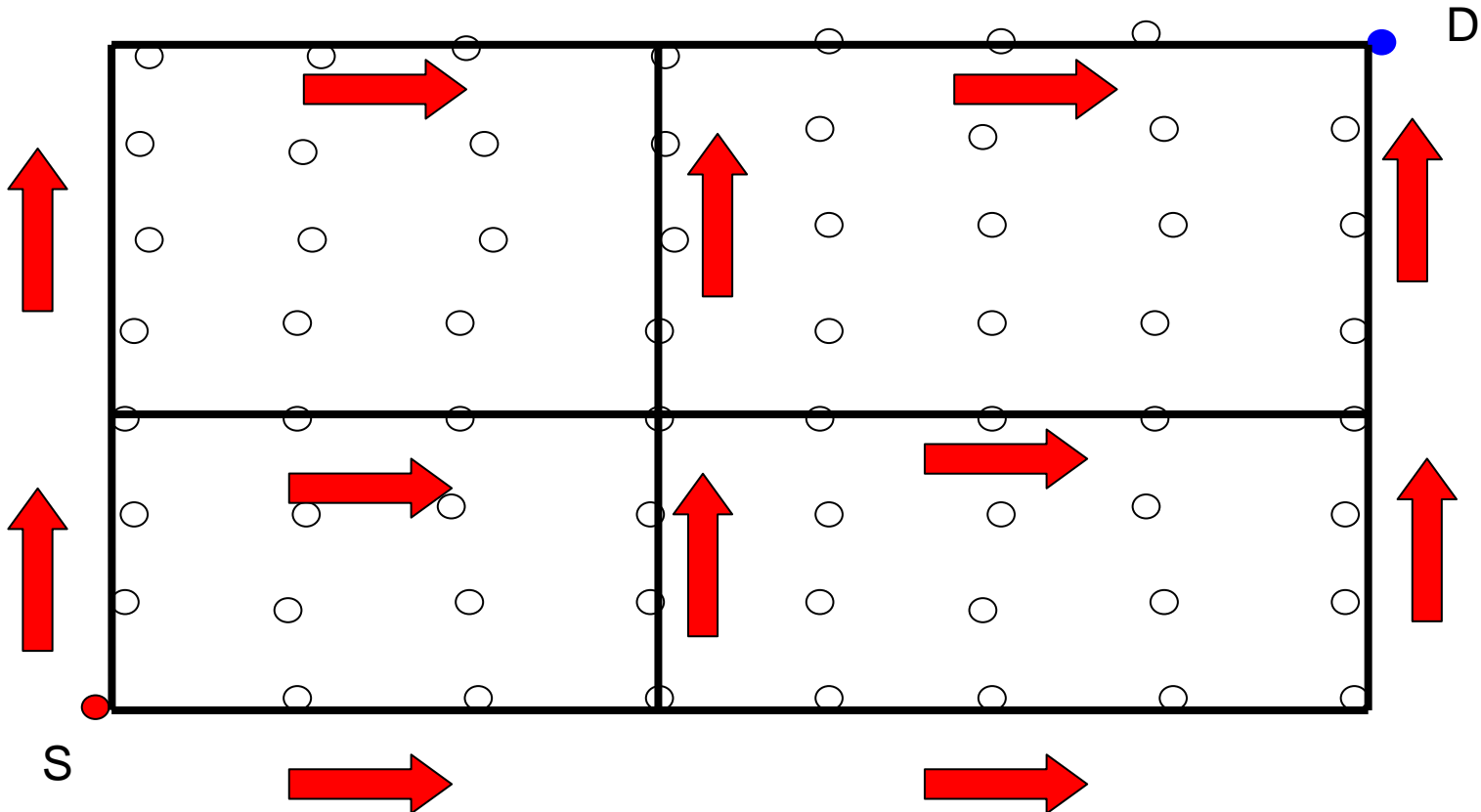
- We assume that n sensor nodes are placed uniformly in a square area.
- We choose units of length such that the size of area is $\sqrt{n} * \sqrt{n}$, i.e. on the average every unit area contains one single sensor.



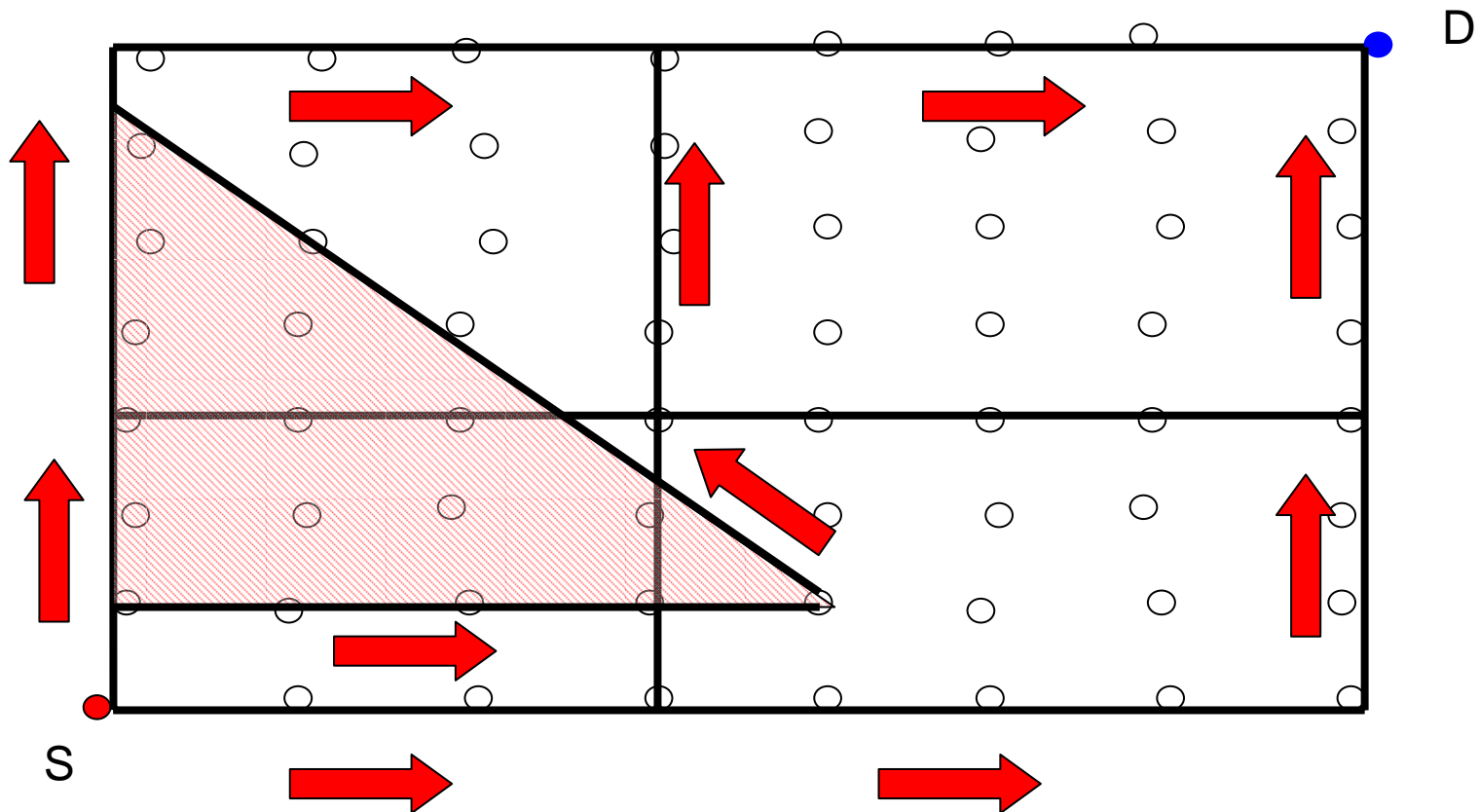
The Idea of this Paper(1/2)

Navigation using uniform Skeleton Graph

- Step1: Establishing the **Skeleton Graph**
- Step2: Using **shortest path algorithm** to find the safe path

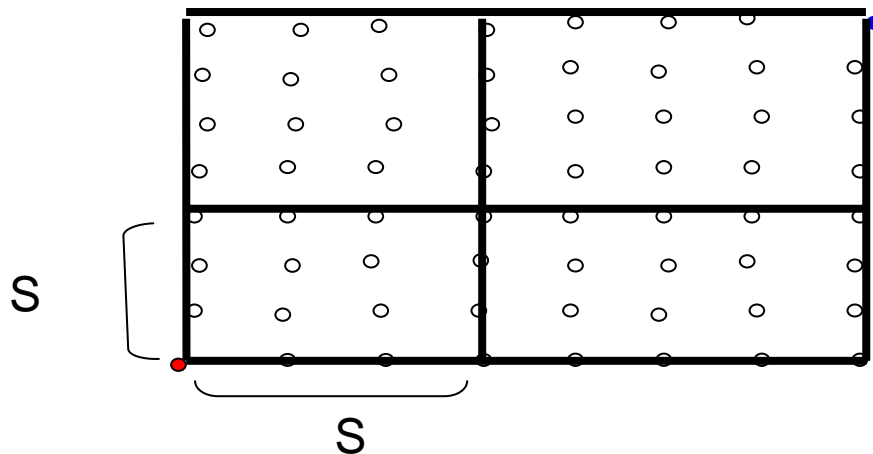


The Idea of this Paper(2/2)



How to construct the Skeleton Graph

- All nodes know that the side length of the grid is S .
- Because each sensor knows its location, it can independently decide if they are in the skeleton graph.

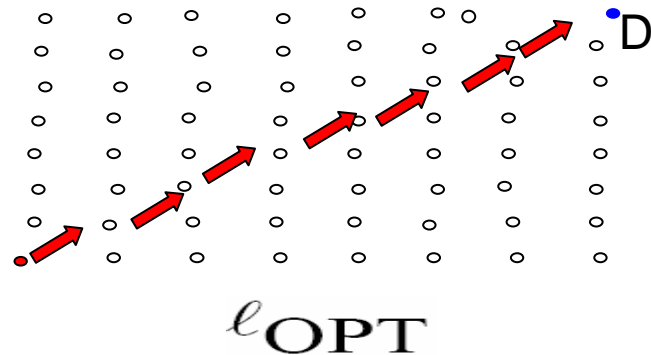
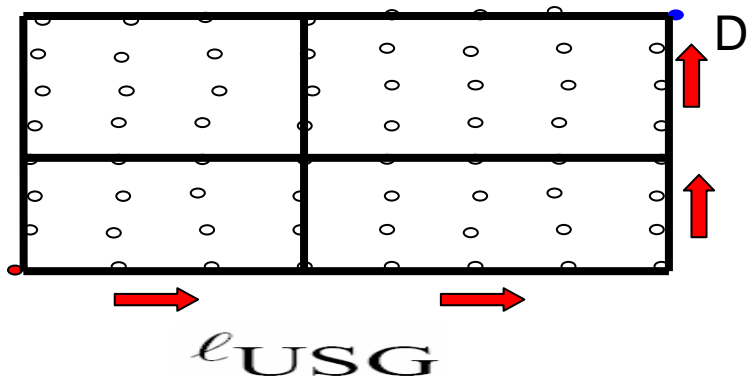


The properties of Skeleton Graph

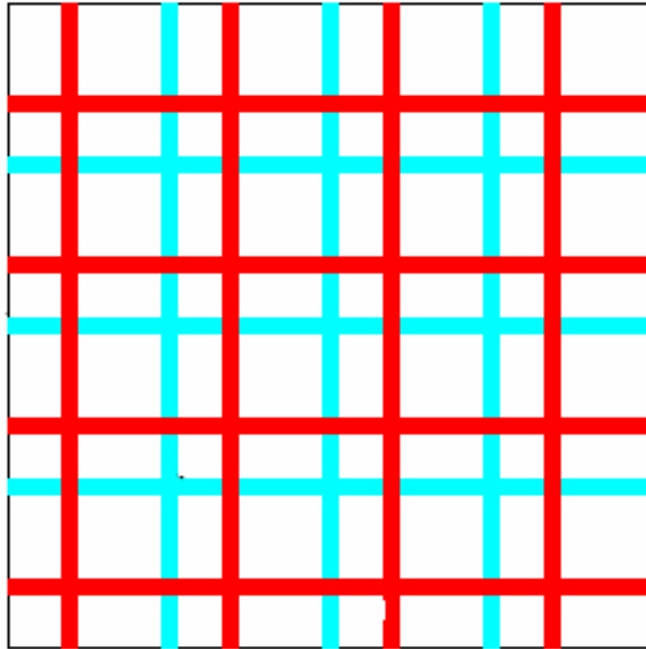
Theorem 1: The communication cost of discovering the shortest path in the uniform skeleton graph is $\mathcal{O}(n^{1/2+\epsilon})$, for any ϵ such that $0 < \epsilon < 1/2$.

Theorem 2: For a path joining any two points located on the streets in uniform skeleton graph,

$$\ell_{\text{USG}} / \ell_{\text{OPT}} \leq 2(1 + c), \quad C \text{ is a constant}$$

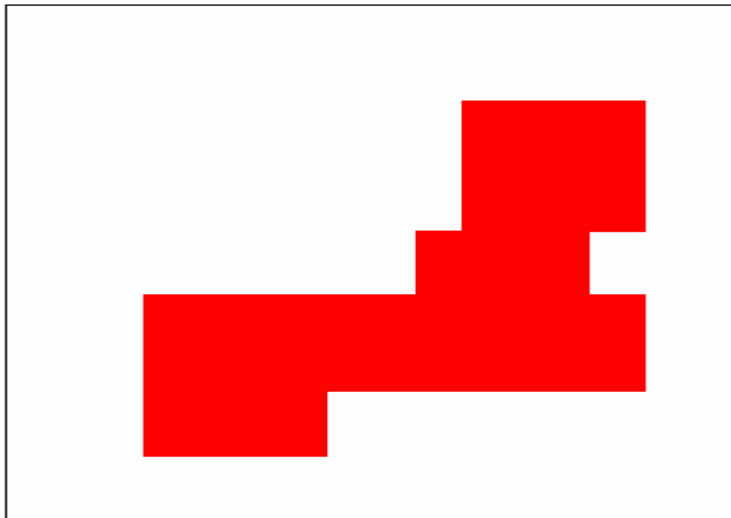


Load Balancing by shifting

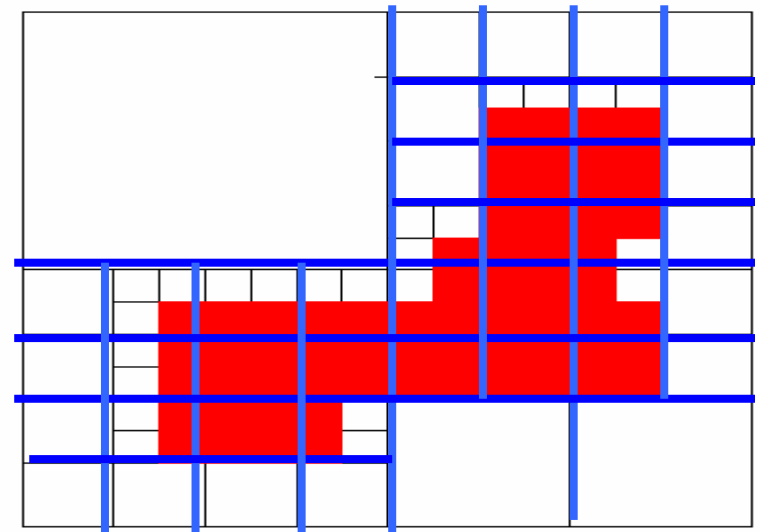


Navigation Using Adaptive Skeleton Graph

- We assume that danger zone boundary is axis aligned.
- The process of generating adaptive skeleton graph is very similar to quadtree mesh generation.

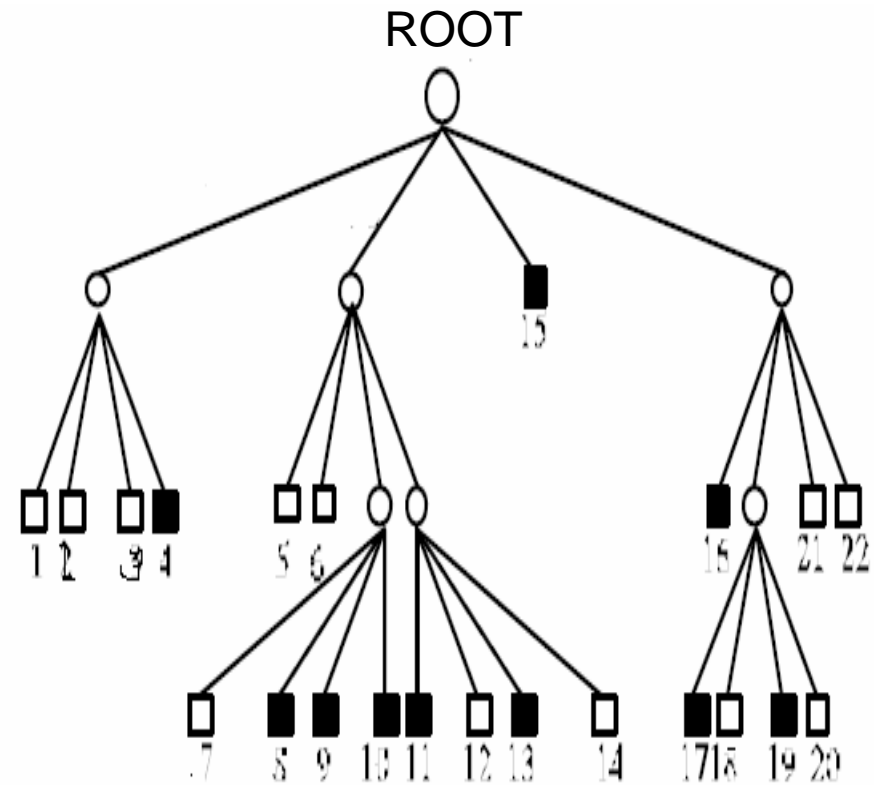
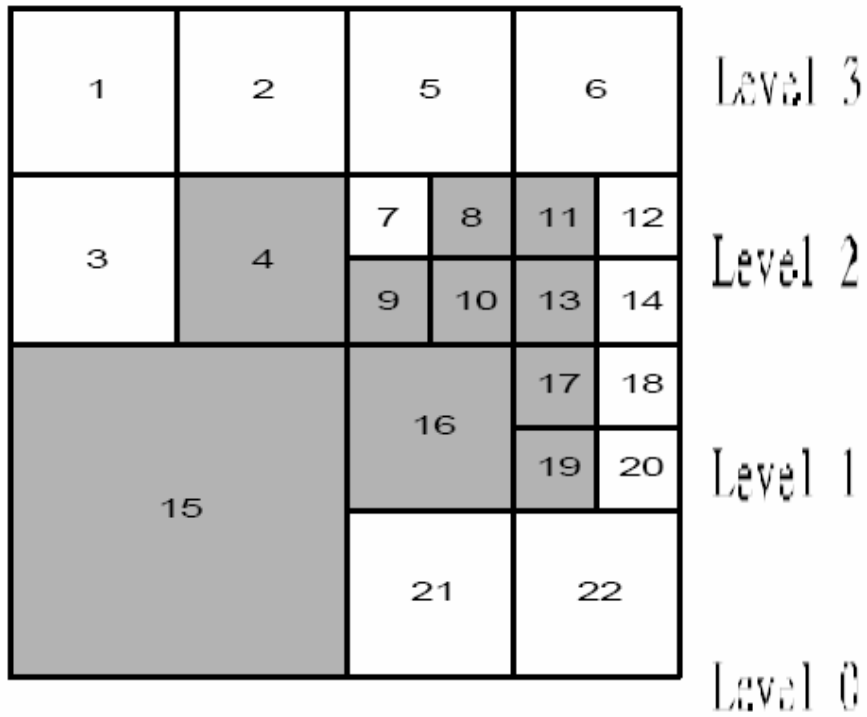


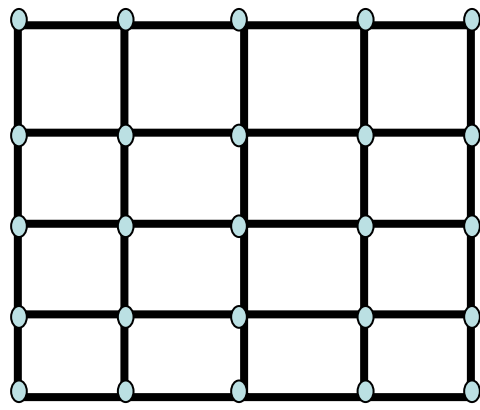
(a)



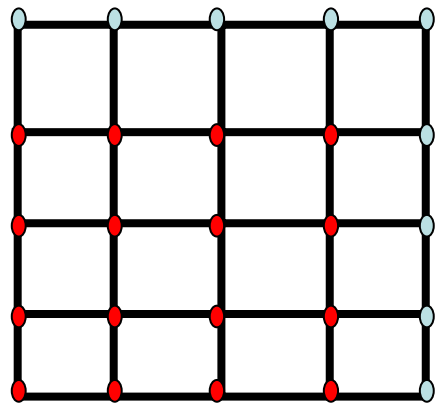
(b)

Quadtree



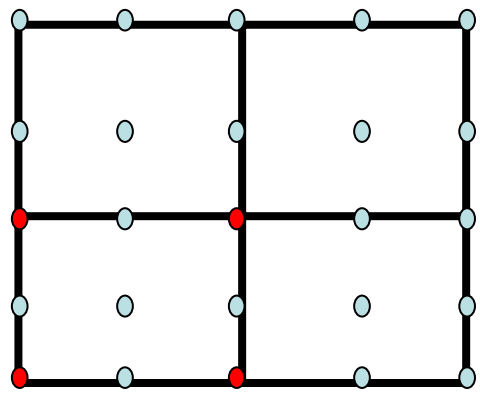


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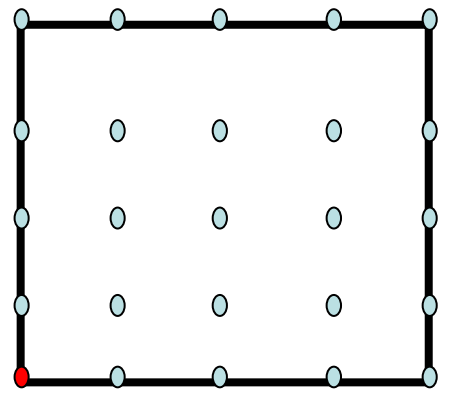
level 0

+

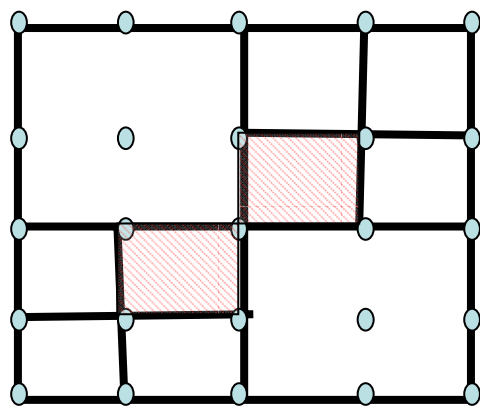
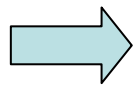
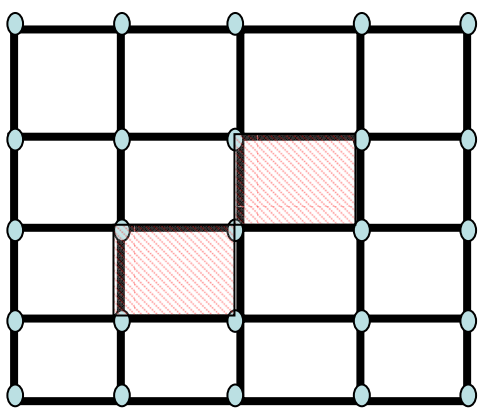


level 1

+



root



The Properties of Adaptive Skeleton Graph

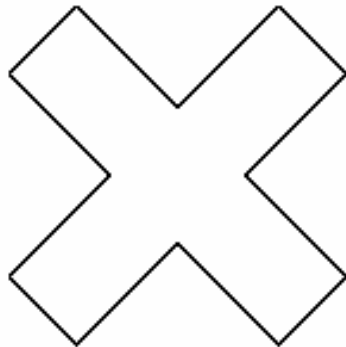
Theorem 4: The communication cost of discovering the shortest path in the adaptive skeleton graph is $\mathcal{O}(n^{1/2} \log n)$.

Theorem 5: For a path joining any two points located on the streets in the adaptive skeleton graph,

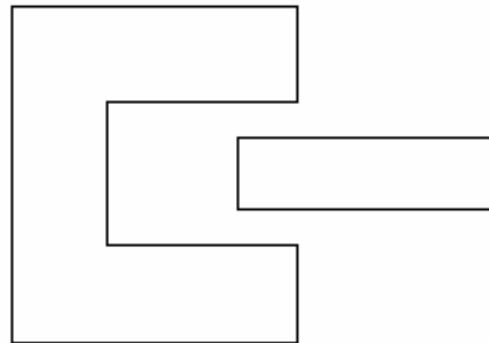
$$\ell_{\text{ASG}} / \ell_{\text{OPT}} \leq 2.$$

Simulation

- Placing n sensor nodes in a $\sqrt{n} * \sqrt{n}$ area.
- Simple and complex danger zone shapes used to test shortest path algorithm.

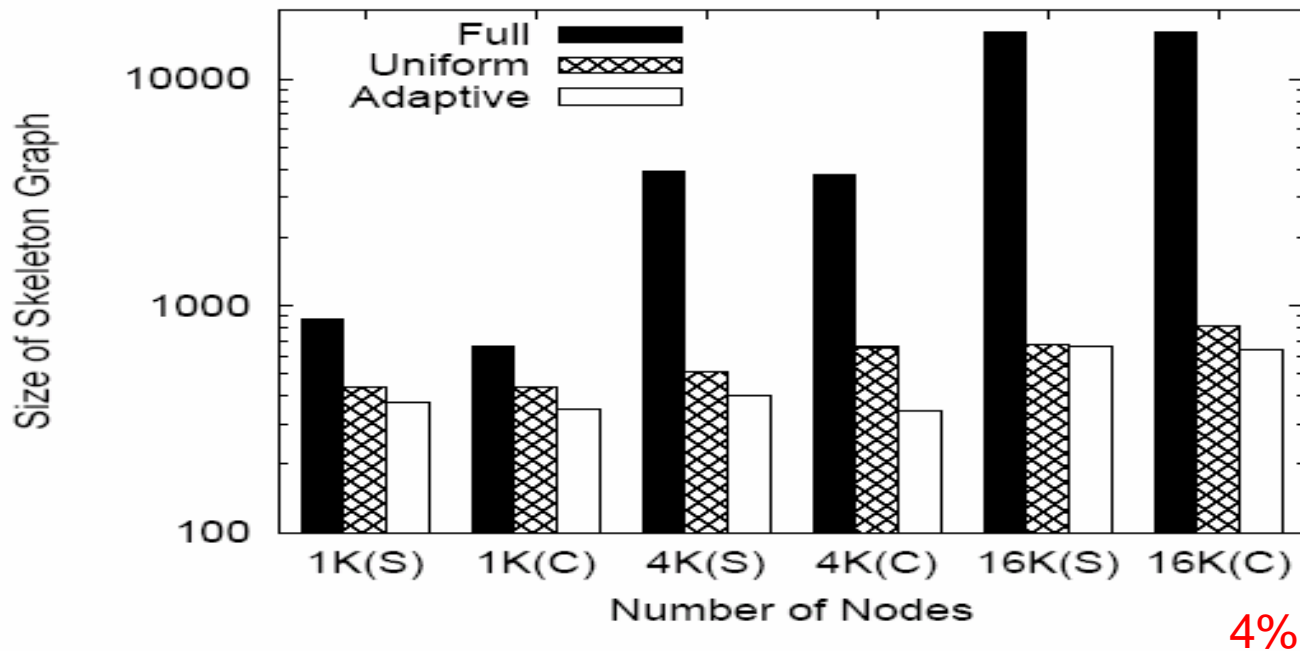


Simple



Complex

Simulation



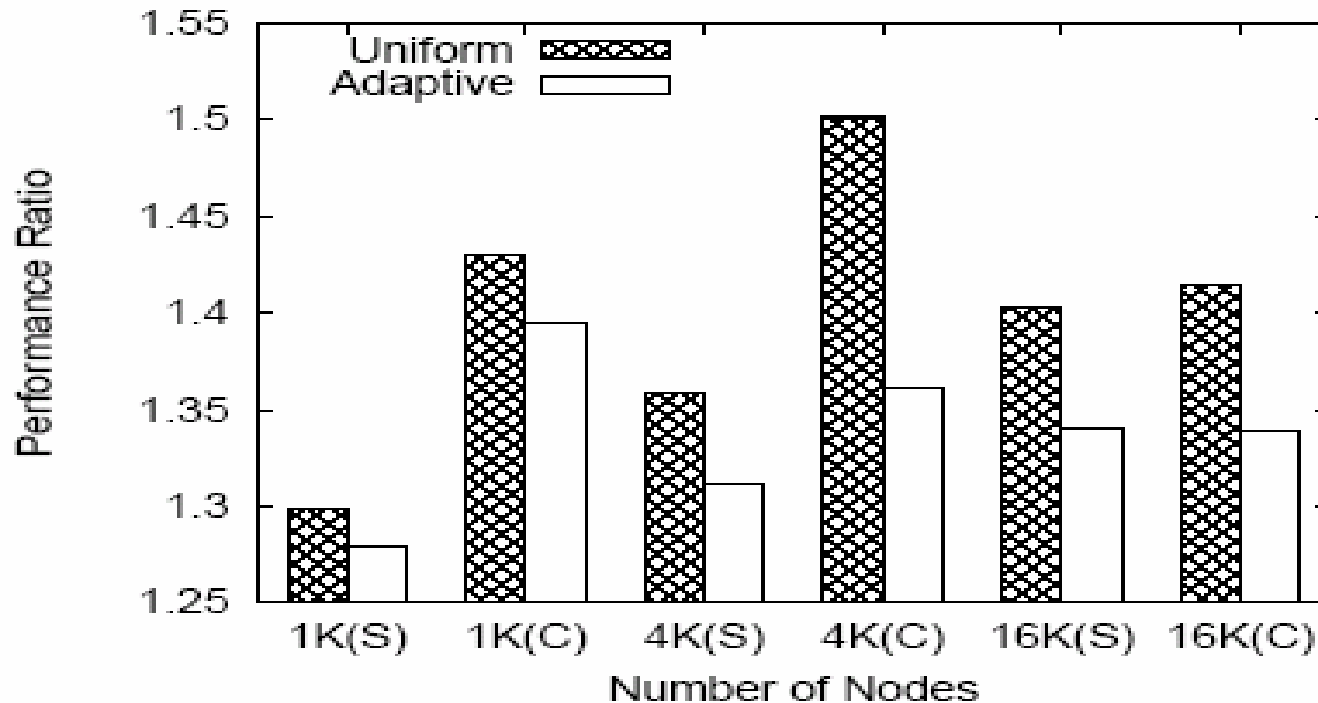
1K(S) means 1024 nodes with simple danger zone.

1K(C) means 1024 nodes with complex danger zone.

The path length performance ratio = l_{SG}/l_{OPT}

l_{SG} :the lengths of the approximate skeleton graph path

l_{OPT} :the lengths of the optimal path



Conclusions

- The problem of finding shortest path on a sensor network can be solved approximately with low communication cost using skeleton graphs.
- The uniform skeleton graph is superior to the adaptive graph in terms of its simplicity and load distribution.