Managing a Broadcast Infrastructure in Ad Hoc Networks in Presence of Mobility : A New Algorithmic Framework **IEEE VTC 2007** Presented by Ming-Chieh Li November 15, 2007

## Outline

- Introduction
- Related works
- Definitions and Preliminaries
- The algorithmic framework
- Simulation results
- Conclusions

# Introduction(1/2)

- In mobile wireless ad hoc networks, it is essential to have a broadcast infrastructure for disseminating data of broadcast nature and control information.
- Backbone nodes forward broadcast messages for others. Backbone formation is tightly related to the concept of connected dominating set (CDS).

Introduction(2/2)

 Backbone management in presence of mobility has been studied in a less extent.

In mobile networks, backbone update involves a tradeoff between reliability and energy efficiency.

# Related works(1/2)

In [4], they showed that mobility is the major threat to efficient self-pruning protocols.

In [18], the authors proposed to use two transmission ranges (power levels), the longer one for data transmission, and the shorter one for determining neighborhood, in order to improve backbone reliability.

# Related works(2/2)

In [8], at every node, its decisions of joining and leaving the backbone are made based on 2-hop topology information, which is exchanged periodically between nodes.

# The framework's features

- Nodes do not need to acquire topology information other than its local one-hop environment.
- No location information is required.
- Decisions of joining and leaving the backbone are made individually.

#### Definitions and Preliminaries

- The networks are composed of
  - Backbone nodes (BNs) &
  - Non-backbone nodes (NBNs)
- A backbone node has two possible states
  Regular state (RBN)
  Pruning state (PBN)

#### Definitions and Preliminaries

- An NBN may change its state to RBN if some of its neighbors cannot reach each other. A BN *u* provides reachability to its neighbors.
- If these neighbors may be able to reach each other via BNs other than u, u can prune itself and become an NBN.

# An example of changing states



- Each node broadcasts probe messages regularly via the current backbone.
  - The probe message contains
    - Source
    - TimeStamp
    - TTL
    - Attribute => true or false
- When a BN received the message
  - □ If it is PBN, rebroadcast the message with attribute = true
  - Else, rebroadcast the message



Fig. 1. Scenarios of backbone connectivity between two nodes

- Every node has a database containing sources (within k hops) from which probe messages have been received. A node keeps at most two entries (of different attributes) for any source.
- Every node shares its reachability database with one-hop neighbors.

### An example of redundant PBNs



Node b and c will become RBNs.

An RBN enters the PBN state from time to time in order to examine whether or not it can prune itself from the backbone.

So node b or c will prune itself

### Simulation results

- Algorithm parameters
  - 50 mobile nodes
  - 750m x 700m
  - Each node is moving following the random waypoint model
  - Initial TTL of a probe message = 4
  - Probing interval = 1 second
  - Database exchange interval = 1 second

#### Simulation results

- Ratio 1 is the proportion of node pairs that are connected by the backbone in relation to the total number of node pairs computed as 1/2 |V |(|V | -1).
- Ratio 2 shows the degree of connectivity of the underlying graph G, that is, the number of connected node pairs in relation to 1/2 |V |(|V | -1).

#### Low-density scenario



#### High-density scenario



### Conclusions

We have presented an algorithmic framework for managing a broadcast backbone in mobile ad hoc networks.

The simulation results show that the proposed framework is effective in adapting backbone to topology change.