

# Weak State Routing for Large Scale Dynamic Networks

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# Outline

- Introduction
- Assumptions
- Weak State Routing
- Simulation
- Conclusions

# Introduction

- In a large scale dynamic networks (large scale MANET), traditional routing protocols require high overhead to keep some information.
- This paper proposes Weak State Routing to solve this problem.
- The state information is weak, i.e. interpreted not as absolute truth, but as probabilistic hints.
- Nodes only have partial information about the region a destination node is likely to be.

# Assumptions

- All nodes know their positions on a 2-D plane(GPS).
- By using single hop beacon messages, the nodes also know their neighbors and their positions.
- All nodes have uniform omnidirectional antennas.
- There is no ACK after transmitting.

# Weak State Routing

- Weak State Routing (WSR) is based on link-state geographic routing protocol.
- WSR doesn't maintain the state of each node, since there are many and large size beacon messages.
- It only maintains the state of each region.

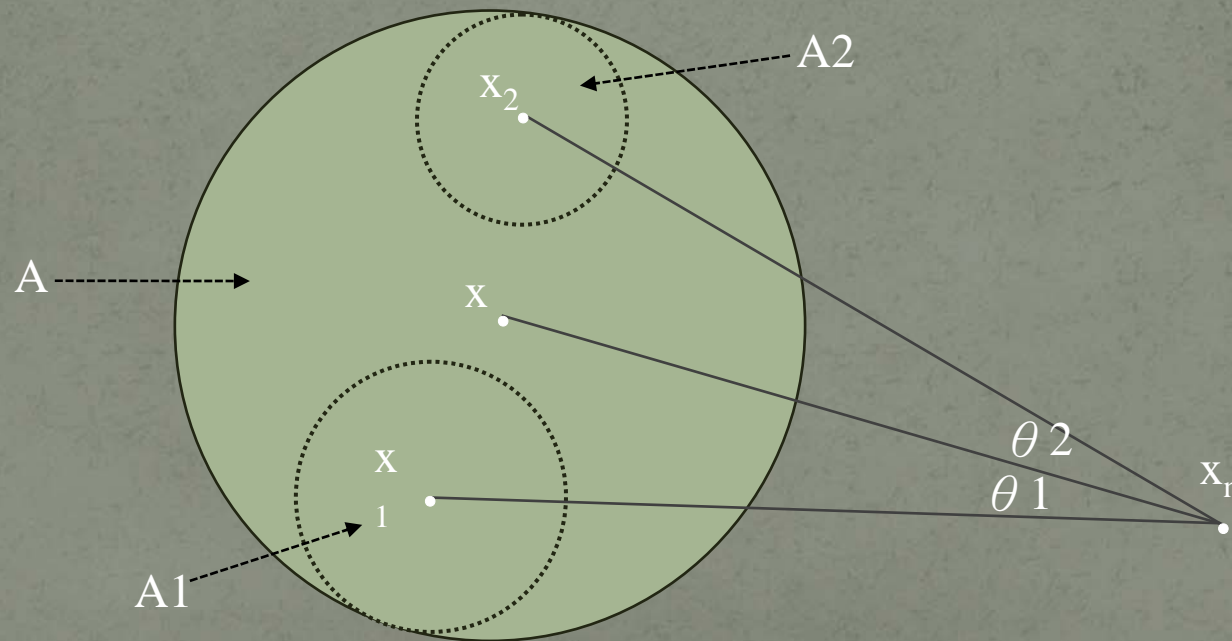
# Weak State Routing

- Every node has a table which maintains the states of the entire network.
- A weak state is composed of four parts :

Component Notation	Description
S	Weak <b>Bloom Filter</b> containing the <b>SetofIDs</b> portion
X	X coordinate of the center of the <b>GeoRegion</b>
Y	Y coordinate of the center of the GeoRegion
R	Magnitude of the radius of the GeoRegion

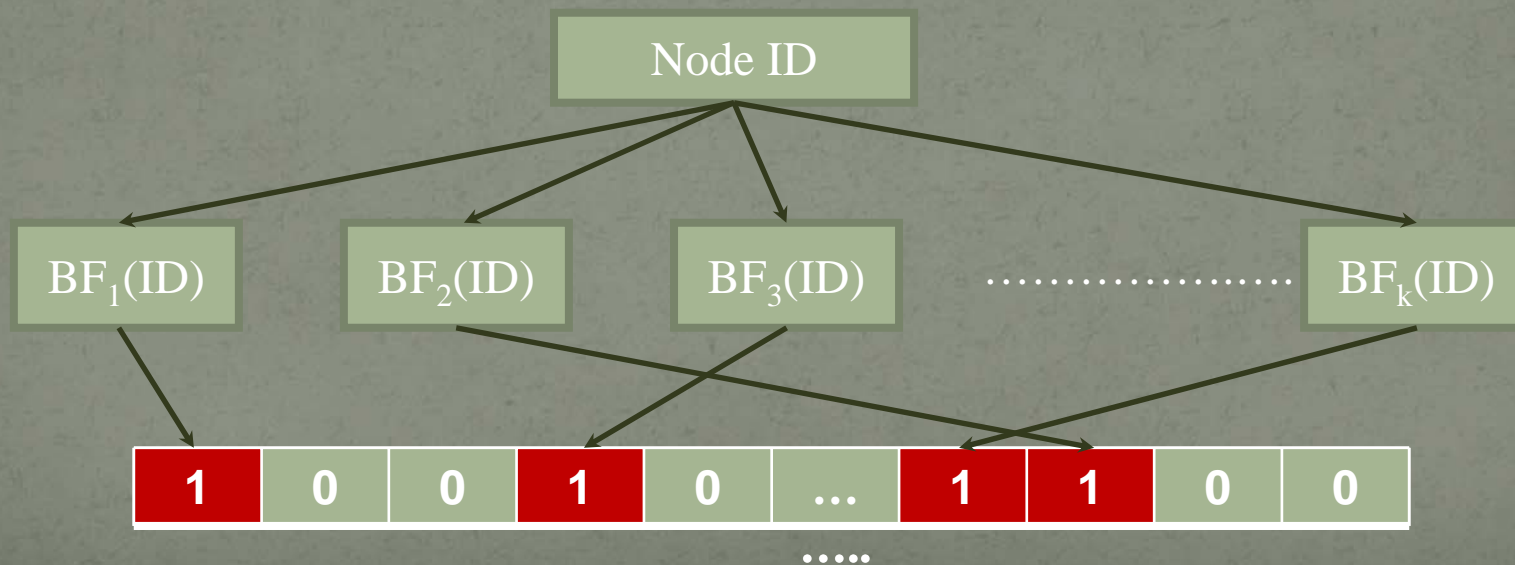
# Weak State Routing

- In order to decrease the beacon message overhead, it decreases the size of the message by using bloom filter and tries to combine some states to one state.



# Weak State Routing

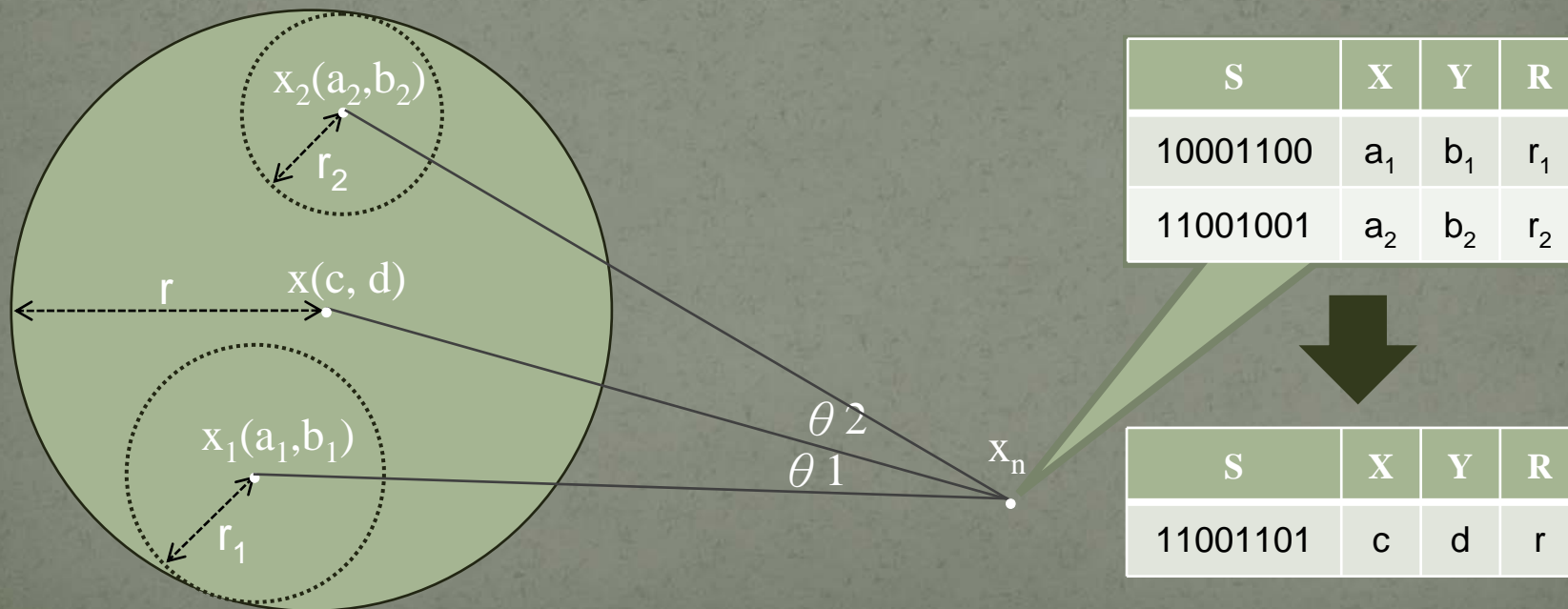
- Use bloom filter to calculate the virtual ID (SetofID) :
  - There are  $k$  hash functions  $BF_1 \sim BF_k$  in each node.
  - Every hash function produces a position in virtual ID ( $n$  bits,  $n > k$ ).
  - Setting the bits at these positions to 1.





# Weak State Routing

- If any two states can combine, then :
  - The new SetofID = SetofID<sub>1</sub> | SetofID<sub>2</sub>
  - The new GeoRegion is the smallest circle that contains both GeoRegions.

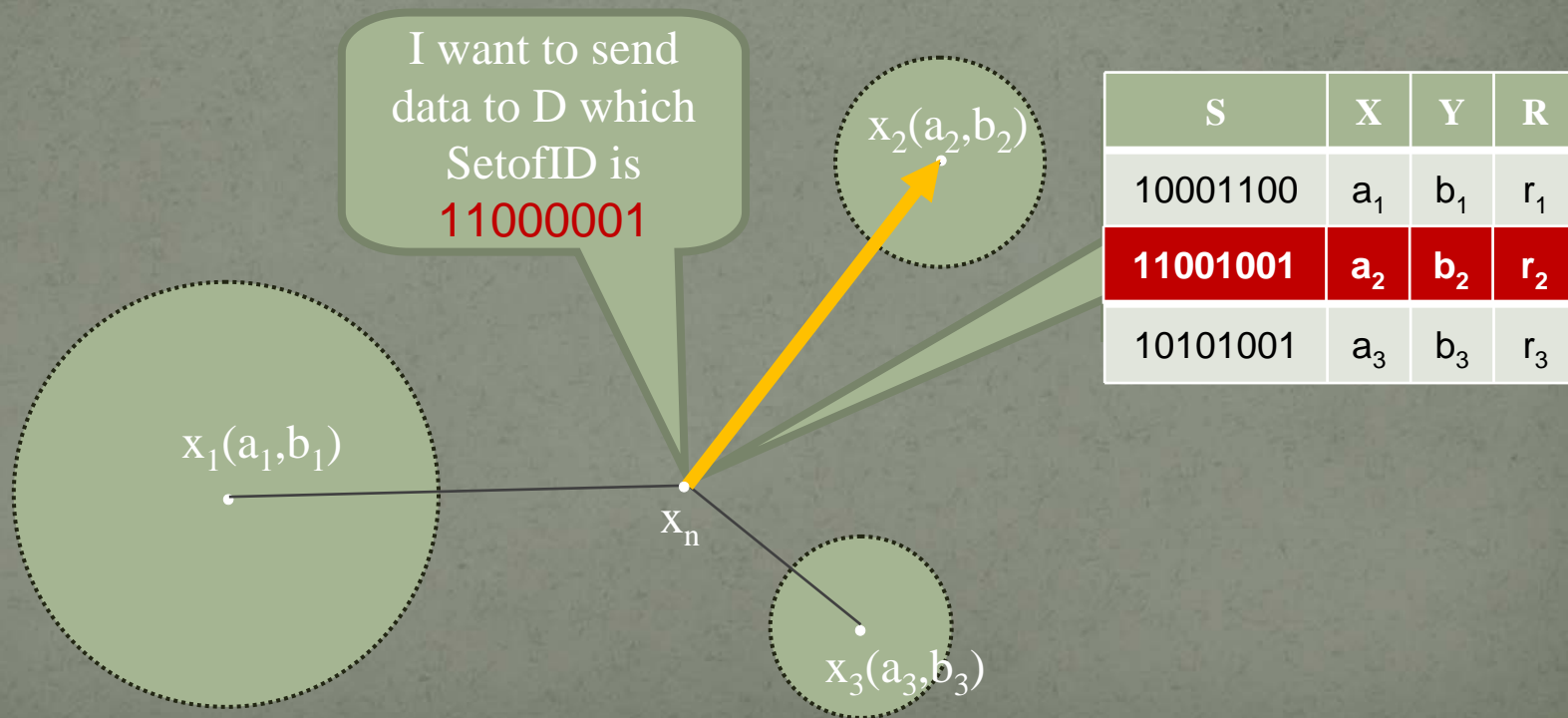


# Weak State Routing

- To combine two states to one state, the two states must satisfy some properties as follows :
  - Checking the GeoRegion radius after combining is smaller than or equal to threshold,  $R$ .
  - $\theta_1 + \theta_2$  is smaller than or equal to threshold,  $\theta$ .
  - Checking the total number of bits set to 1 in SetofID is smaller than or equal to threshold,  $B$ .

# Weak State Routing

- The source transmits data to the next node which is closest to highest confidence area (it's similar to geographic routing).



# Weak State Routing

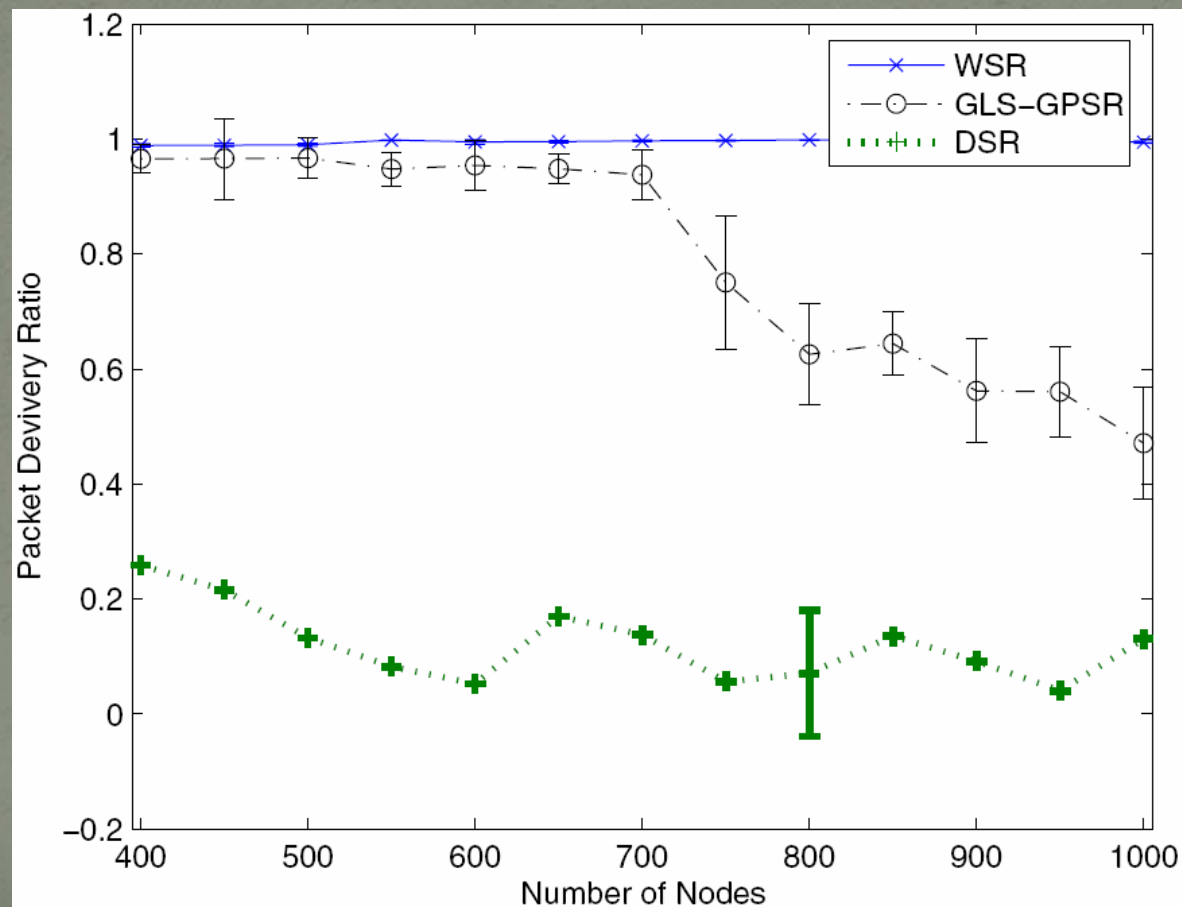
- At each time interval, the bits in each SetofID set to 1 are reset to 0 by a fixed probability  $p$  in order to adapt to the dynamic networks.
- Once the number of 1 is below a threshold value, WSR removes the state since the state is too old.
- The data packet is sent to a random direction if there are multiple states which have the same confidence.

# Simulation

- NS2
- 2500m x 2500m
- 500~1000 nodes
- Dynamic networks
- $v_{\min}=5\text{m/s}$ ,  $v_{\max}=10\text{m/s}$

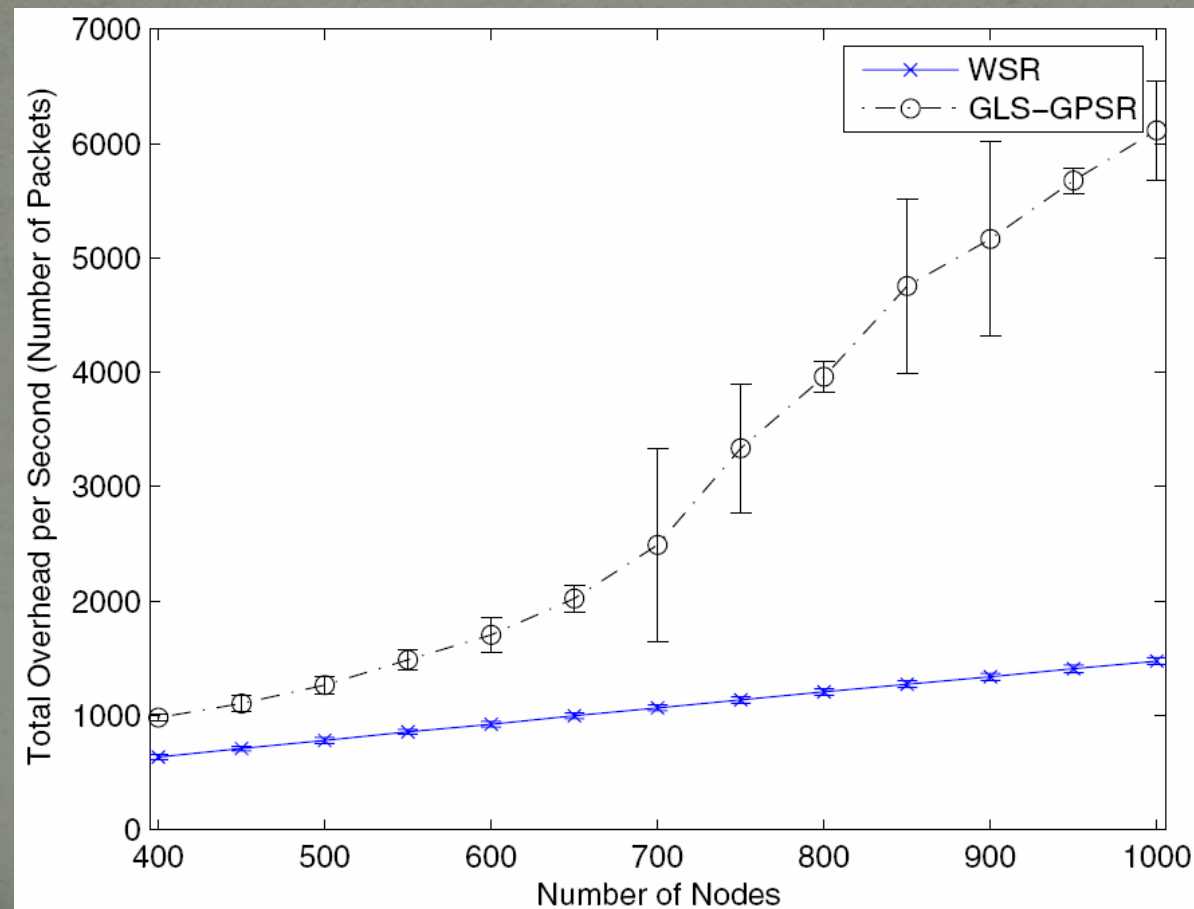
# Simulation

- Packet delivery success rate



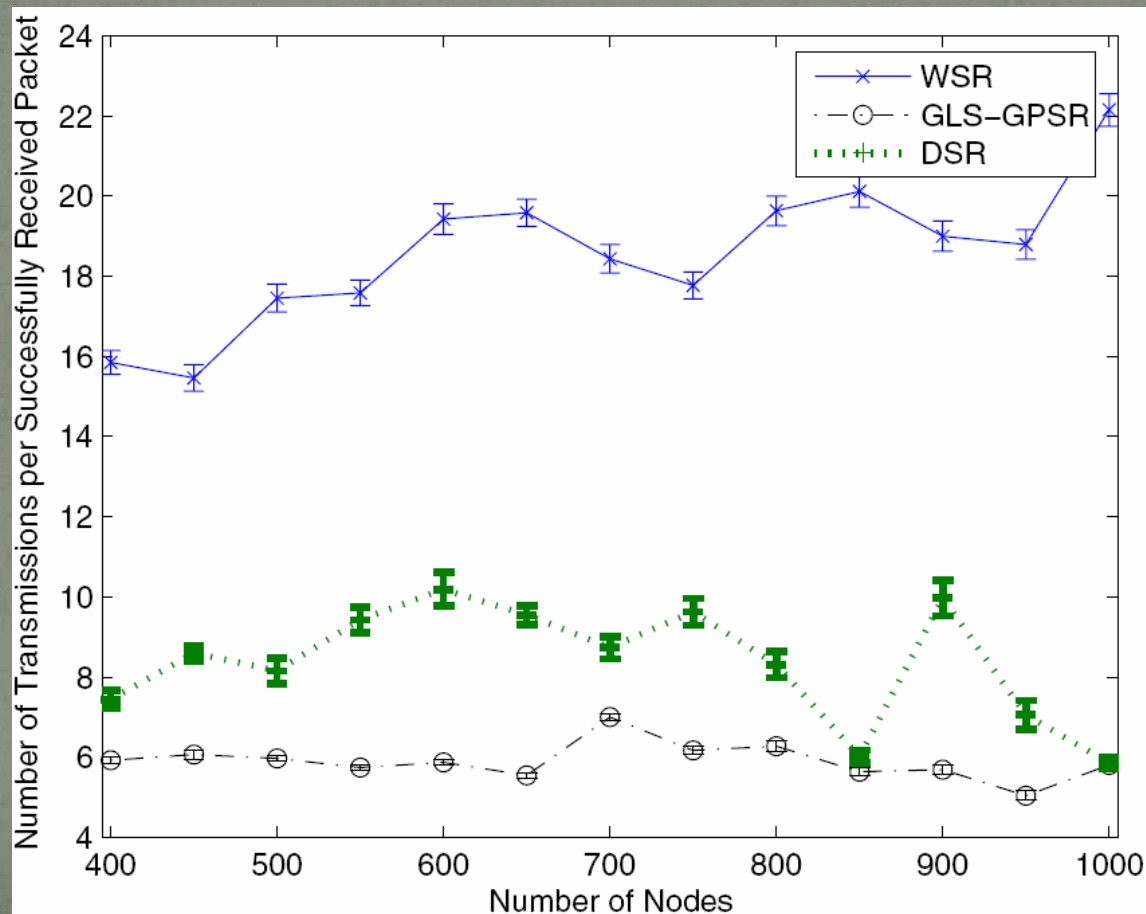
# Simulation

- Total overhead per second



# Simulation

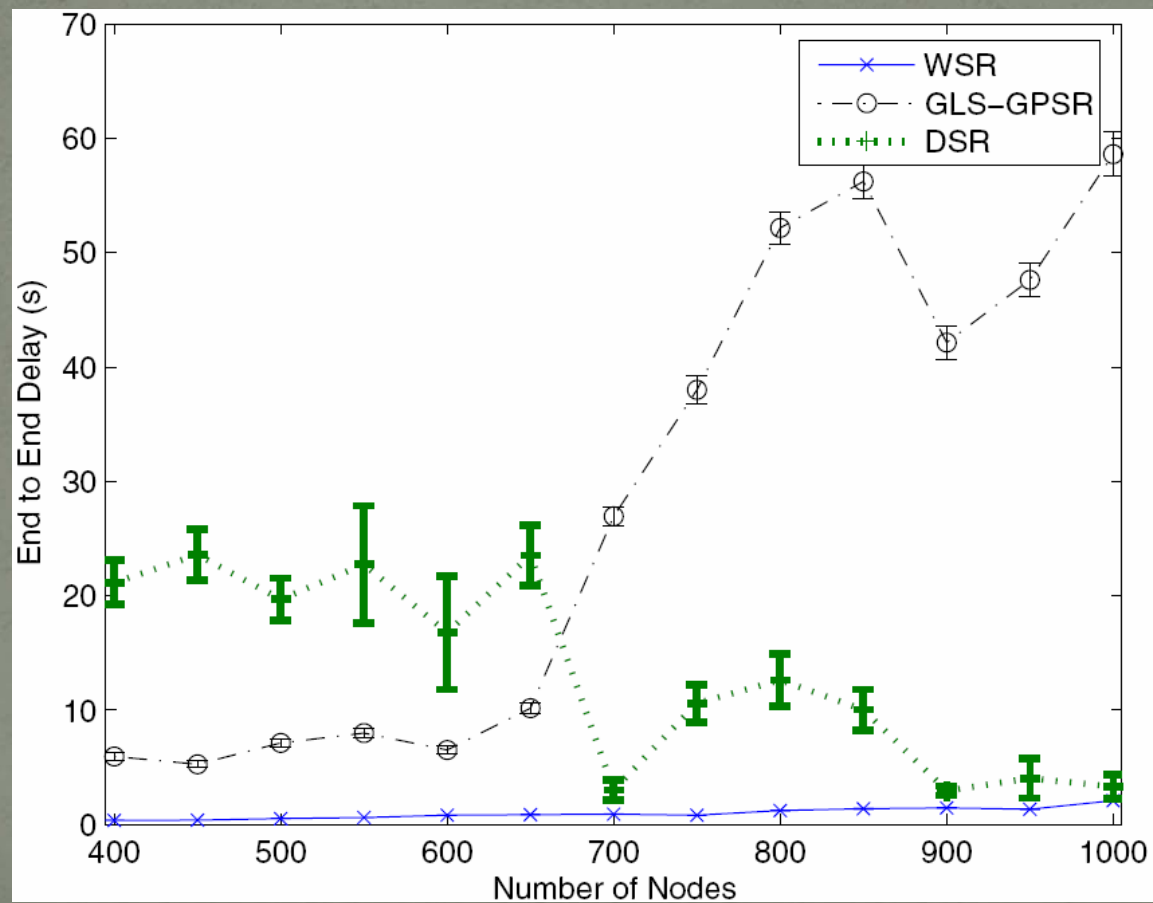
- Number of transmissions per successfully received packet



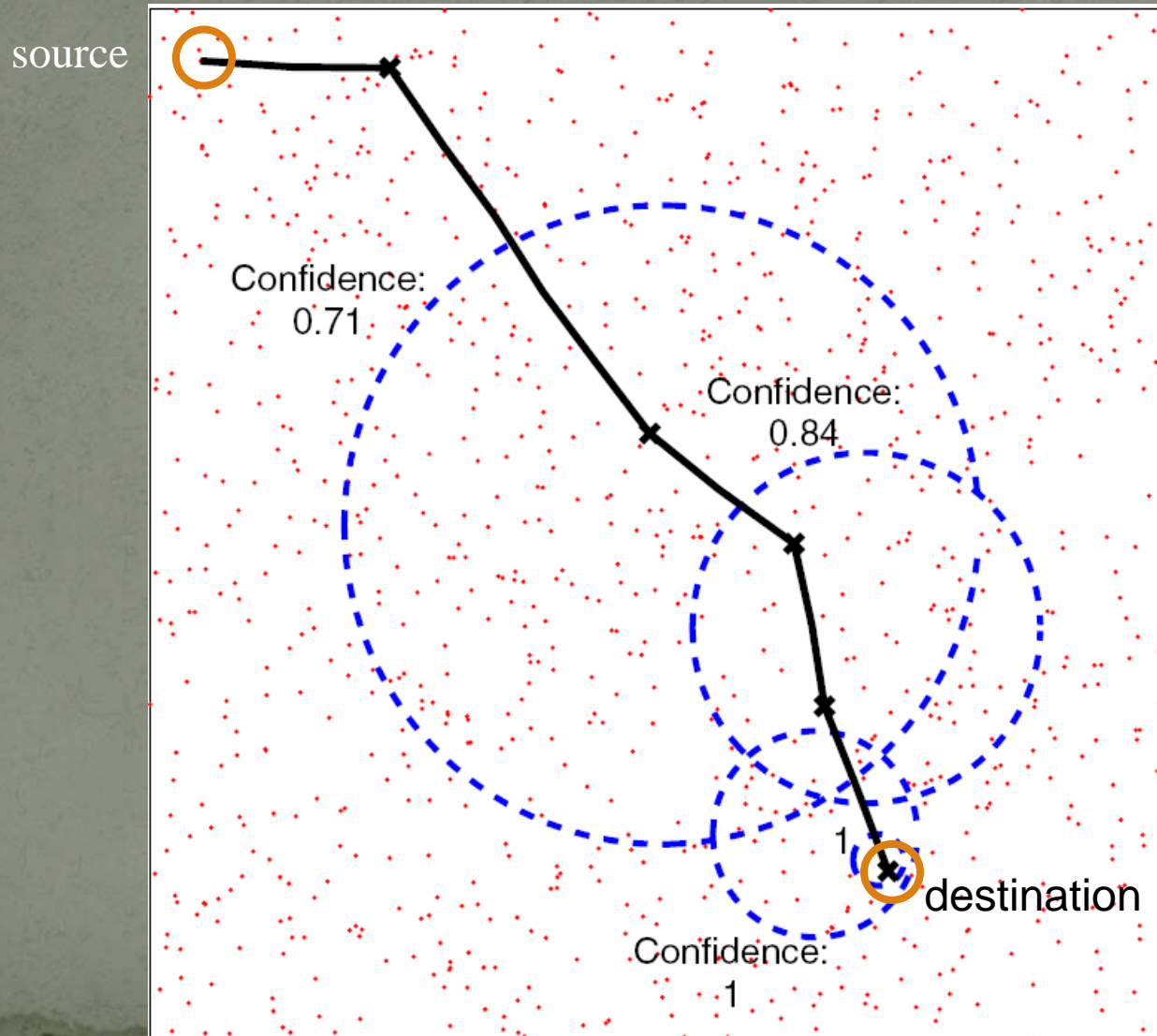


# Simulation

- End to end delay



# Simulation



# Conclusions

- WSR decreases the size of beacon messages and adapts to large scale dynamic networks.
- Using bloom filter, this mechanism reduces the size of the table which every node maintains.
- In simulation, WSR offers a high packet delivery ratio, more than 98%.
- It also retains high reachability, low overhead and delay.