

Efficient Retrieval of User Contents in MANETs

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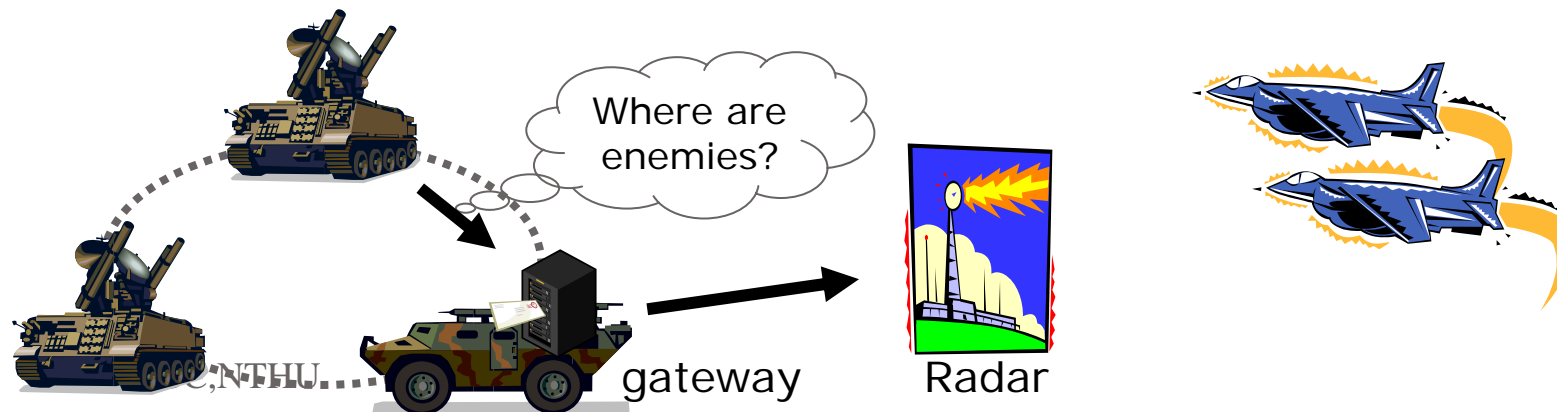
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outline

- 1.Introduction
- 2.System Assumptions
- 3.Eureka
- 4.Simulation
- 5.Conclusion

Introduction

- MANET environment
 - Wireless nodes are free to move, join, leave the network-factors
 - Highly-dynamic network system
- MANET applications often require to use resource/services at a gateway or at another user devices



Introduction

- MANET as a peer-to-peer network
 - User nodes not only require **content delivery** but also act as **content providers**
- A visible application of peer-to-peer MANET is in the field of vehicular networks
 - Car-to-Car Peer-to-Peer (C2P2)



Introduction

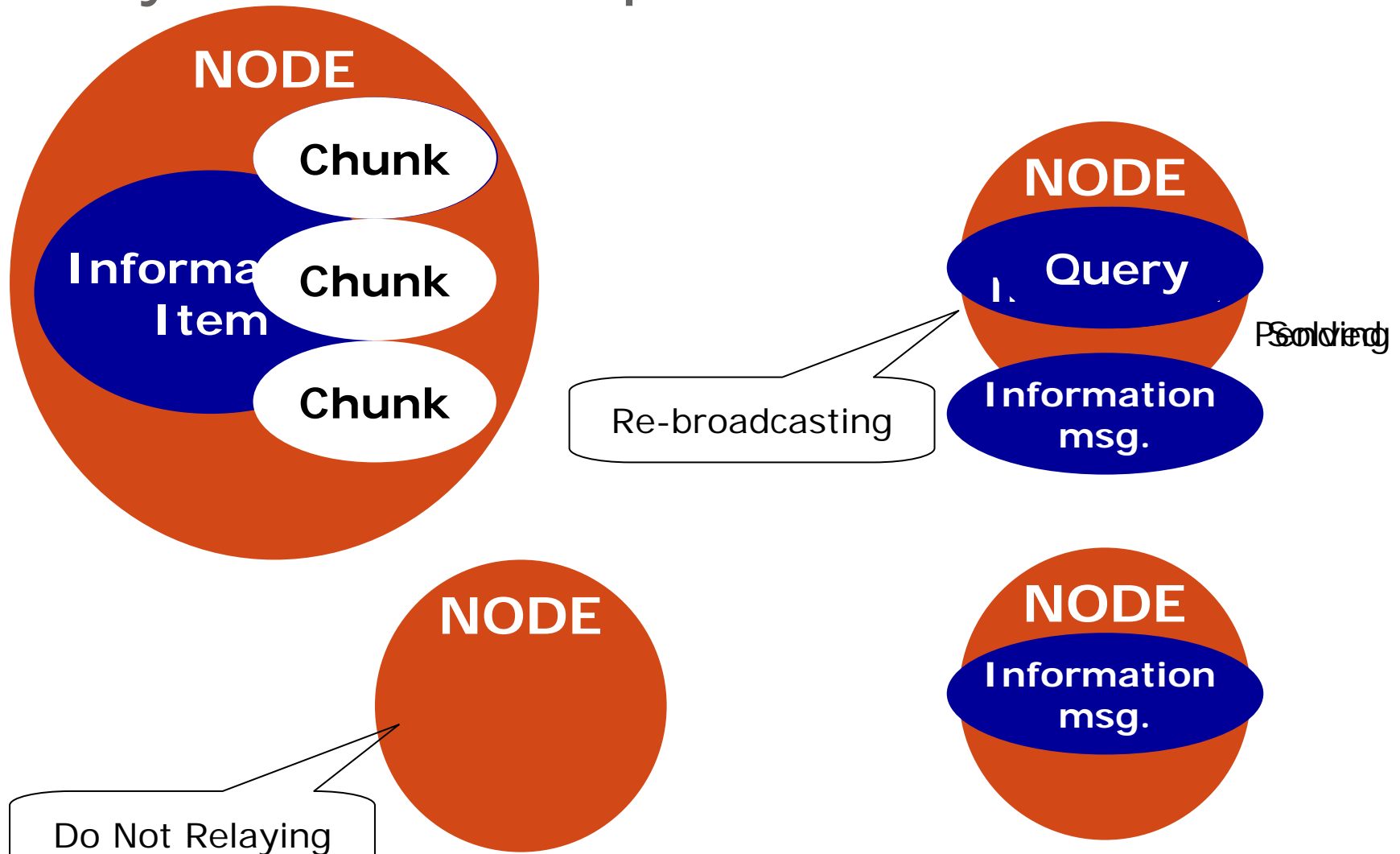
- Solution: Eureka
 - Purpose
 - Reduce query/messages overhead from broadcast storms in a cooperative MANET environment
 - Concept
 - Exploit the information density
 - Methodology to identify regions of the network where the required information is more likely to be stored and steer queries there

System Assumptions

- Basic Features

- Finite set of information items: each item has an identifier $0 \leq i \leq N$ and it is divided into chunks
- Nodes query other nodes for missing information item with rate λ . Each node generates queries for as many as C chunks of the missing items
- Queries are broadcast and relayed by other nodes
- Nodes with all, or part, of the requested information chunks reply to the query. Chunks are returned through an application-driven unicast path
- Information chunks retrieved by the requesting node are locally cached and then dropped at a rate of μ chunks per second

System Assumptions



System Assumptions

- Message exchange
 - Queries
 - Broadcast by requesting and relaying nodes
 - Before relaying a query, a node stores its routing details (ID, src add, last add) and set status for the missing chunk to “pending”
 - Responses
 - MAC-layer unicast transmission by relaying nodes.
 - Next-hop chosen by routing tracks stored at relay node applications
 - Promiscuous mode are allows overhearing nodes to toggle the status of pending queries (pending-> solved)
 - Avoiding the relay of duplicated information msg.

Eureka

- Eureka and Information density
 - Provide each node in the network with an estimate of information density in its proximity
 - Difference between local and other nodes' estimates are used to decide whether a query must be forwarded
 - Thus, queries travel toward areas where it is more likely that information is found

Eureka

- Information density estimation
 - Fully distributed process, Run by all nodes
 - Information Density Function $\delta_i(x,y)$
 - Spatial density of information item i , cached at nodes around point(x,y)
 - We measure the information density in *copies/m²*(in case of uni-dimensional topologies such as a highway scenario, we consider $\mathcal{O}(x)$ measured in *copies/m*).

^

- Information density estimation
 - Two contributions
 - $S_{i,j}^l$: locally-computed density sample from generated, overheard and received information messages of item i at step j
 - $S_{i,j}^d$: density sample computed from advertised samples by neighboring nodes
 - A Moving Average(MA) filter weights the contributions and returns the density estimate $\hat{\delta}_i(\mathbf{x}, \mathbf{y})$
 - It is not important to us that the estimates match the absolute values of the actual density

Eureka

- Three techniques to counter the drawbacks of flooding
 1. Queries are issued with a Time-To-Live(TTL)
 2. Relay nodes delay query forwarding by a Query Lag Time, in the hope that a response is returned by neighboring nodes

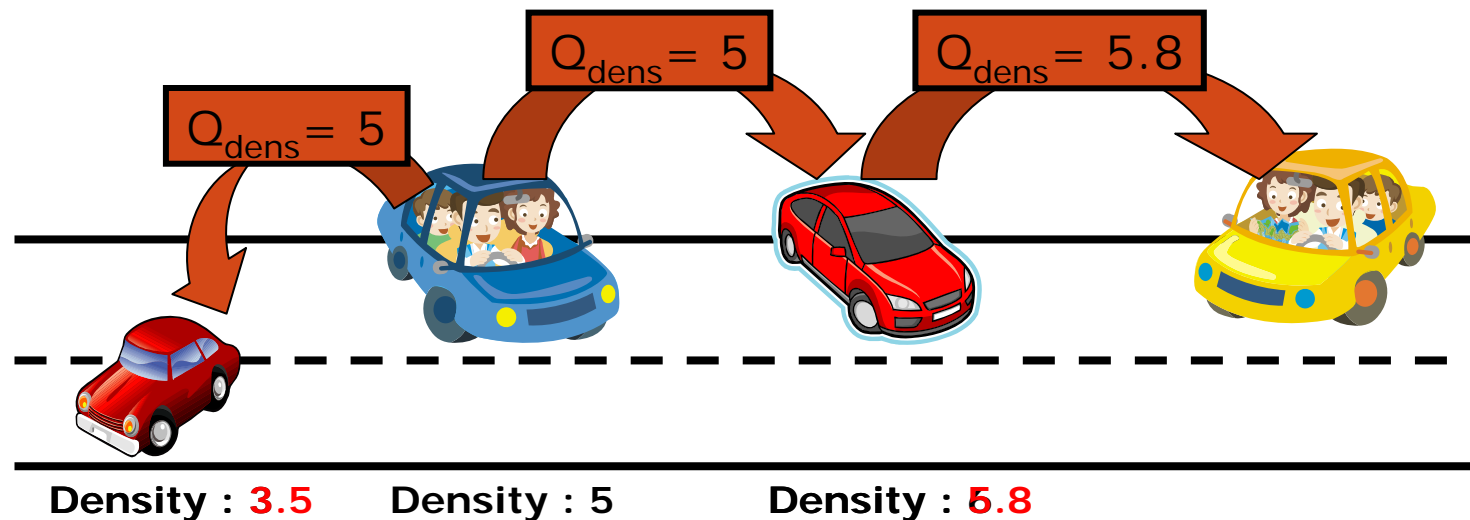
Mitigated Flooding

Eureka

- Three techniques to counter the drawbacks of flooding
 1. Queries are issued with a Time-To-Live(TTL)
 2. Relay nodes delay query forwarding by a Query Lag Time, in the hope that a response is returned by neighboring nodes
 3. Queries include information density as seen by issuing node; relay nodes forward query only if they “see” higher information density

Eureka

- Three techniques To counter the drawbacks of flooding
 3. Queries include information density as seen by issuing node; relay nodes forward query only if they “see” higher information density



It updates its density estimate
But, does **not** forward the query

HSCC,NTHU

It updates its density estimate
and forward the query

Information Density Estimations

- *A. Local information density sample*
 - $S_{i,j}^1$: locally-computed density sample from generated, overheard and received information messages of item i at step j

$$1 - \frac{h_Q - 1}{TTL}$$

- h_Q is equal to the number of hops covered by the query
- Range between $\frac{1}{TTL}$ to 1

Information Density Estimations

- B. *Distributed information density sample*

$$s_{i,j}^d(n) = \frac{\sum_{m \in \mathcal{M}_{i,j}(n)} s_{i,j}^l(m)}{|\mathcal{M}_{i,j}(n)|}$$

- $\mathcal{M}_{ij}(n)$ is the set of neighbor nodes which advertised their local sample to node n , for information item i and sampling step j

Information Density Estimations

- *C. Overall information density sample*

$$s_{i,j}(n) = \frac{s_{i,j}^l(n) + s_{i,j}^d(n)}{2}$$

- *D. Finer information density estimate—using MA filter*

$$\begin{aligned} \hat{\delta}_{i,j}(n) = & \sum_{k=1}^{W-1} (1 + \alpha^W - \alpha^{W-k}) s_{i,j-k}(n) \\ & + (1 + \alpha^W - \alpha) \sum_{k=W}^j \alpha^{k-W+1} s_{i,j-k}(n) \quad (1) \end{aligned}$$

$$\hat{\delta}_{i,j}(n) = \sum_{k=1}^{W-1} (1 + \alpha^W - \alpha^{W-k}) s_{i,j-k}(n) + (1 + \alpha^W - \alpha) \sum_{k=W}^j \alpha^{k-W+1} s_{i,j-k}(n) \quad (1)$$

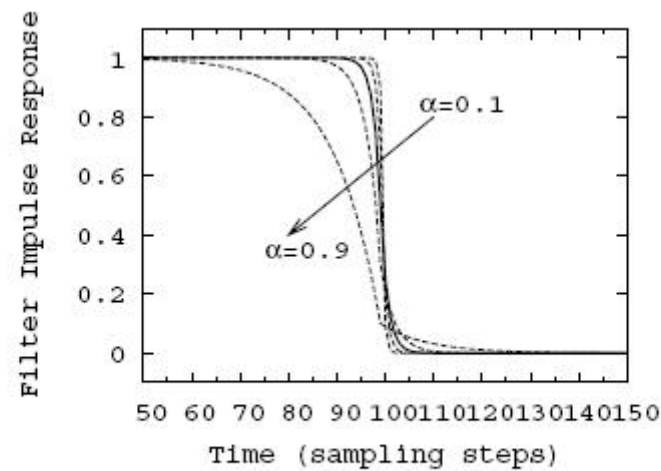


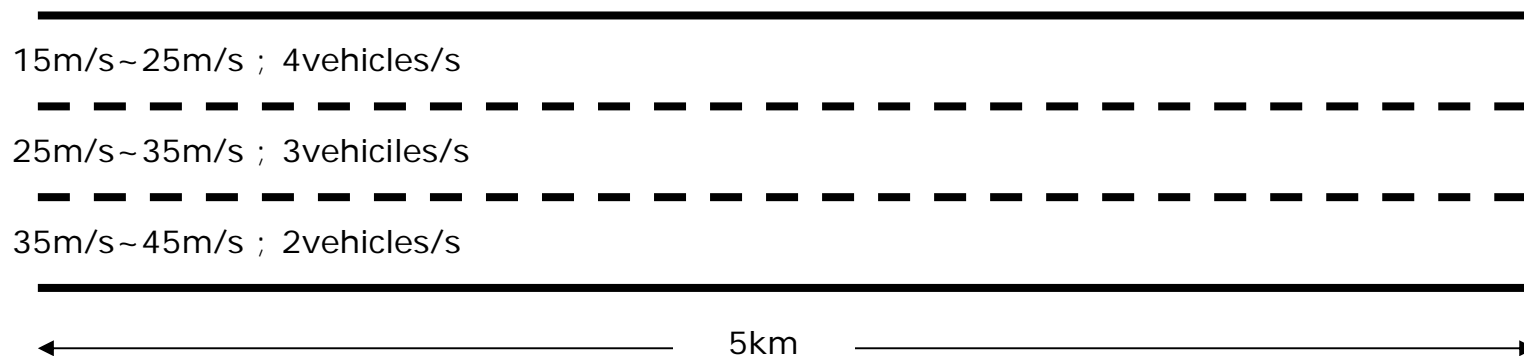
Fig. 2. MA filter impulse response, with $W = 100$ and varying α

Simulation

- Performance evaluation
 - Simulation with ns-2
 - Scenario parameter
 - Information set cardinality, $N=[1,25]$
 - Each information item divided into up to 30 chunks
 - Queries generation rate, $\lambda=[1,6]$ queries/ms
 - Cached information drop rate, $\mu=[5,50]$ drops/ms
 - Queries Time-to-Live, TTL=10
 - Topologies
 - Highway
 - Urban
 - Nodes have a 802.11 MAC layer, radio range : 100m

Simulation

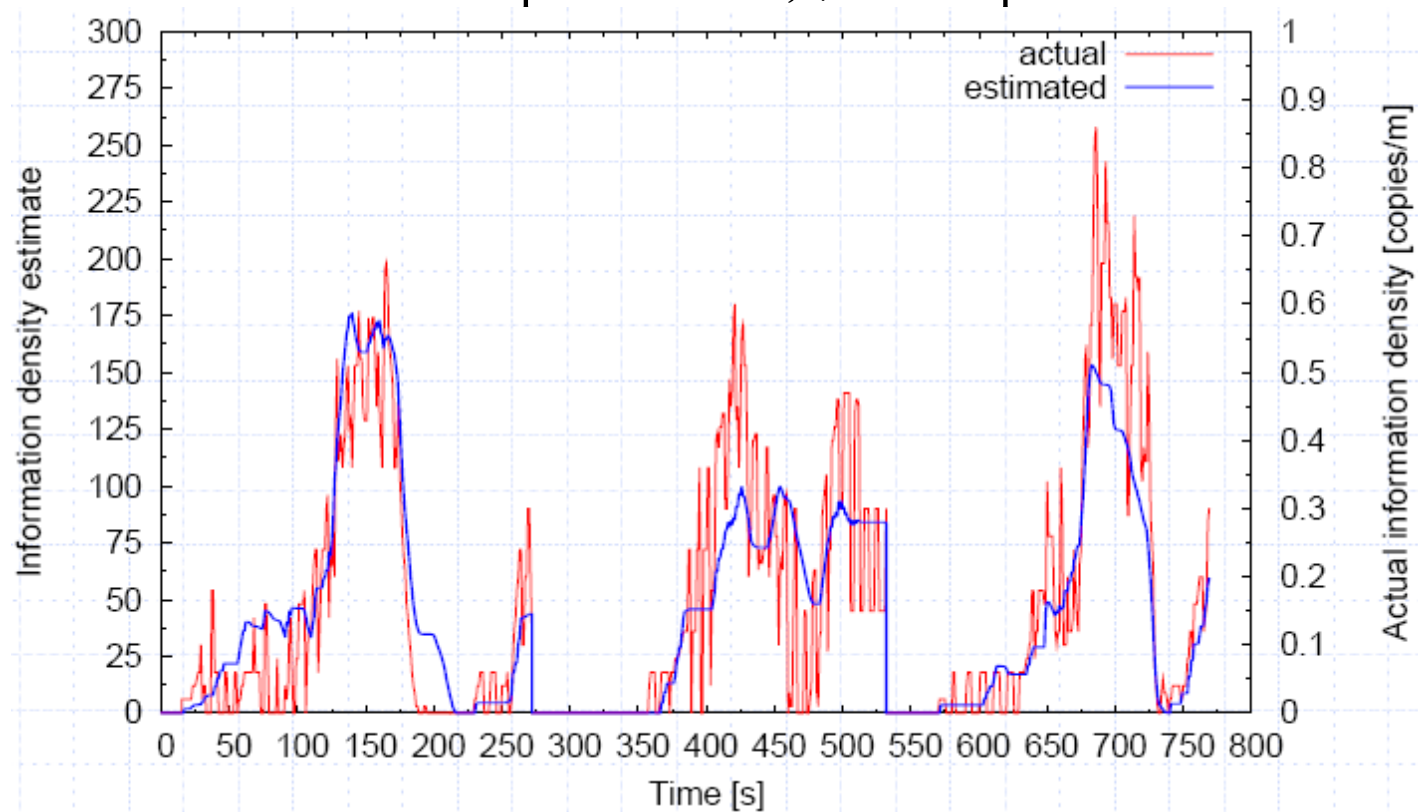
- Road Scenario : Highway
 - 5km straight, unidirectional road with three parallel lanes at different speeds
 - Lane changes allowed when overtaking
 - Speeds ranging from 15m/s to 45m/s
 - An “information gateway” node is located halfway along the road



Simulation

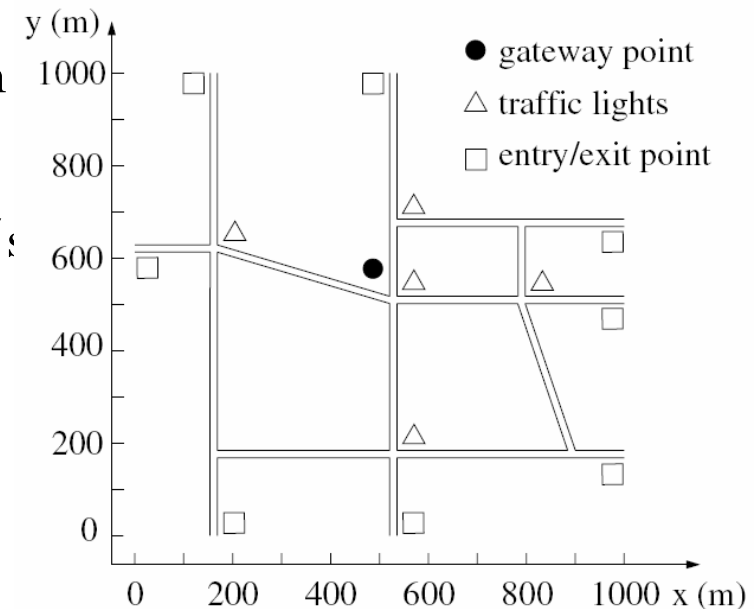
- Estimation Accuracy – Highway

$$\lambda = 3 \text{ queries/ms} ; \mu = 25 \text{ queries/ms}$$



Simulation

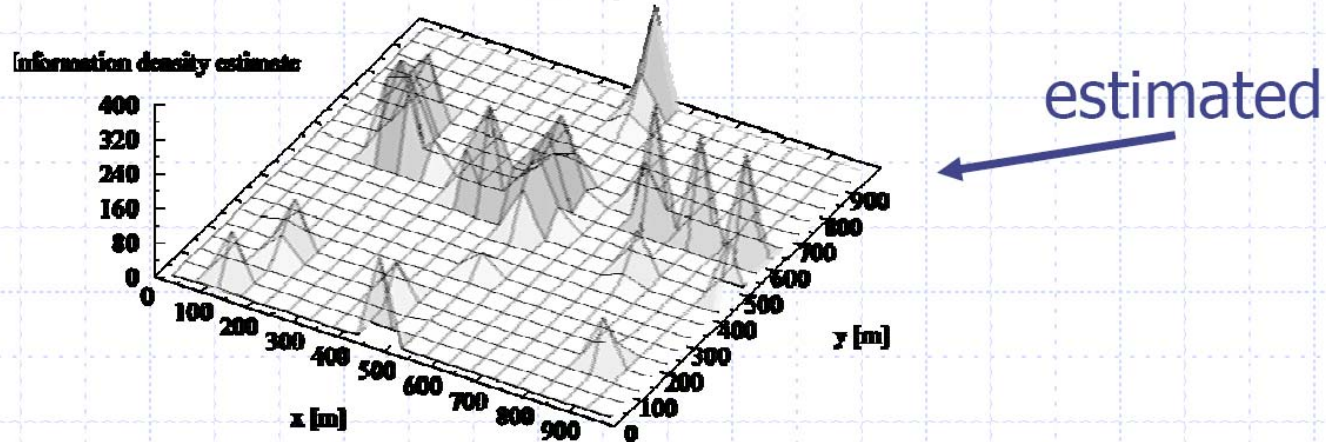
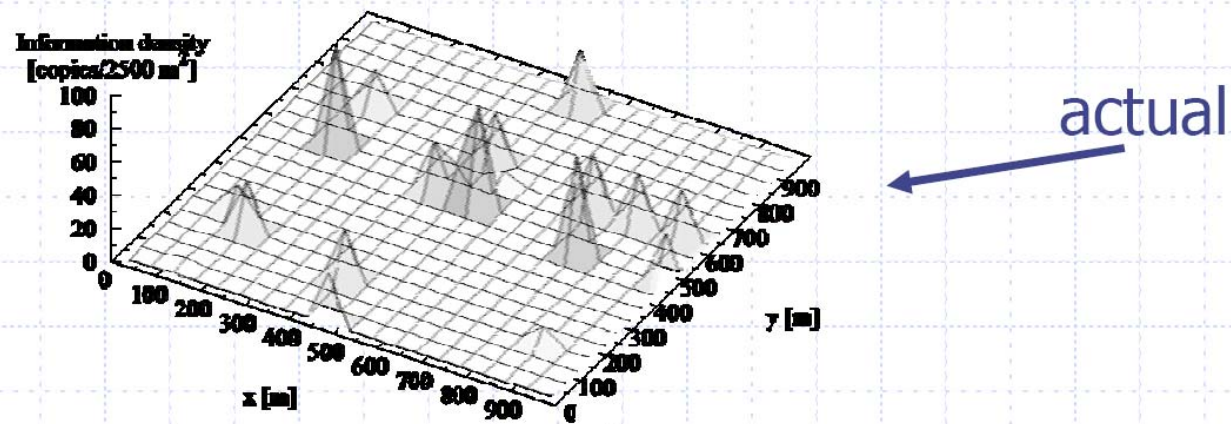
- Road Scenario : Urban
- **CanuMobiSim**
- City sections with traffic lights and stop signs at intersection.
- Vehicles enter/exit at random from entry/exit points
- Speeds ranging from 10m/s to 20m/s
- On average, 70 vehicles at 5.8m/s
- One “information gateway”

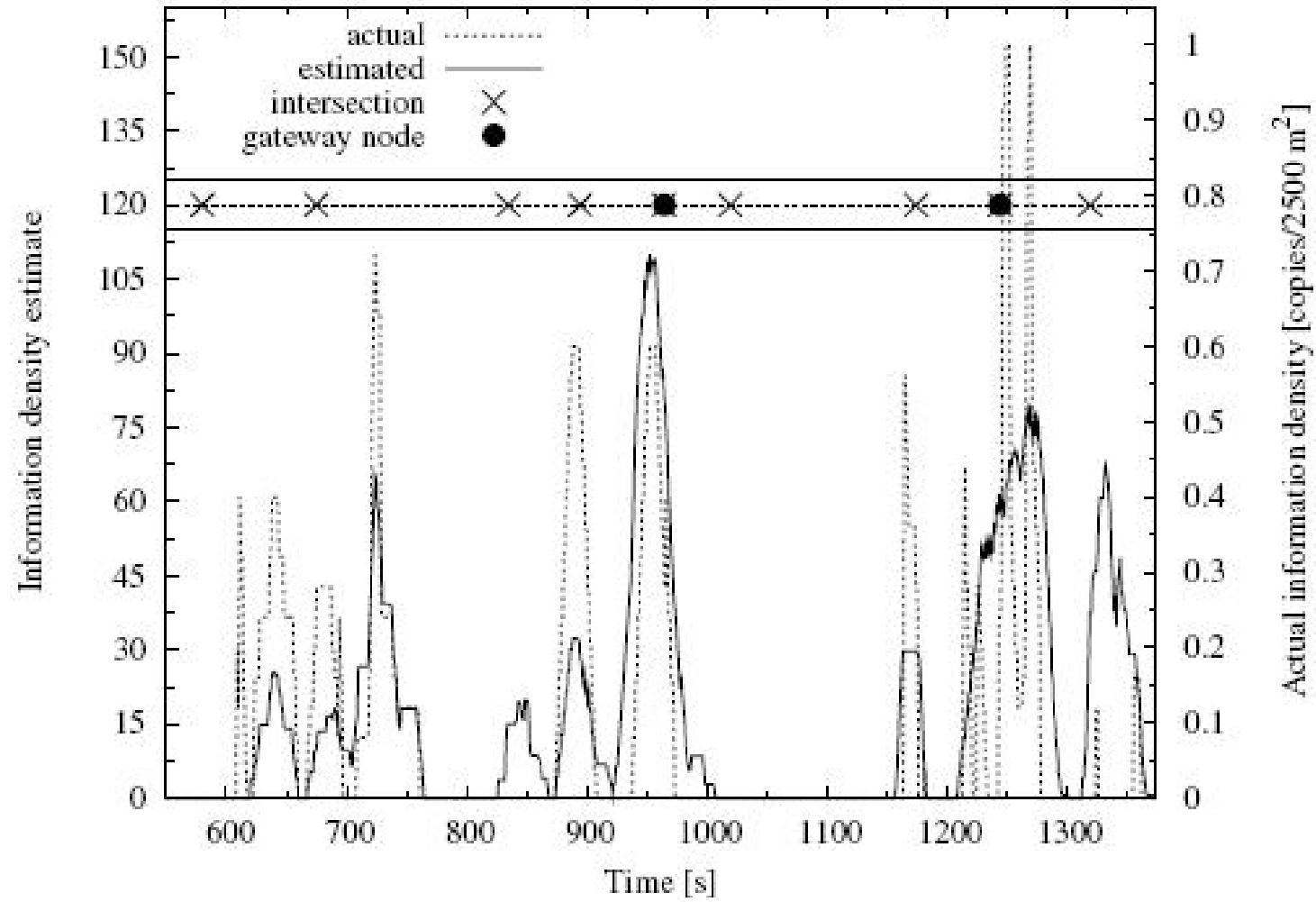


Simulation

- Estimation Accuracy – Urban

$$\lambda = 3 \text{ queries/ms} \cdot 11 = 25 \text{ queries/ms}$$





Conclusion

- Eureka provides effective overhead reduction
- User use it to direct queries toward areas where the requested information is denser
- Broadcast storms are prevented and congestion is reduced

Thank You !