

Analysis of Multi-Hop Emergency Message Propagation in Vehicular Ad Hoc Networks

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OUTLINE

- INTRODUCTION
- NETWORK AND CHANNEL MODEL
- THEORETICAL ANALYSIS
- DISCUSSION AND SIMULATION
- CONCLUSIONS



INTRODUCTION(1/2)

- Safety applications are attracting a lot of attention because of improving driver's awareness of surrounding environment.
- We can improve reliability of 1-hop emergency message by devoting more resources to safety-related message dissemination, like setting the transmit power level [13, 14].

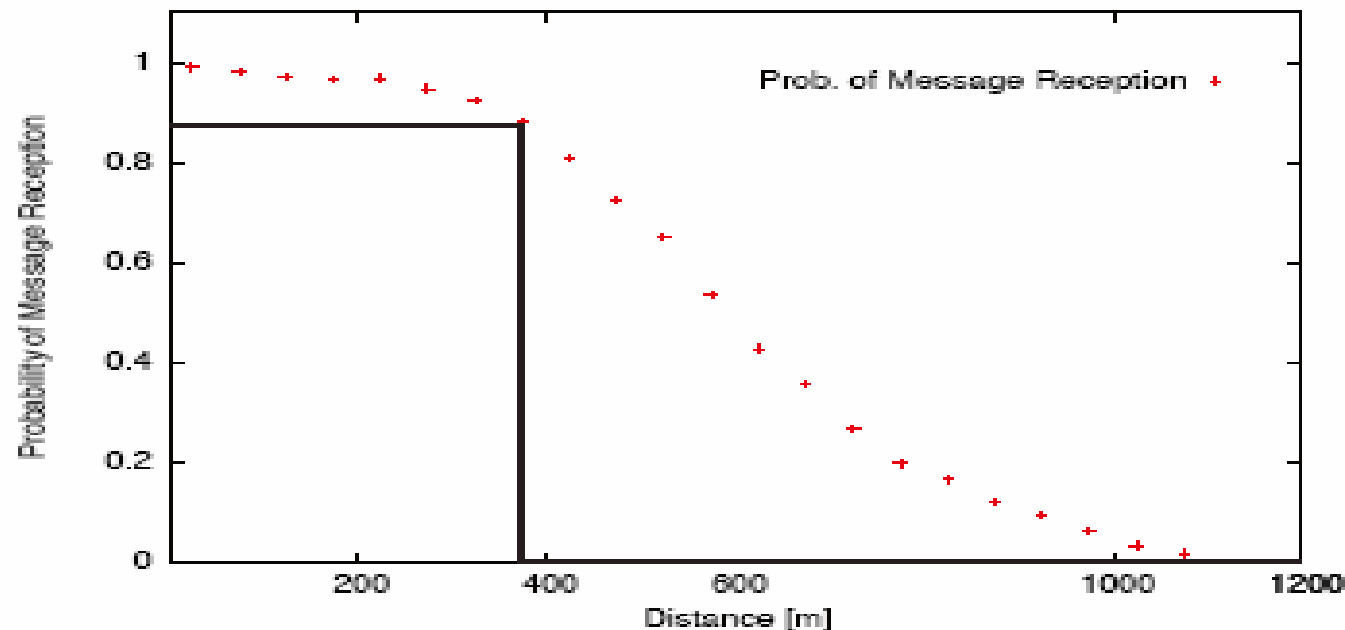


INTRODUCTION(2/2)

- Goals of this paper
 - Analyze multi-hop emergency message protocols.
 - Derive lower bounds on $P(d, t, p)$.
 - Investigate the tradeoff between safety-level and emergency message resource wastage.
 - **The relative advantage of having an increased p tends to decrease as d increases.**
 - **The efficacy of increased p displays a clear dependence on the traffic conditions.**

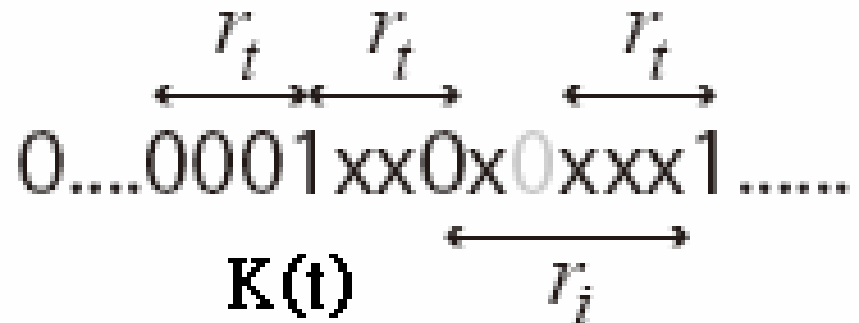
NETWORK AND CHANNEL MODEL

- All cars are equally spaced along a line : $S_n(t)$.
- Interference between concurrent transmission.
- Channel model :



THEORETICAL ANALYSIS(1/8)

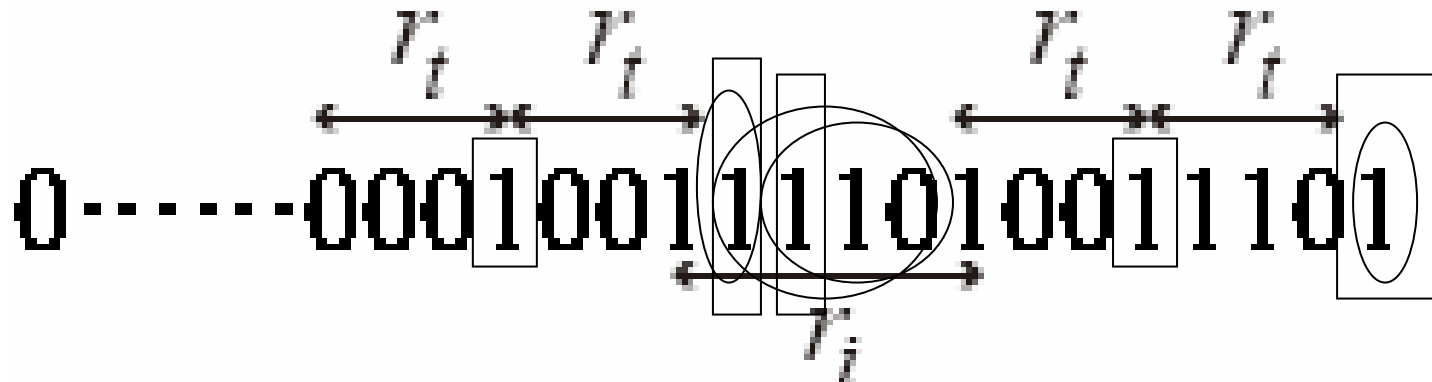
- IDEALIZED strategy :
 - Reliability :
 - Turns every internal 0-node into a 1-node with probability p at each round.
 - Fast :
 - Every node i such that $i > k$ and $i - k \leq r_t$ assigns probability p of turning into a 1.



THEORETICAL ANALYSIS(2/8)

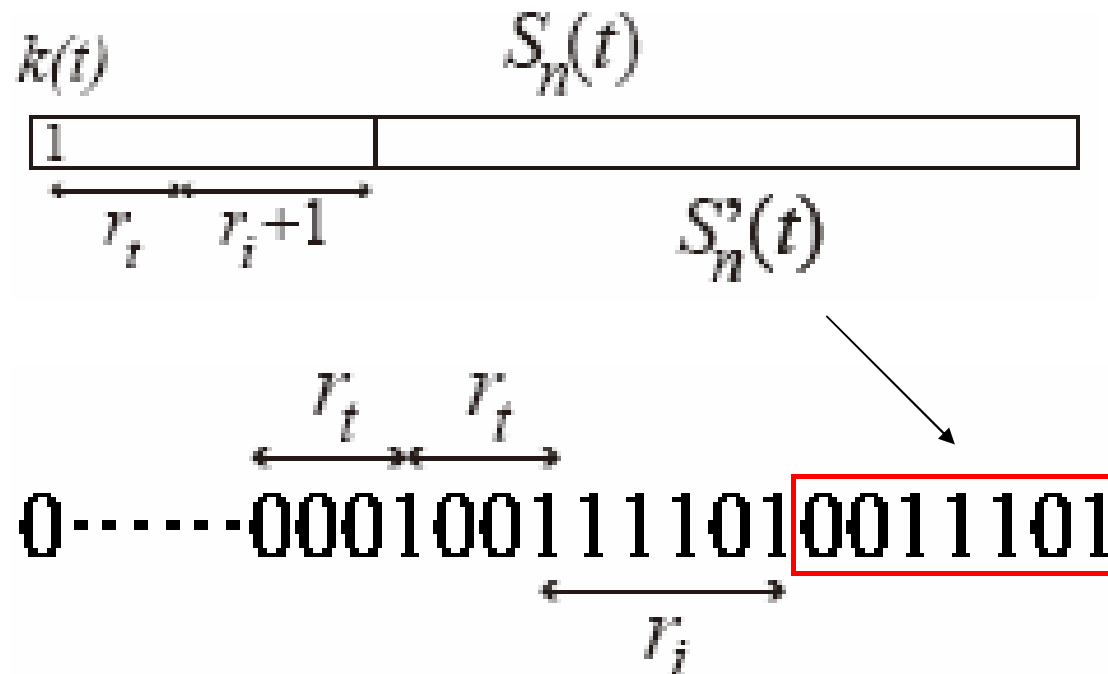
■ Global strategy :

- 1 stage has h rounds, $h = 2$ if $r_t = r_i$, and $h = 3$ if $r_t < r_i \leq 2r_t$.
- Input : $S_n(t)$, output : set T_h of transmitters for round $t + h$.



THEORETICAL ANALYSIS(3/8)

- IMGlobal strategy :



THEORETICAL ANALYSIS(4/8)

■ Approximation bounds : (α, β)

□ α : Gives chance at least p/α of turning into 1 to all internal 0-nodes at each round.

□ β : $E_s(k(t)) \geq \frac{E(k(t))}{\beta}$

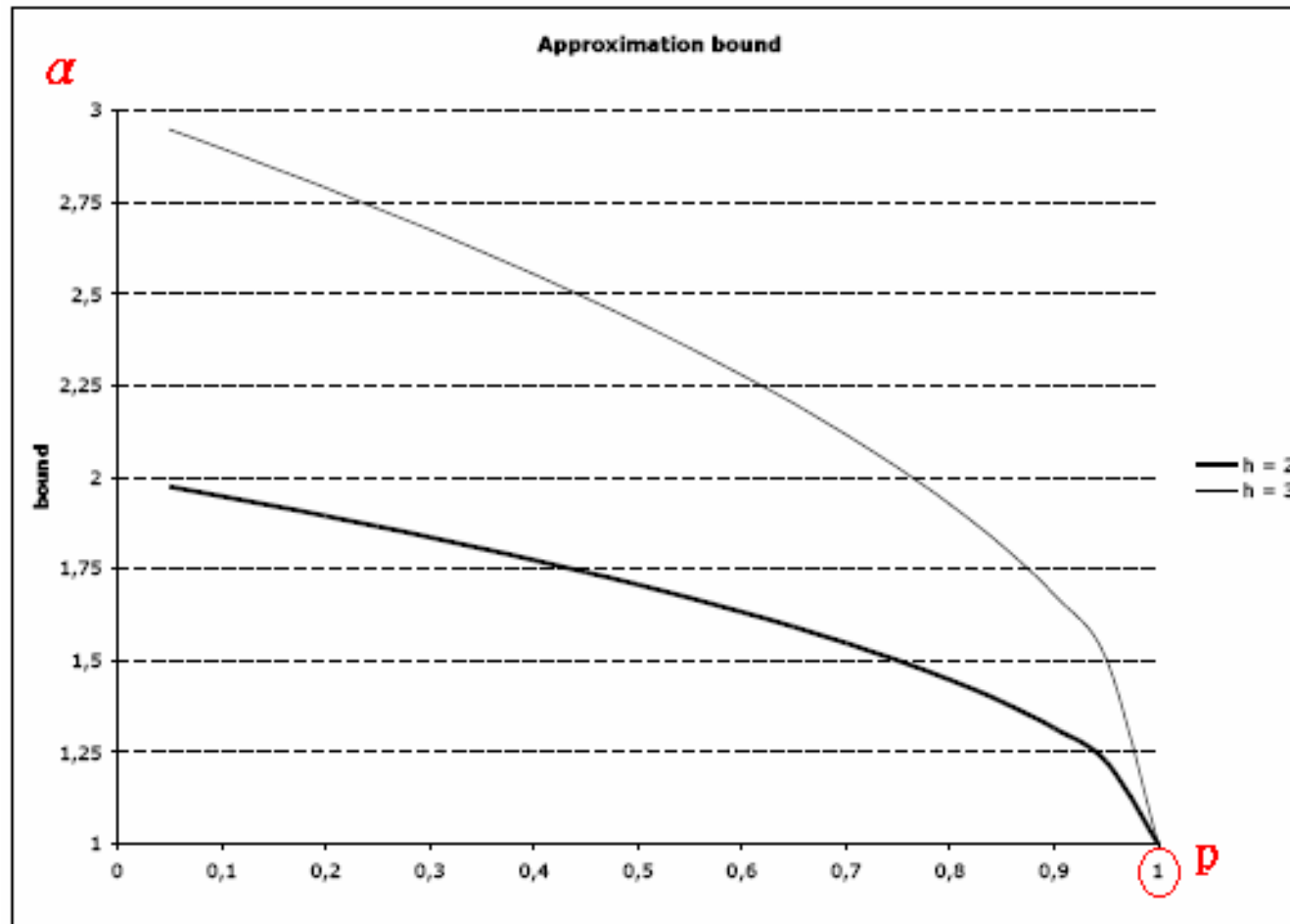
□ GLOBAL : $\left(\frac{p}{1 - \sqrt[h]{1-p}}, h \right)$

$$(1-q)^h \geq (1-p)$$

$$q \leq 1 - (1-p)^{\frac{1}{h}}$$

□ IMGLOBAL : $\left(\frac{p}{1 - \sqrt[h]{1-p}}, 1 \right)$

THEORETICAL ANALYSIS(5/8)



THEORETICAL ANALYSIS(6/8)

- Time-constrained reception probability : $P(d,t,p)$

$$P(\bar{d}, \bar{t}, p) = \sum_{h=1}^{\bar{t} - \frac{\bar{d}}{r_T}} \boxed{Prob(Succ(p, h))} \cdot Prob(EStart(\bar{d}, \bar{t} - h))$$

- Lower bound of $Prob(Start(d,t))$:

$$\boxed{1 - (1 - \bar{p})^h}$$

- Define $X(t) = k(t+1) - k(t)$.

$$E(p, r_T) = \sum_{h=1}^{r_T} h \cdot p \cdot (1 - p)^{r_T - h} = r_T + 1 - \frac{1 - (1 - p)^{r_T + 1}}{p}$$

$$\begin{aligned} V(p, r_T) &= \sum_{h=1}^{r_T} h^2 \cdot p \cdot (1 - p)^{r_T - h} - E(p, r_T)^2 \\ &= (1 - p) \left[1 - (2r_T + 1)p(1 - p)^{r_T} - (1 - p)^{2r_T + 1} \right] / p^2 \end{aligned}$$

$$\mathbf{X} = \sum_{i=1}^k X_i$$

$$Prob(\mathbf{X} \leq E(\mathbf{X}) - \lambda) \leq e^{-\frac{\lambda^2}{2 \cdot \sum_{i=1}^k E(X_i^2)}}, \text{ for any } \lambda > 0$$

$$Prob(\mathbf{k}(\mathbf{t}) \leq t \cdot E(p, r_T) - \lambda) \leq e^{-\frac{\lambda^2}{2 \cdot t \cdot E(\mathbf{X}(\mathbf{t})^2)}}$$

$$Prob(Start(\bar{d}, \bar{t})) = Prob(\mathbf{k}(\bar{\mathbf{t}}) > \bar{d} - 1)$$

$$Prob(\mathbf{k}(\bar{\mathbf{t}}) > \bar{t} \cdot E(p, r_T) - (\bar{t} \cdot E(p, r_T) - \bar{d} + 1))$$

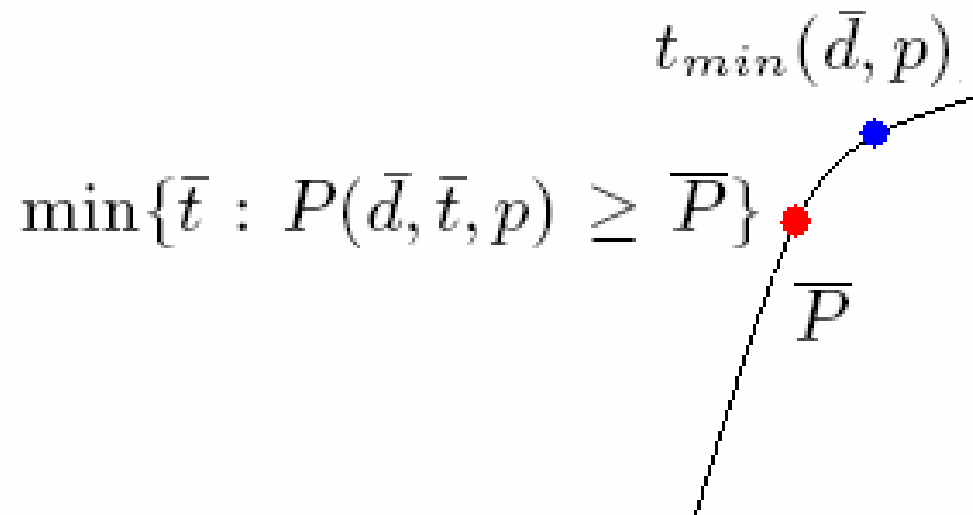
$$\bar{t} > \frac{\bar{d}-1}{E(p, r_T)}. \text{ Then,}$$

$$Prob(Start(\bar{d}, \bar{t})) \geq 1 - e^{-\frac{(\bar{t}E(p, r_T) - \bar{d} + 1)^2}{2\bar{t}E(\mathbf{X}(\bar{\mathbf{t}})^2)}}$$

$$\lambda > 0$$

THEORETICAL ANALYSIS(7/8)

- Upper bound of t : $\min\{\bar{t} : P(\bar{d}, \bar{t}, p) \geq \bar{P}\} \leq t_{min}(\bar{d}, p)$



THEORETICAL ANALYSIS(8/8)

- Define $\bar{h} = \left\lceil \frac{\ln(1-\sqrt{\bar{P}})}{\ln(1-p)} \right\rceil \longrightarrow \text{Prob}(\text{Succ}(p, h)) \geq \sqrt{\bar{P}}$ for each $h \geq \bar{h}$.

$$P\left(1 - e^{-\frac{((\bar{t}-\bar{h}) \cdot E(p, r_T) - \bar{d} + 1)^2}{2(\bar{t}-\bar{h}) \cdot E(\mathbf{X}(t)^2)}} \geq \sqrt{\bar{P}} \mid \text{EStart}(\bar{d}, \bar{t} - h)\right)$$

$$\bar{t} - \frac{\bar{d}}{r_T}$$

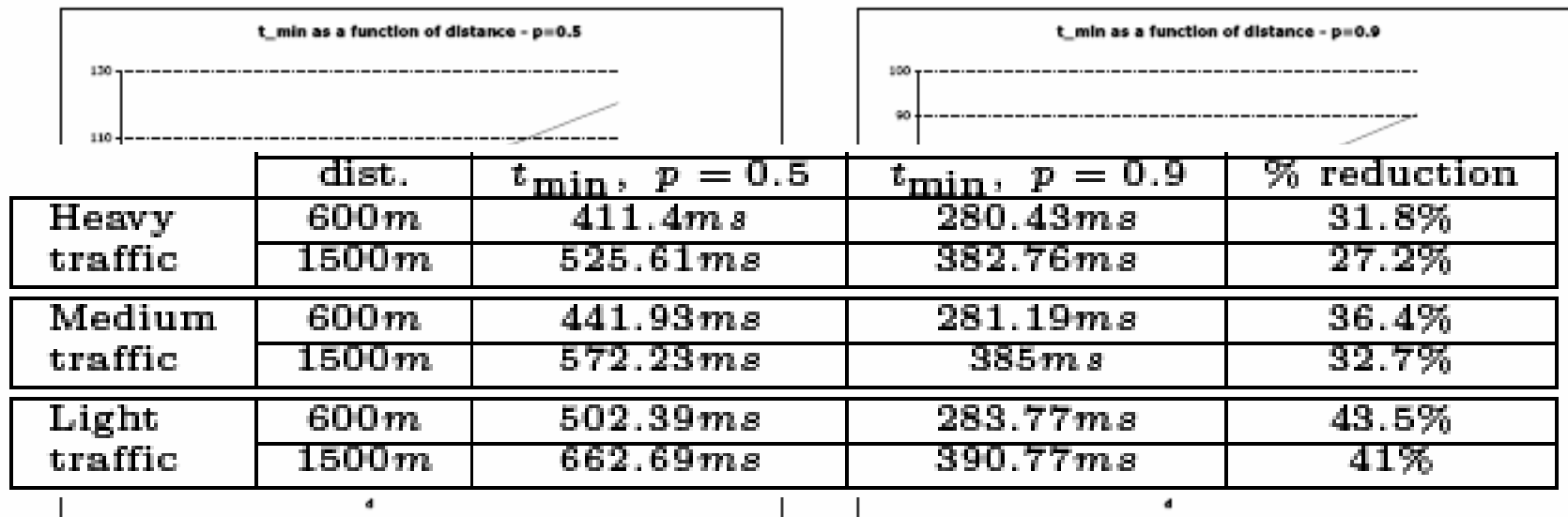
$$t_{\min}(\bar{d}, p) = \frac{\bar{h}E(p, r_T)^2 + (\bar{d} - 1)E(p, r_T) - E(\mathbf{X}(t)^2) \ln(1 - \sqrt{\bar{P}}) + \sqrt{E(\mathbf{X}(t)^2) \ln(1 - \sqrt{\bar{P}}) (-2\bar{d}E(p, r_T) + 2E(p, r_T) + E(\mathbf{X}(t)^2) \ln(1 - \sqrt{\bar{P}}))}}{E(p, r_T)^2}$$

$$\geq \sqrt{\bar{P}} \cdot \sum_{h=\bar{h}}^{\bar{t} - \frac{\bar{d}}{r_T}} \text{Prob}(\text{EStart}(\bar{d}, \bar{t} - h)) = \sqrt{\bar{P}} \cdot \boxed{\text{Prob}(\text{Start}(\bar{d}, \bar{t} - \bar{h}))}$$

$$\geq \sqrt{\bar{P}}$$

DISCUSSION AND SIMULATION(1/3)

■ Dependence on distance :



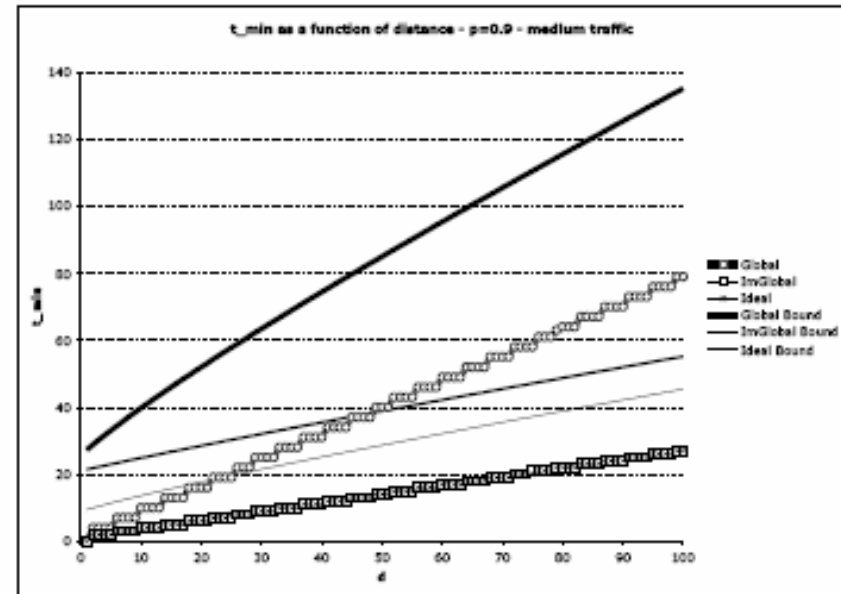
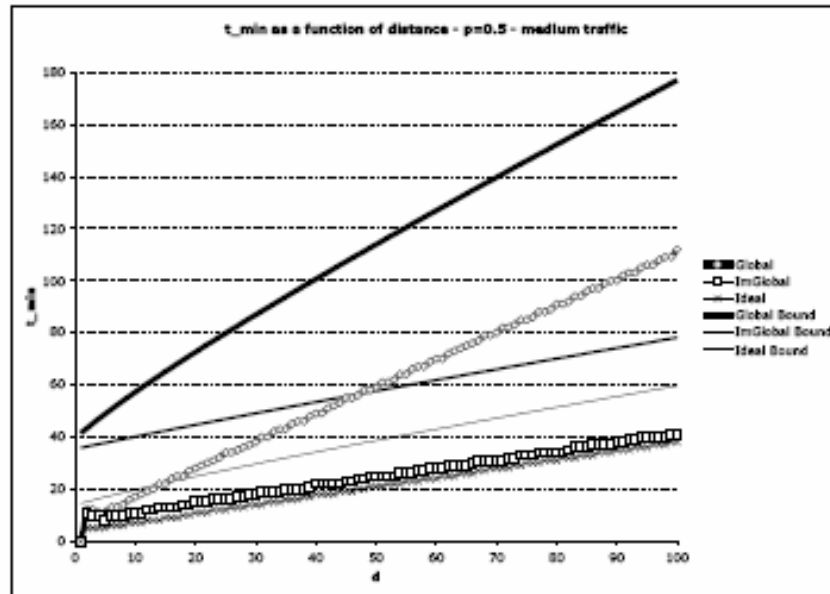
Values of t_{min} for fixed values of p and varying values of \bar{d} . Medium traffic scenario

DISCUSSION AND SIMULATION(2/3)

- Dependence on traffic conditions :

	dist.	$t_{\min, p = 0.5}$	$t_{\min, p = 0.9}$	% reduction
Heavy traffic	600m	411.4ms	280.43ms	31.8%
	1500m	525.61ms	382.76ms	27.2%
Medium traffic	600m	441.93ms	281.19ms	36.4%
	1500m	572.23ms	385ms	32.7%
Light traffic	600m	502.39ms	283.77ms	43.5%
	1500m	662.69ms	390.77ms	41%

DISCUSSION AND SIMULATION(3/3)



Values of t_{min} for fixed values of p and varying values of \bar{d} . Medium traffic scenario.

- Upper bound is accurate.
- When p is high, IMGLOBAL performs as IDEALIZED.



CONCLUSIONS

- Beneficial effect of Increased 1-hop reliability tends to decrease as the distance from the initiator increases.
- Benefit on multi-hop reliability of having high 1-hop reliability tends to decrease as car density increases.
- The dissemination strategy has a major impact on multi-hop reliability, especially when p is high, fast backward is more important.