

# A Distributed LSP Mechanism to Reduce Spare Bandwidth in MPLS Networks



ICC'2002

報告人: 唐崇實, 9/13/2002



# Outline

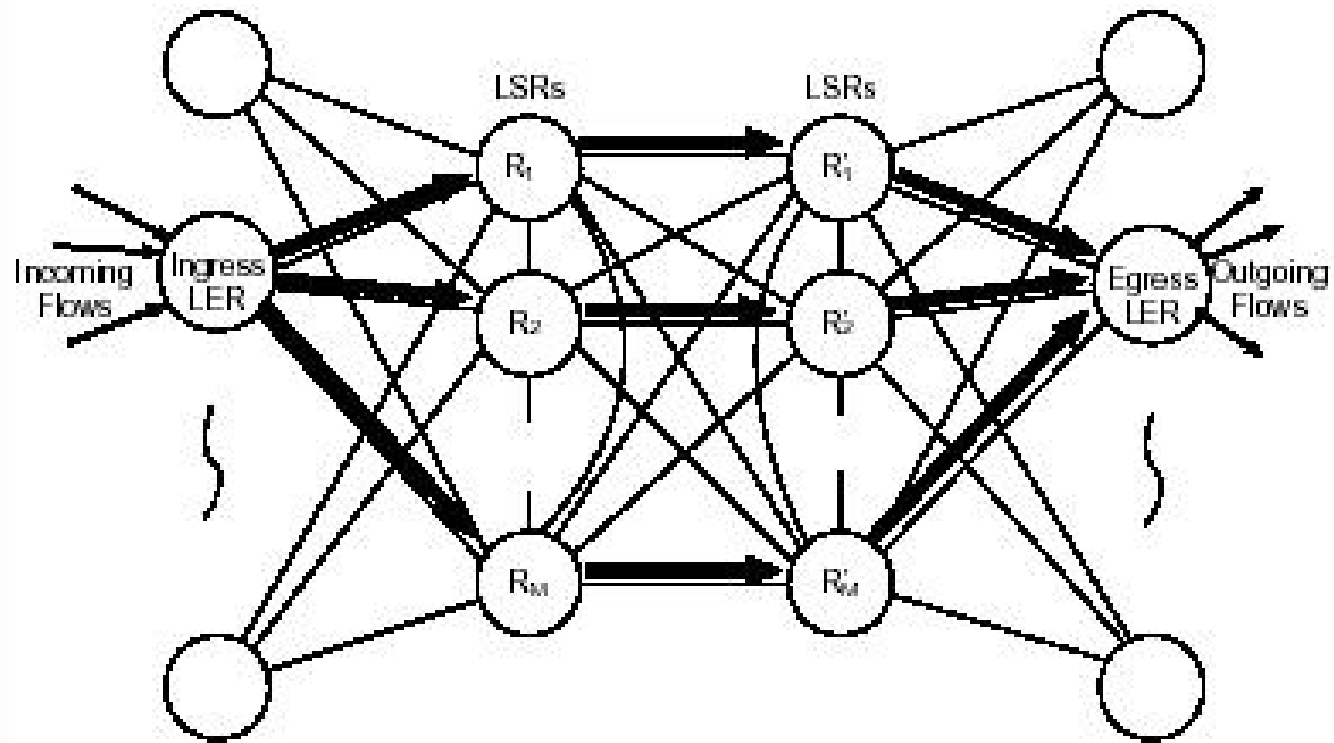
- Key Idea: Distributed LSP
- LSP Assigning Mechanism
- Spare Bandwidth Demand
- Effect of LSP Partitioning
- Conclusions



## Key Idea

- **Distributed LSP (D-LSP)**
  - An LSP is partitioned into several sub-LSPs.
  - Each sub-LSP is distributed to a different node-disjoint route.
  - The incoming traffic is also partitioned and spread over the sub-LSPs at the ingress LSR.
  - The backup LSP for a D-LSP is established on another node-disjoint route.
  - The amount of spare bandwidth is equal to the bandwidth demanded by a sub-LSP

# A D-LSP Example





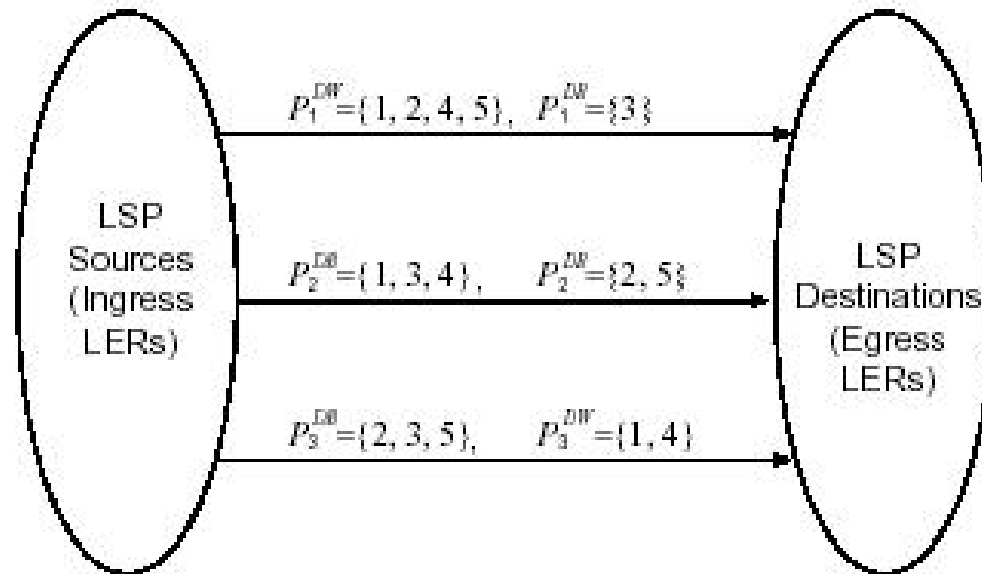
## LSP Assigning Mechanism (1/3)

- For D-LSP, assume that
  - $L$  D-LSPs originates from ingress LSRs to egress LSRs through the  $M$  parallel paths
  - $P_m^{DW}$  : a set of “working D-LSPs” of which a sub-LSP is assigned to the node-disjoint route  $m$
  - $P_m^{DB}$  : a set of D-LSPs for which a backup LSP is assigned to the node-disjoint route  $m$
  - Backup LSPs are distributed over node-disjoint routes uniformly 
$$|P_m^{DB}| = \begin{cases} \lfloor L / M \rfloor, & \text{on } M - R \text{ routes} \\ \lceil L / M \rceil, & \text{on } R \text{ routes} \end{cases}$$

## LSP Assigning Mechanism (2/3)

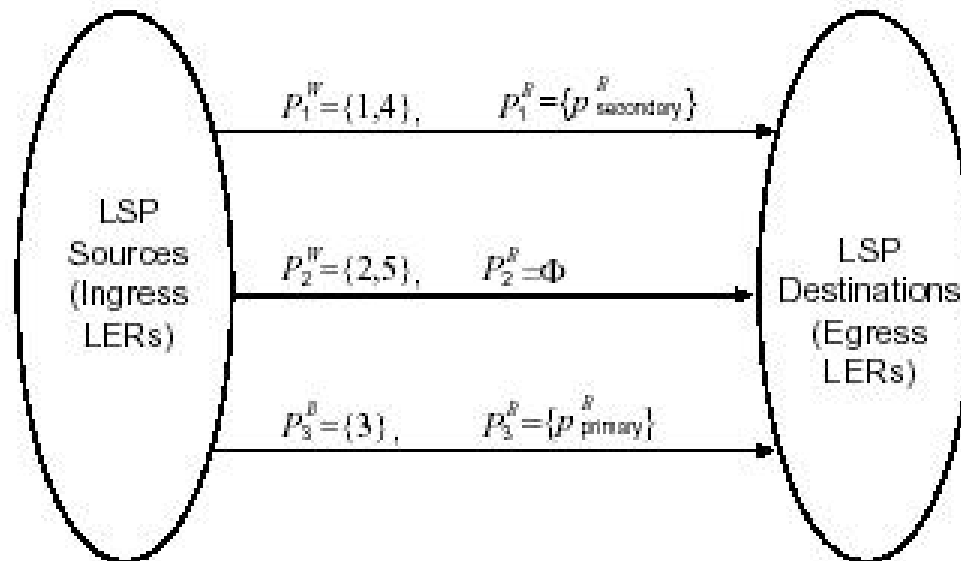
- For each D-LSP, there is only one backup LSP

$$\sum_{m=1}^M |P_m^{DB} \cap \{p\}| = 1, \quad \text{for } p = 1, \dots, L$$



# LSP Assigning Mechanism (3/3)

- For conventional LSP
  - At least two backup LSPs are required



# Spare Bandwidth Demand

- For D-LSP, the total bandwidth demanded is

$$B_{total}^D = LMC'$$

- For conventional LSP, the total bandwidth is

$$B_{total}^C = (L + (2Q + \min(R,1)))C, \text{ where } L = QM + R$$

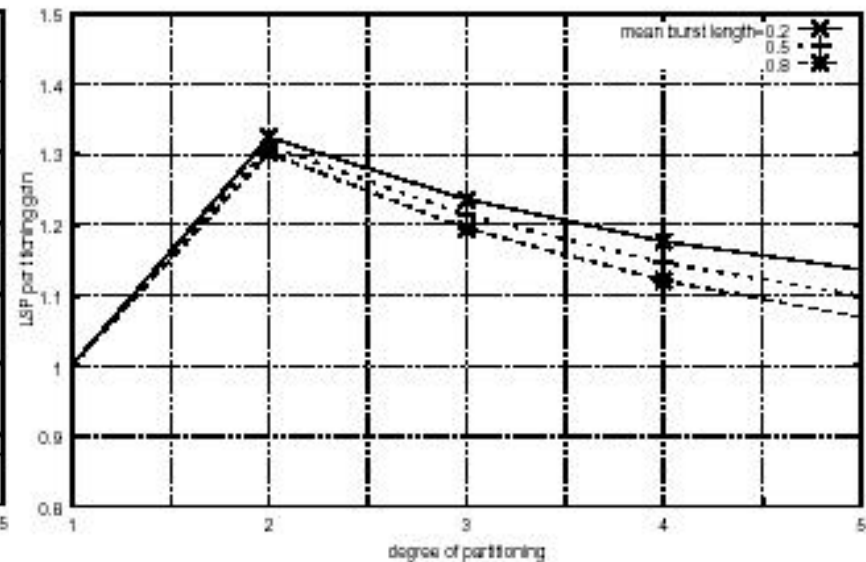
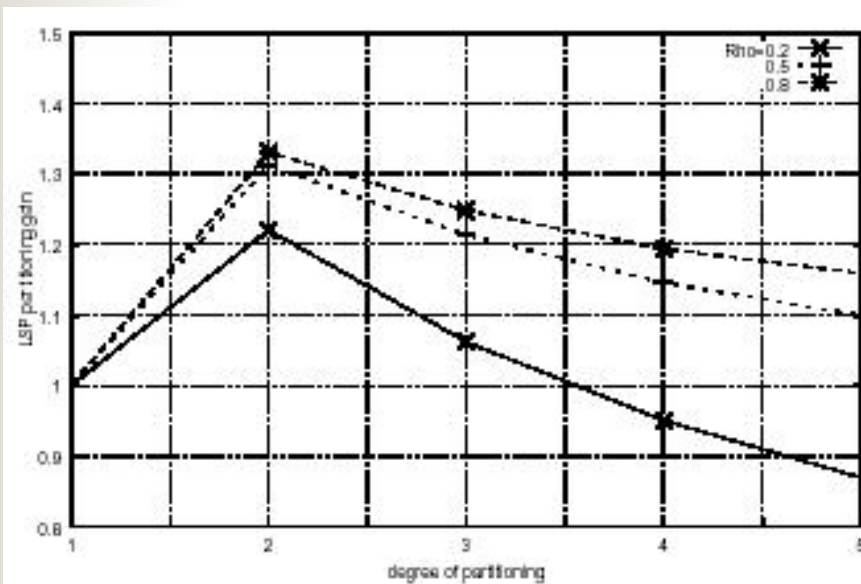
- LSP Partitioning Gain

$$\begin{aligned} \gamma &= \frac{\text{bandwidth demand of conventional LSP mechanism}}{\text{bandwidth demand of D - LSP mechanism}} = \frac{B_{total}^C}{B_{total}^D} \\ &= (L + 2Q + \min(R,1))C / LMC' \end{aligned}$$



# Effect of LSP Partitioning

- Considering ON-OFF traffic model ( $r, b, \rho$ )  
 $r$ : peak rate;  $b$ : mean burst length;  $\rho$ : flow source util.





# Conclusions

- A very intuitive and effective mechanism for reducing the spare bandwidth demand is proposed.
- There is a trade-off between LSP partitioning and statistical multiplexing gain.
- One possible extension of this study can be investigating the D-LSP mechanism for multi-QoS leveled MPLS networks.