

# Adapting Peer-to-Peer Topologies to Improve System Performance

Presented by  
Chi-Hung Chao



# Outline

- Introduction
- Topology Evolution Algorithm
- Balancing Bias Algorithm
- Experiment
- Analysis and Observation
- Conclusion and Discussion

# Introduction

- Improving the performance of unstructured Peer-to-Peer file sharing system.
- Controlling the routing protocol and network topology simultaneously.

# Introduction

- Flooding based search algorithms are very sensitive to the number of edges in the network graph.

# Topology Evolution Algorithm

- Edge Thinning
  - Reducing the total number of links
- Diameter Folding
  - Adding some useful links

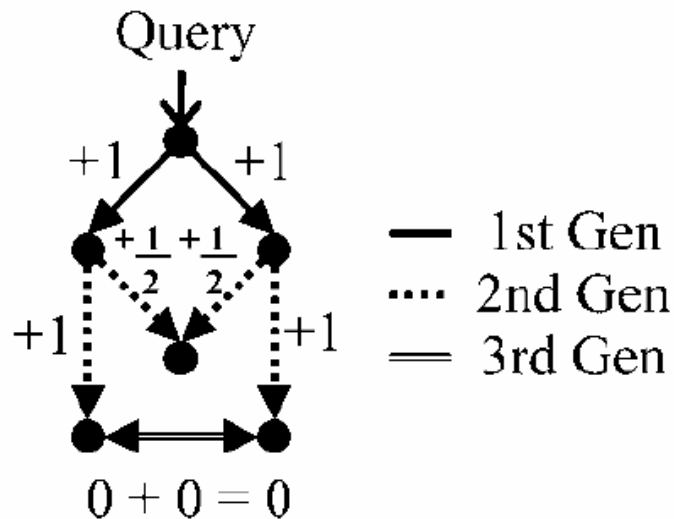
# Topology Evolution Algorithm

- Edge Thinning
  - Reducing search cost by removing the most redundant links without disconnecting the graph.
  - Assigning a “score” to each link within the network to determine how useful it was across all searches.

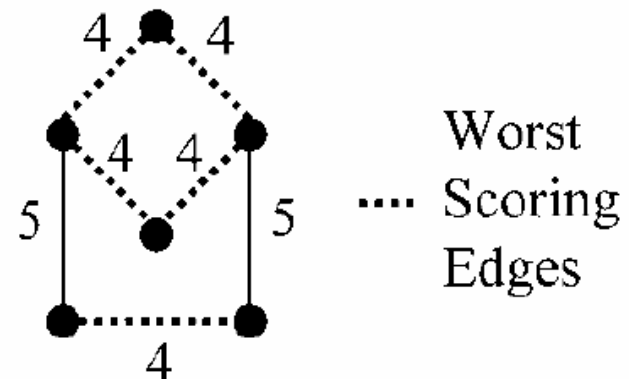
# Topology Evolution Algorithm

- Edge Thinning

Edge Scoring for a Query



Candidate Edges for Removal  
(Possible cumulative scores 0-6)



# Topology Evolution Algorithm

- Edge Thinning
  - Late message
    - Redundant messages received at a later time.
  - Even-length cycles also result in simultaneous redundant message delivery, but no late message are passed.
  - Some bandwidth is used without any benefit to search in these odd-length cycles.



# Topology Evolution Algorithm

- Edge Thinning
  - This calculation-removal cycle can be repeated until the graph becomes disconnected, at which point the last edge removed can be re-inserted.

# Topology Evolution Algorithm

- Diameter Folding
  - After having created a way to remove the most redundant links, we began considering how and when we might add useful links.

# Topology Evolution Algorithm

- Diameter Folding

- By breadth-first search protocol, the worst-case time for a search to reach all other nodes is the same as the graph diameter.
- One way to add shortcut links between all node-pairs whose minimum path length is the diameter.

# Topology Evolution Algorithm

- Diameter Folding
  - So we perform the Diameter Folding algorithm, and then perform thinning from the folded topology.
  - But an edge would often be added that was immediately removed by the next thinning step because the folding step was creating an odd cycle within the graph.

# Topology Evolution Algorithm

- Diameter Folding

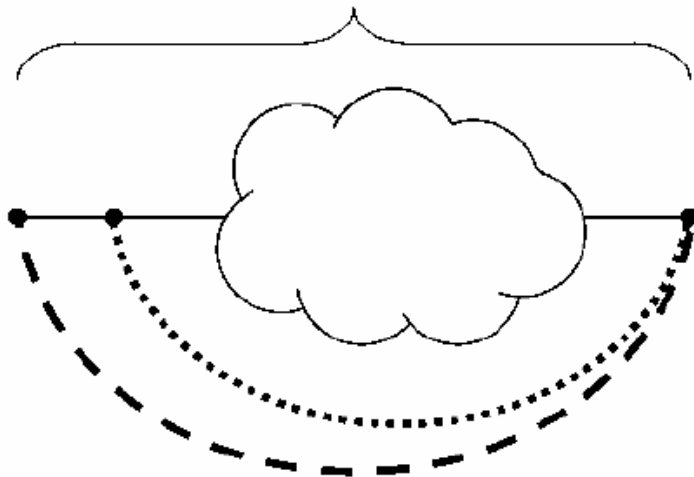
- Ensuring that the new cycle that is being created is always of even length
  - Folding operation selects a pair of nodes at the diameter and connects them directly if the diameter is odd.
  - If the diameter is even, connects one of them to a neighbor along the path.

# Topology Evolution Algorithm

- Diameter Folding

Greatest Odd-Length Paths for Folding

Diameter



-- Odd Diameter Fold

..... Even Diameter Fold

# Balancing Bias Algorithm

- To balance nodes degrees
- In thinning
  - Remove the poor scoring links between nodes with high degrees.
- In Folding
  - Add links between distant nodes with low degree.

# Experiments

- Experiment 1
  - Fixed number of links.
- Experiment 2
  - Allowed the number of links to increase or decrease.



# Experiment 1

Graph Size		Target Mean Degree	Actual Mean Degree	Initial Degree Statistics			Degree statistics after 5000 steps of alternating thin and fold					
							Without Balancing Bias			With Balancing Bias		
Nodes	Links			Min	Max	Std	Min	Max	Std	Min	Max	Std
32	66	5	4.1	1	8	2.0	2	17	4.4	2	8	1.4
64	178	6	5.6	1	10	2.1	3	33	7.6	2	9	1.7
128	455	7	7.1	2	15	2.8	3	54	9.7	4	15	2.0
256	1020	8	8.0	1	17	2.9	4	59	10.5	4	17	2.7
512	2371	9	9.3	1	20	3.1	5	37	5.6	6	21	2.3

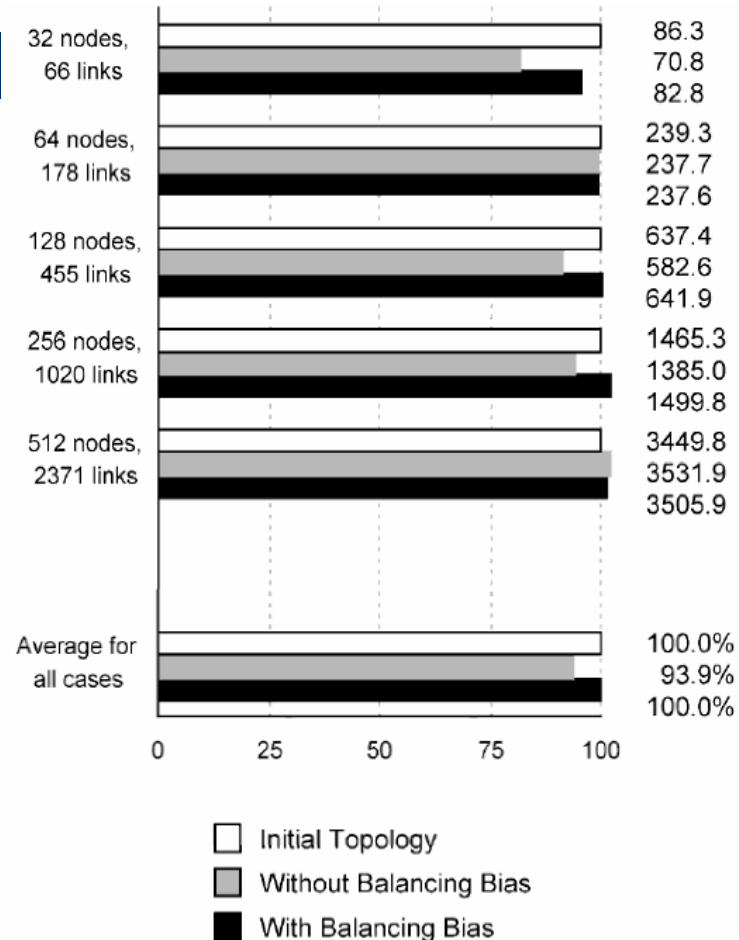
NOTE: Random graphs generated by probabilistically deciding whether to include each possible link, with the probability determined by target mean degree. Here, the target degrees mimic a base two hypercube of the same size.

# Experiment 2

Initial Graph Topology	Degree			Path			Search		
	Min	Mean	Max	Min	Mean	Max	Min	Mean	Max
Random Regular: degree 3	3	3.0	3	1	5.1	9	219	226.9	235
	1	3.0	5	1	4.6	7	213	222.6	233
Random Regular: degree 7	7	7.0	7	1	2.7	4	594	634.5	666
	1	3.0	7	1	4.2	7	210	221.9	232
Random Net	1	7.1	13	1	2.7	5	618	645.6	673
	1	3.0	13	1	3.8	7	218	225.9	233

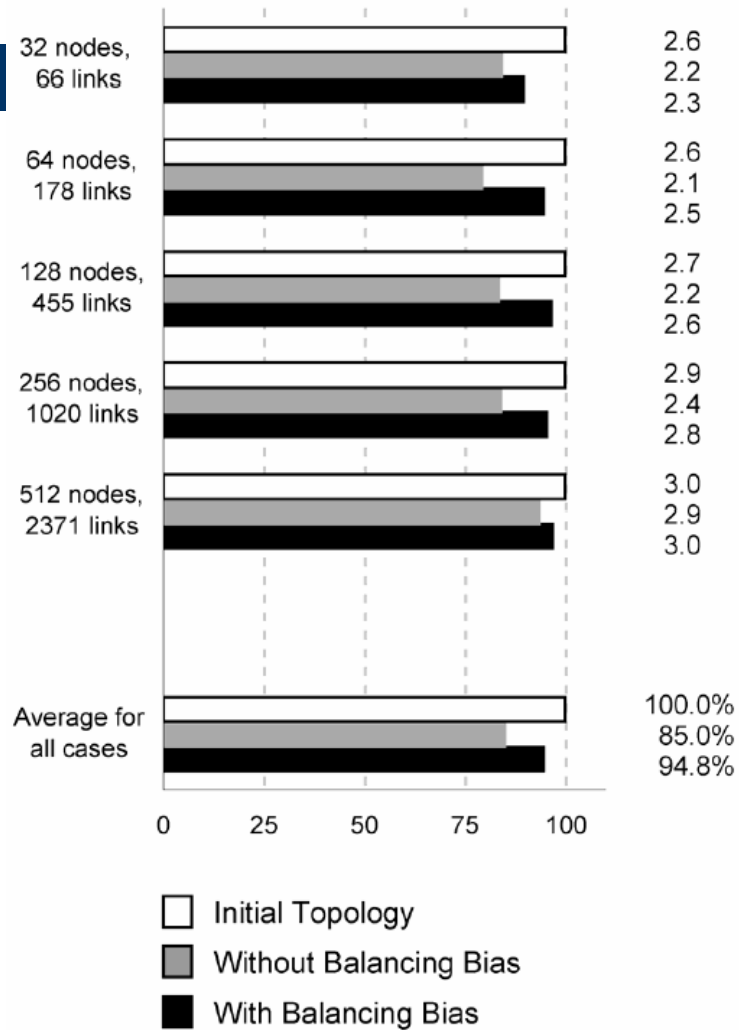
# Analysis and Observations

- Bandwidth Usage



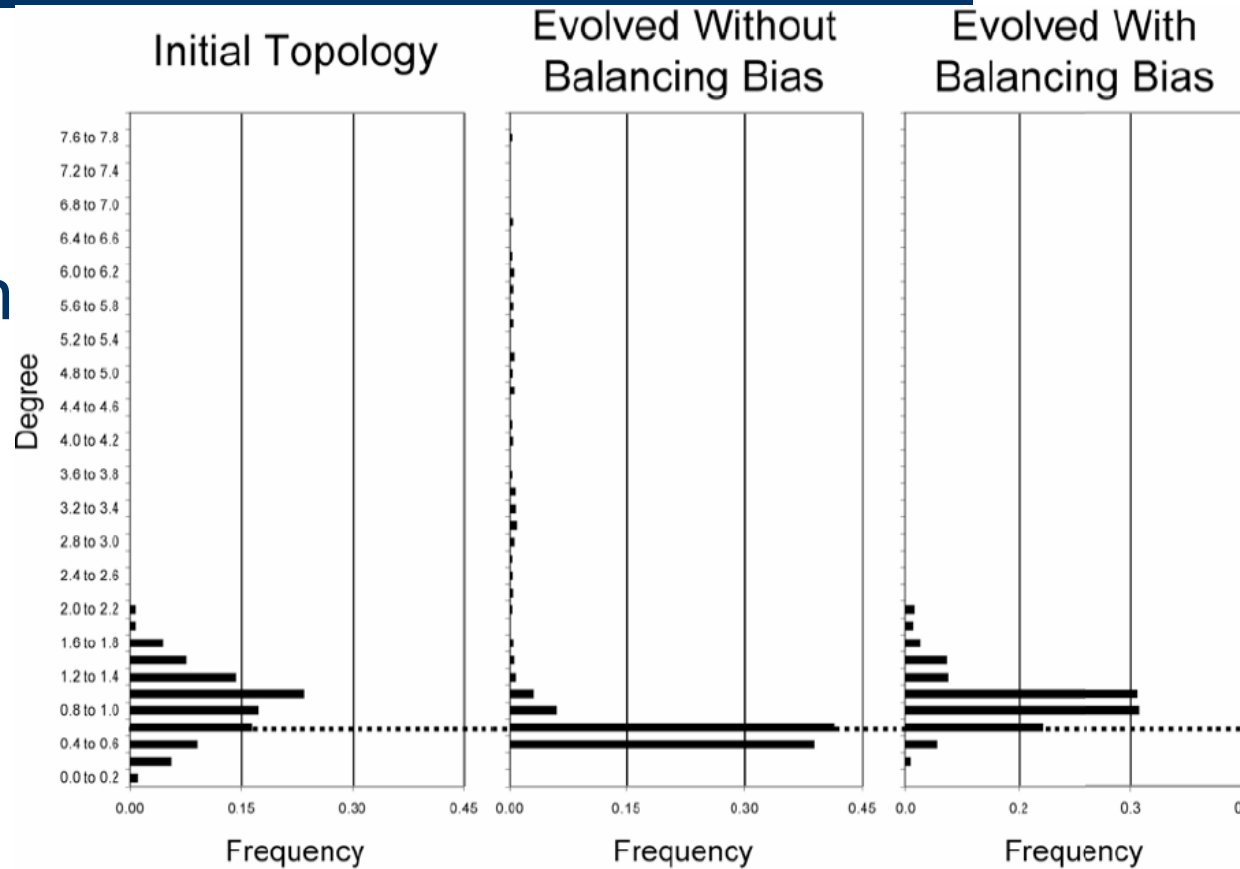
# Analysis and Observations

- Search Time



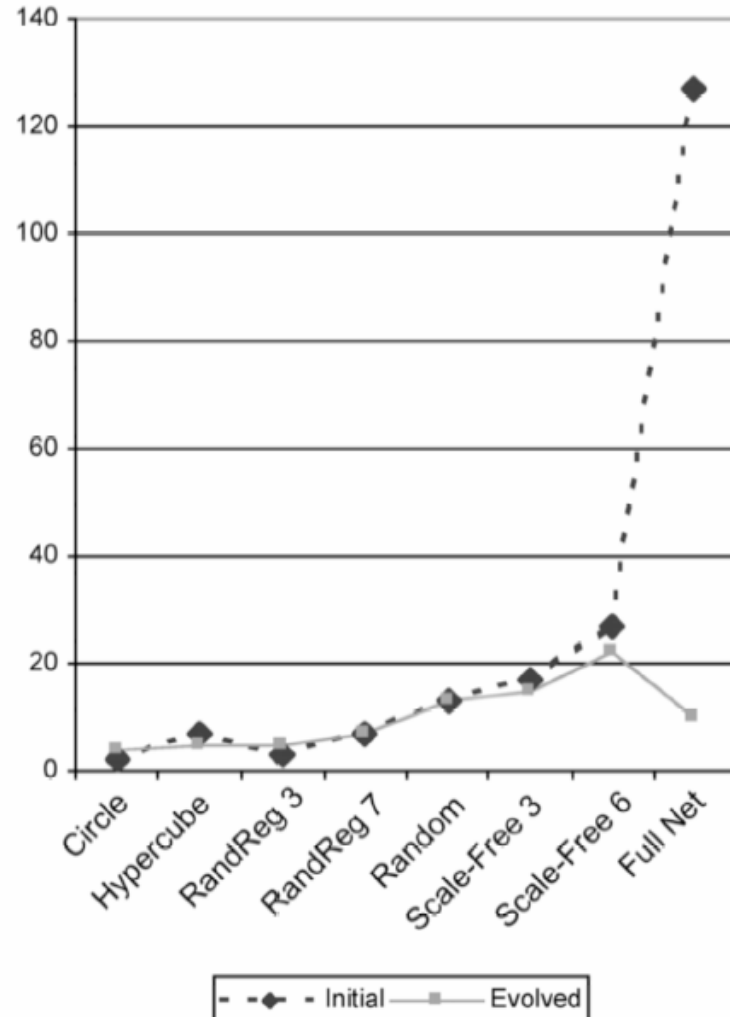
# Analysis and Observations

- Average Degree Distribution



# Analysis and Observations

- Maximum Degree



# Conclusions and Discussions

- These algorithm will lead to network performance improvements.
- But the algorithm need to be performed in a stable environment, and the bandwidth consumption of performing this algorithm maybe large.