Secure and Reliable Decentralized Peer-to-peer Web Cache

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Outline

- Decentralized Web Cache
- Performance Evaluation
- Conclusion

- Web caching is a well developed scheme for improving the performance of web browsing.
- Users' requests can be satisfied by the local cache or the cache of a nearby proxy, instead of a remote web server.
- Performance Studies[3,11,22].
- Various proxy caching techniques and replacement[1,4,12,14,22].
- Cooperative caching architectures and mechanism[5,15,21,23].

- One class of the cooperative caching strategies is based on the client-side caching.
 - In a LAN environment, the clients' cache space can be organized and utilized to reduce the total outgoing traffic.
 - A client may get web contents from other clients.
 - Simulation results indicate that sharing cache among clients can significantly decrease the external traffic.
 - However, client-side caching has some fundamental problems:
 - Availabilities of the client platforms are unpredictable.
 - The privacy of users' access activities should be protected.

- Centralized peer-to-peer web caching protocol[24].
- Centralized approaches incur a higher management cost, especially in a large domain.



- Decentralized peer-to-peer web cache mechanism – Squirrel[9].
- Squirrel considers two document look-up protocols, home store and directory.

Home store



Directory



- Privacy issues are not considered in either of the look-up protocols.
- Though these schemes yield a lower external bandwidth, they incur a much higher internal traffic.
- The high hit ratio in these schemes are obtained based on some idealistic assumptions that may not be satisfied in a real environment, such as low failure rate and some necessary operations when nodes join or leave the system.

- This paper presents a web caching model based on the decentralized peer-to-peer architecture.
- Using a hybrid policy that combines existing home store and directory approaches.
- The combined approach eliminates the problems in the original home store and directory schemes.
- Developing new algorithms and policies to resolve certain problems that arise due to the combination.
- A new cache replacement policy based on the hybrid approach has been developed.
- Encoding the web contents and hiding IP information to protect clients' privacy.

Scheme 1	no client cache share, proxy cache
Scheme 2	client cache share, no proxy cache
Scheme 3	client cache share, proxy cache

- The home store method represents scheme 1.
- The directory approach yields scheme 2.
- The new design in this paper is based on scheme 3.

Splitting the client cache space into two parts.

- Local cache and Home node cache
- □ *psize* denote the Home node cache size
- Isize denote the Local cache size
- Both *psize* and *lsize* are arguments that can be set by users.

Access Request Processing

- While a client C requests a web object, it first checks its local cache. If the object is not found, a request message will be forward to the home node H.
- If the home node has cached a copy, it sends the object to C. Otherwise, H randomly pick a delegate D from the directory for the transferring. If the directory is empty, the requesting client will fetch the object from the original server.



Cache Replacement

- Remove pages with expired TTL.
- When the home node evicts a web object it hosts, it informs all holders of that object so that they can increase the priority of that object in their local cache.
- On the other hand, when a client evicts a hot page, if it is the last copy or one of the last few copies, it sends the page to the home node.

Node Load Smoothing

- It's important to keep the load evenly distributed among them.
- In the directory approach, the number of served objects per second can be burst.
- A stricter solution for load balancing is to use a reference counter. Each node keeps an upper bound *rlimit,* which represents the maximum number of nodes it can be referenced by.
 - When a home node H wants to add a client C to its directory, H first sends a message to C. C checks its current value of the reference counter. If it's less then rlimit, C sends an 'accept' response to H and increases the reference counter. Otherwise, C sends ''deny response to H.

Privacy and Security Issues

- In web caching systems, a node may not want the others to trace its access history.
- In the centralized scheme, this could be done by the trusted proxy servers. However, it's more complex in this self-organizing approach.
- (IP, Hash(URL) / (P)_k)
 K=f(URL)

Performance Evaluation

Access Frequency



Performance Evaluation

Performance

- We assume that each node has the same size of cache space and we assign the size to 10K, 100K, 1M, 5M, 10M and 20M.
- We set *psize* to 30% of the total cache size and TTL to one hour.

Performance Evaluation

Performance



Conclusion

- The experimental study shows that it performs better than the existing schemes in term of hit ratio and node load.
- It may be necessary to partition these objects and cache them on multiple client platforms to avoid occupying a large cache space of a single client and facilitate load smoothing.