Clustered P2P Architectures

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Outline

- Unstructured P2P
- DHT P2P (structured)
- Clustered P2P (hierarchically structured)
- ECSP/EFA/Performance
- Conclusions
- Hybrid P2P Issues

Unstructured P2P

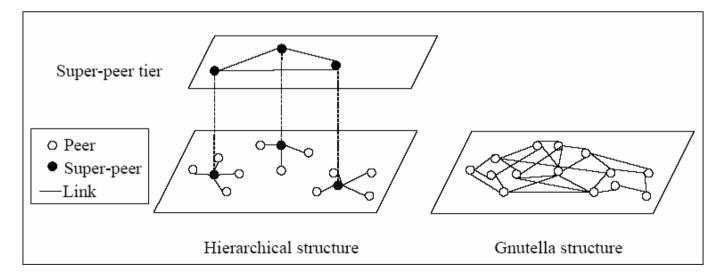
- Gnutella
- Adv.
 - resilience to dynamic peer join/leave
 - no overhead on peer failure
 - support for keyword search
- Disadv.
 - overhead of query messages
 - no guarantee on availability

DHT P2P

- DHT: Distributed Hash Table
- Pastry, Chord, CAN
- Adv.
 - scalable compared with unstructured P2P
 - guaranteed availability
- Disadv.
 - control overhead on frequent peer join/leave
 - only support for exact name match

Clustered P2P

- eDonkey[3], Kazaa[4]
- Clustered P2P systems improve message overhead of unstructured P2P without introducing control overhead as in DHT P2P.

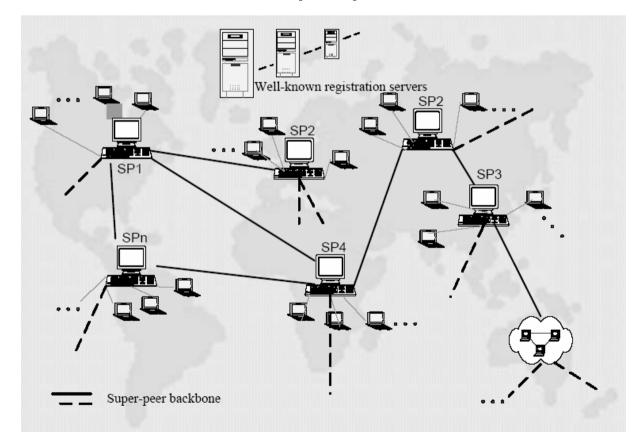


Clustered P2P

- Peers are grouped into clusters and connected to the superpeer of cluster.
- Peers with more resources, higher processing and network capacities are selected as superpeers of clusters.
- Superpeers act as local search hubs, building indices of the files share by each peer connected to them, and proxying search requests on behalf of theses peers.

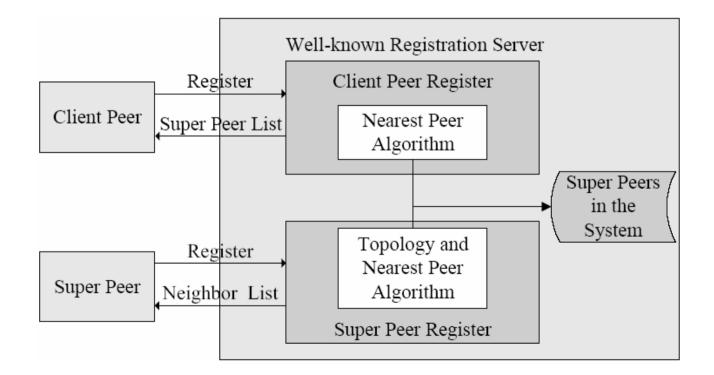
ECSP[1]

• Efficient Clustered Superpeer P2P



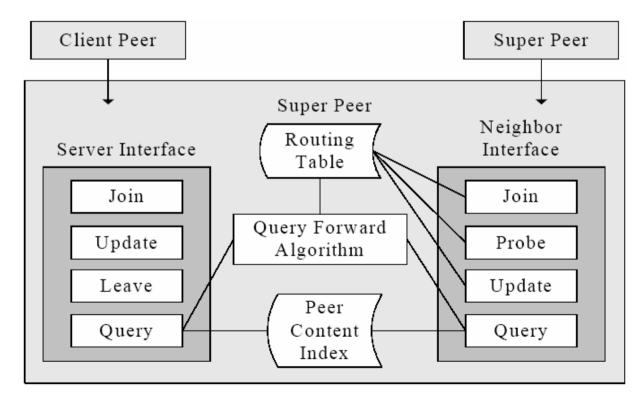
Registration Server

• Registration servers supply yellow page services to all nodes in the network.

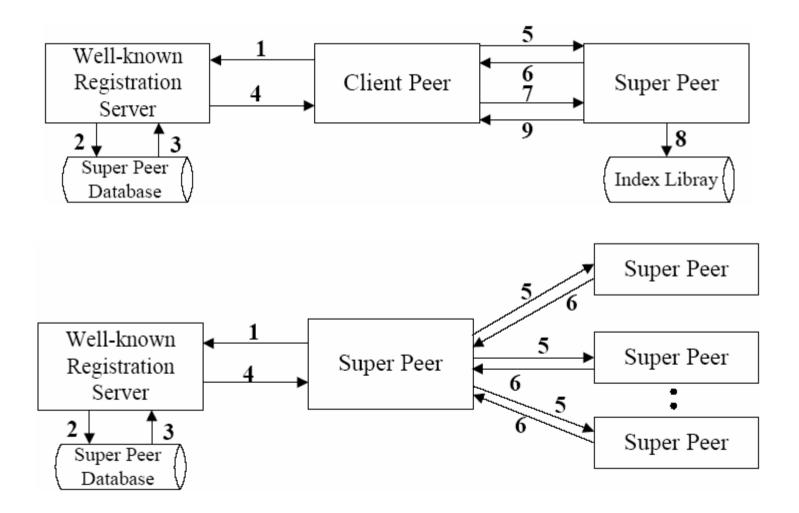


Superpeer

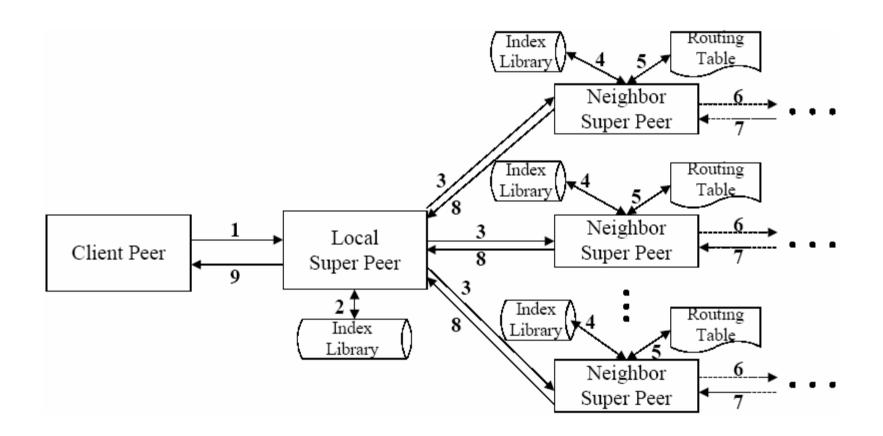
• Superpeers act as cluster leaders and service providers for peers in their clusters.



Joining of Peers



Querying



EFA

- Efficient Flooding Algorithm
- EFA use limited topology information and simple computing to decrease the duplication queries created by flooding.
- If node *v* can anticipate that one of its neighbor *u*, receives query messages from another path, however, then *v* does not forward the query to *u*.

EFA

Symbol	Description
v	Current node
id(v)	Node v's unique id
N(v)	Neighbor set of v
NN(v)	Neighbor's neighbor set of v
fr(u,v)	<i>v</i> is the current node, <i>u</i> is the node which forwards the query to <i>v</i> . $fr(u,v)$ is the forward reaching set of <i>u</i> for the current node <i>v</i> , i.e. the immediate (no more than 2 hops away) set of nodes reached by the local flooding source <i>u</i> .
routing(u,v)	For local source u , current node v 's routing set. For example, if u forwards the query package to v , the set of nodes v forwards is decided by <i>routing</i> (u, v)

EFA

forward(u,v)

/*when node v receives forwarded query from its neighbor u, this algorithm decides how v forwards this query */

If the received query has been received before discard it

else

if u is null / v is the node which initiates the query*/ forward the query to N(v)*

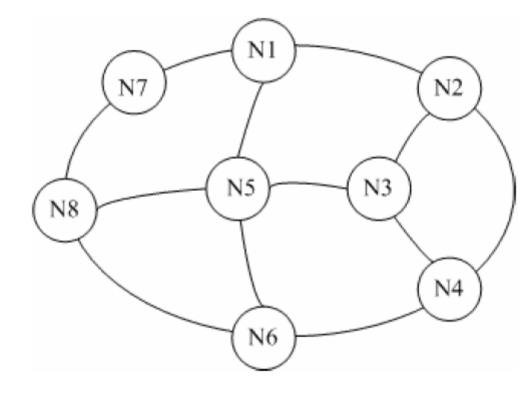
else

forward the query to routing(u,v)

(a)

 $fr(u,v) = N(u) \cup \{ \text{ all } v' \text{ in } NN(u) \mid id(v') \le id(v) \}$ routing(u,v) = all v' in N(v), such that $1. \quad v' \notin fr(u,v) \text{ AND}$ $2. \quad \{N(v') \cap fr(u,v) = \emptyset\} \text{ OR } \{N(v') \cap fr(u,v) = A \text{ AND } (\forall v'' \in A \text{ AND } id(v'') > id(v)) \}$ (\mathbf{b})

Example



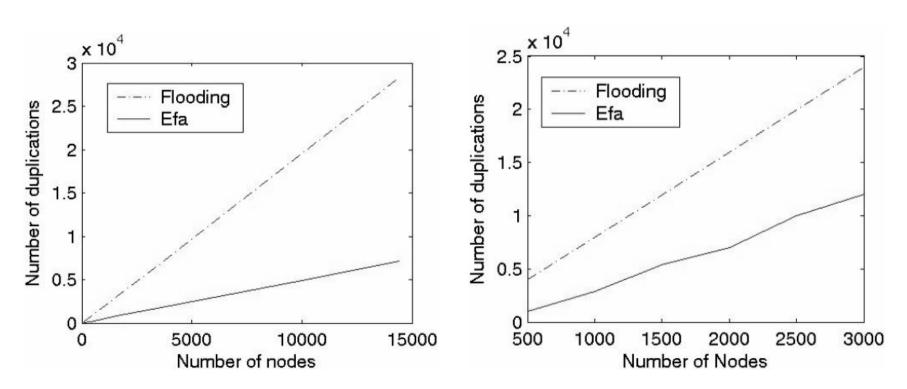
Forward(N1, N5) N(N1)=N2, N7 NN(N1)=N3, N4, N8 fr(N1, N5)=N2, N3, N4, N7

checking 1. = N6, N8 checking 2a. = *NONE* checking 2b = N8

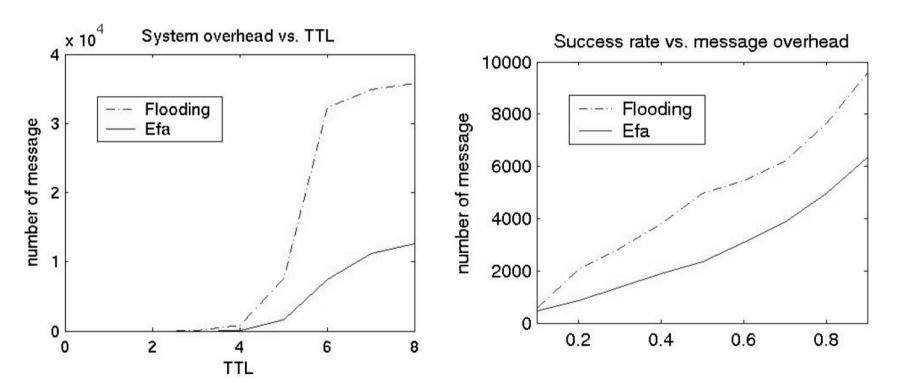
→routing(N1, N5)=N8

Performance

• Grid v.s. Random



Performance



Conclusions

- Hierarchical clustered P2P architecture can improve flooding overhead in pure unstructured P2P.
- EFA prevents duplication instead of discarding duplicate messages as in flooding.
- This work resembles our GP2P project in the two-level architecture.

Hybrid P2P

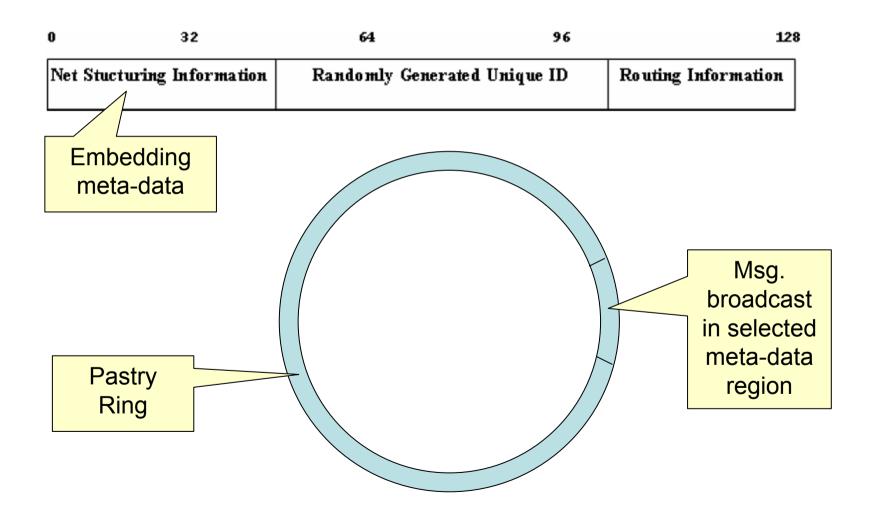
 Hybrid P2P models is proposed in [2] that integrates unstructured and DHT to enable network restructuring and routing behavior adaptation.

CAYLEY TREE OVERLAY

STRUCTURED DECENTRALISED ROUTING / LOOKUP NETWORK

TCP /IP

Reflective Key



Discussions

- How to use meta-data?
 RDF, File types
- How to define meta-data region?
 - Hashing
 - Explicit v.s. implicit
- How to define a peer of certain specialty?
 Interest groups, content analysis
- Multiple overlays v.s. multiple keys?

References

- 1. J. Li and S. Vuong, An Efficient Clustered Architecture for P2P Networks, IEEE AINA'04.
- 2. D. Hughes, I. Warren, and G. Coulson, Improving QoS for Peer-to-Peer Applications through Adaptation, FTDCS'04.
- 3. eDonkey. http://www.edonkey.com/
- 4. KaZaA. http://www.kazaa.com/