

# Cooperative and Opportunistic Transmission for Wireless Ad Hoc Networks

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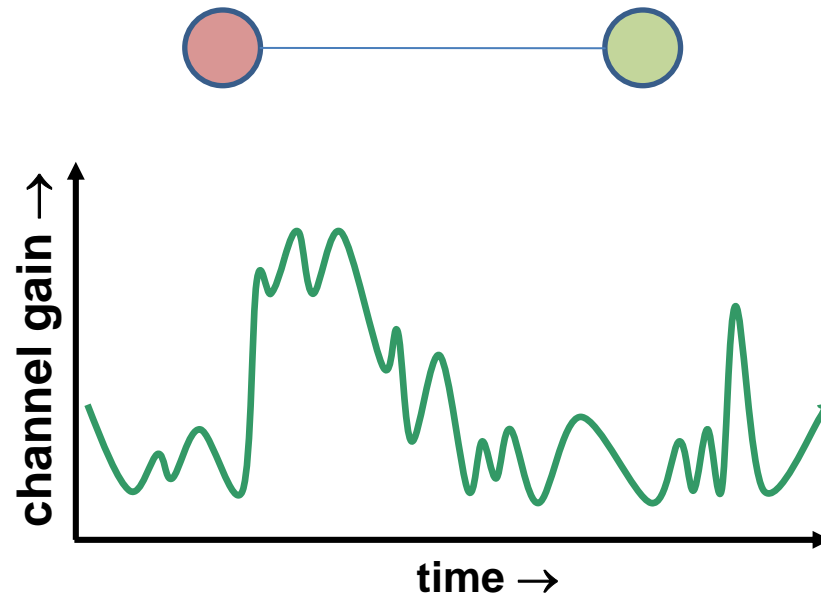
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# Outline

- Introduction
- Distributed Cooperative Rate Adaption (DCRA)
  - DCRA Performance
- Cooperative and Opportunistic Scheduling (COS)
  - COS Performance
- Conclusion

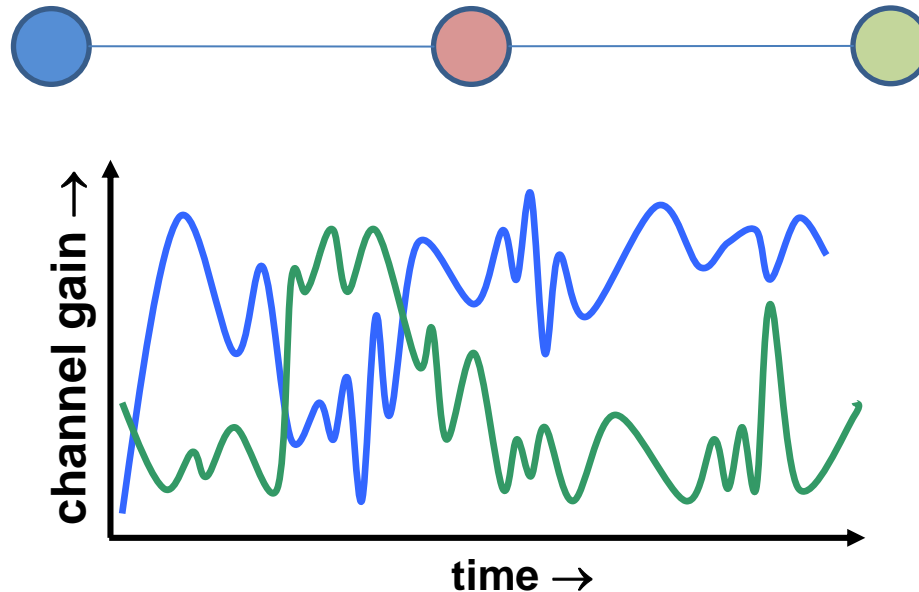
# Introduction

- Opportunistic Transmission
  - Time diversity: Transmit more packets at higher rates when the channel condition is better.



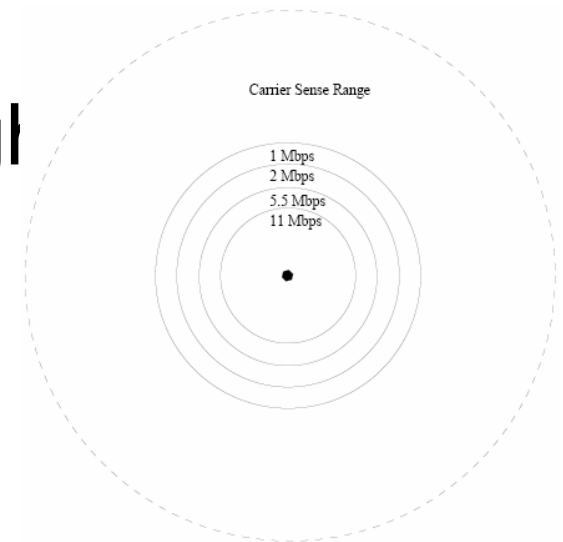
# Introduction

- Opportunistic Transmission
  - Multi-user diversity: Select instantaneously an on-peak receiver with the best channel condition.



# Motivation

- Local decision is not enough
  - Hidden terminal
  - Co-channel interference
- Neighboring transmitter should jointly determine the on-peak flows.
  - Rate selection
  - Flow scheduling

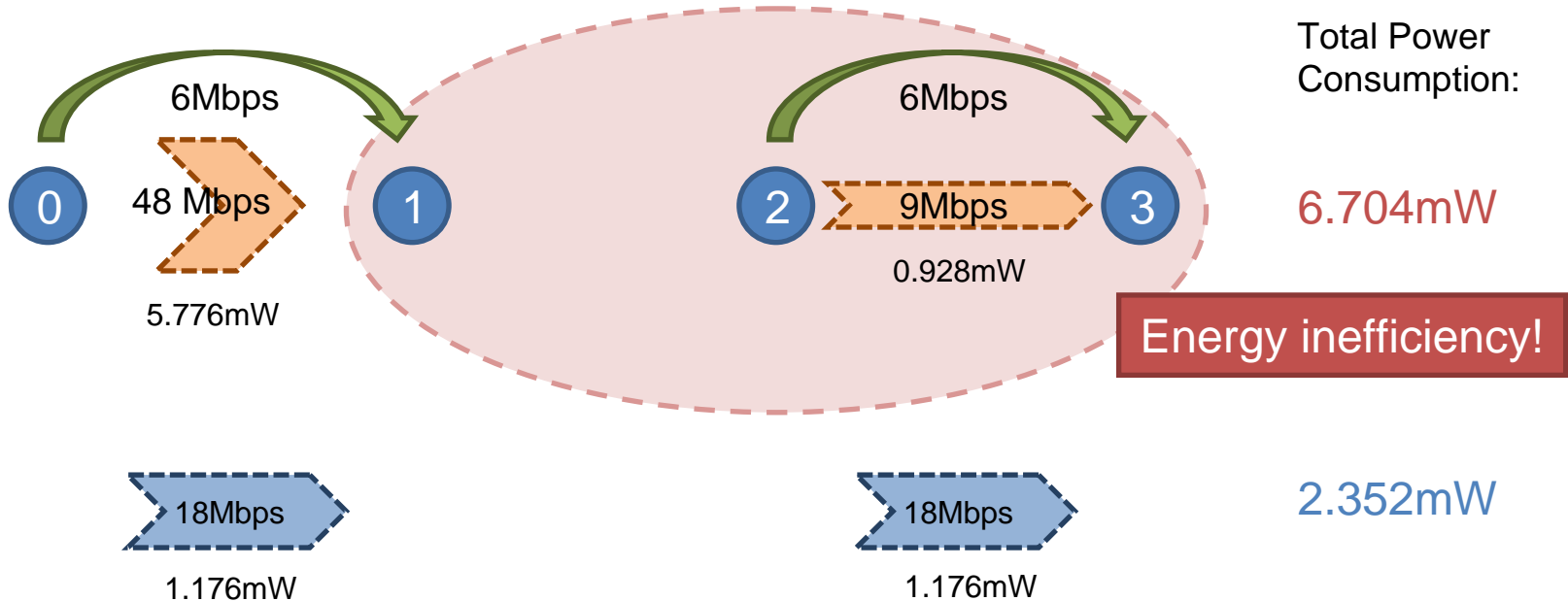


# Design Goals

- Cooperative and opportunistic transmission
- Cooperative rate adaptation (CRA)
  - Energy efficiency
- Cooperative and opportunistic scheduling (COS)
  - Throughput maximization
  - QoS provisioning

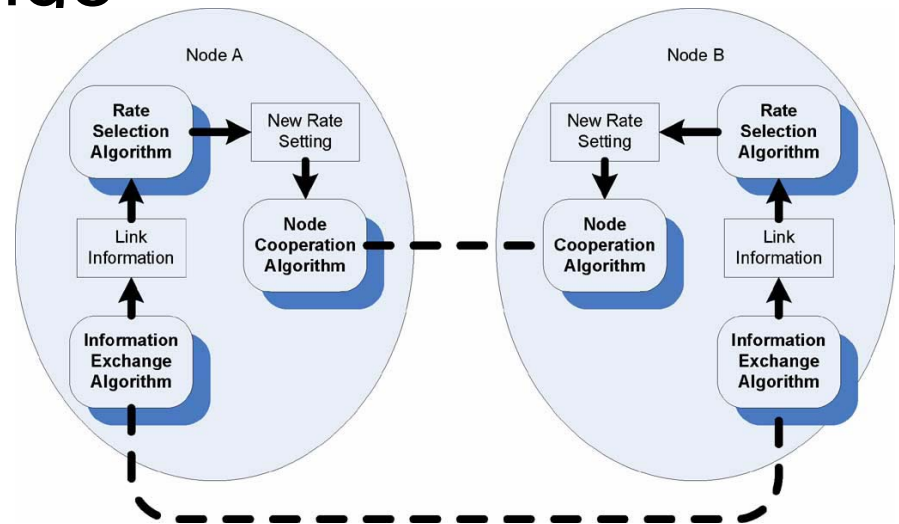
# Inequality of Channel Access

- Inequality of channel access results from hidden terminals.



# Distributed CRA

- Nodes cooperate in rate adaption for high overall energy efficiency.
- Information exchange
- Rate selection
  - (NP-complete)
- Node cooperation



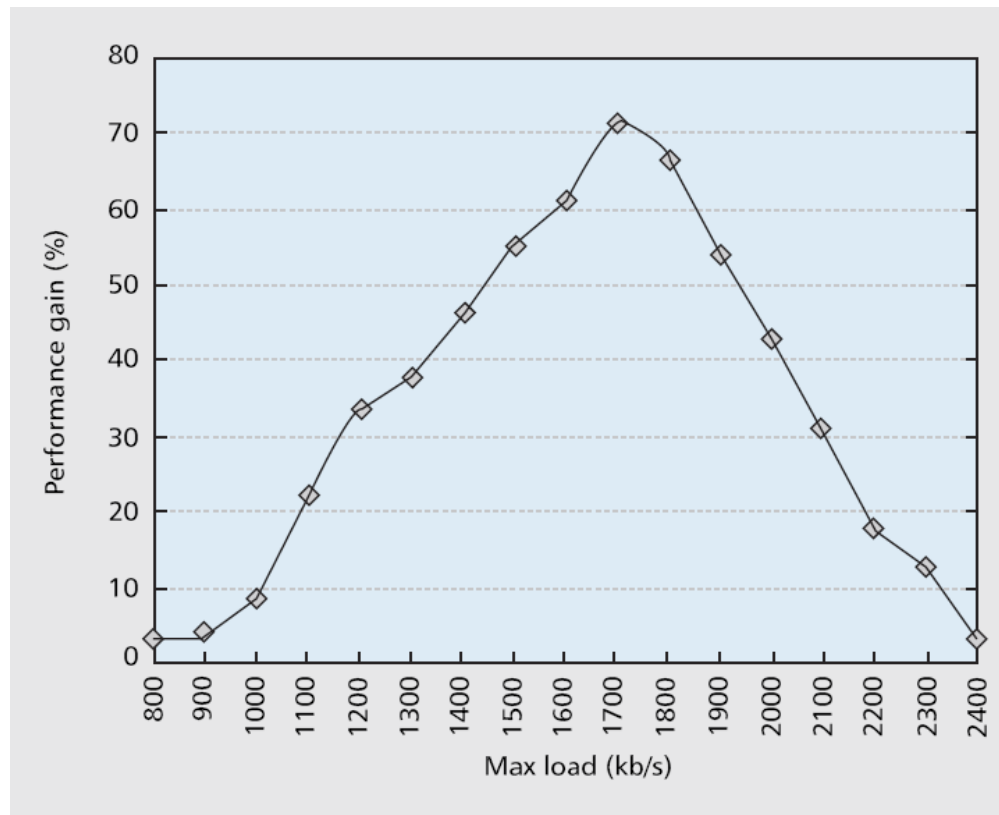


# DCRA (Heuristic Algorithm)

- Step 1 Set the rate for each link in node A's maximum interference range to the highest value as the initial setting.
- Step 2 For each link within A's maximum interference range, select a rate that has the largest  $\Delta E/\Delta T$ , where  $\Delta E$  denotes energy reduction and  $\Delta T$  denotes the channel time increase, as compared to the current setting. Then, choose the link that has the largest  $\Delta E/\Delta T$  among all the links within A's maximum interference range. If we can not find a setting that results in  $\Delta E > 0$ , the algorithm ends.
- Step 3 Check whether the new rate of the link is feasible. If it is feasible, select the new rate setting; otherwise, reset to the previous setting.
- Step 4 Go to step 2.

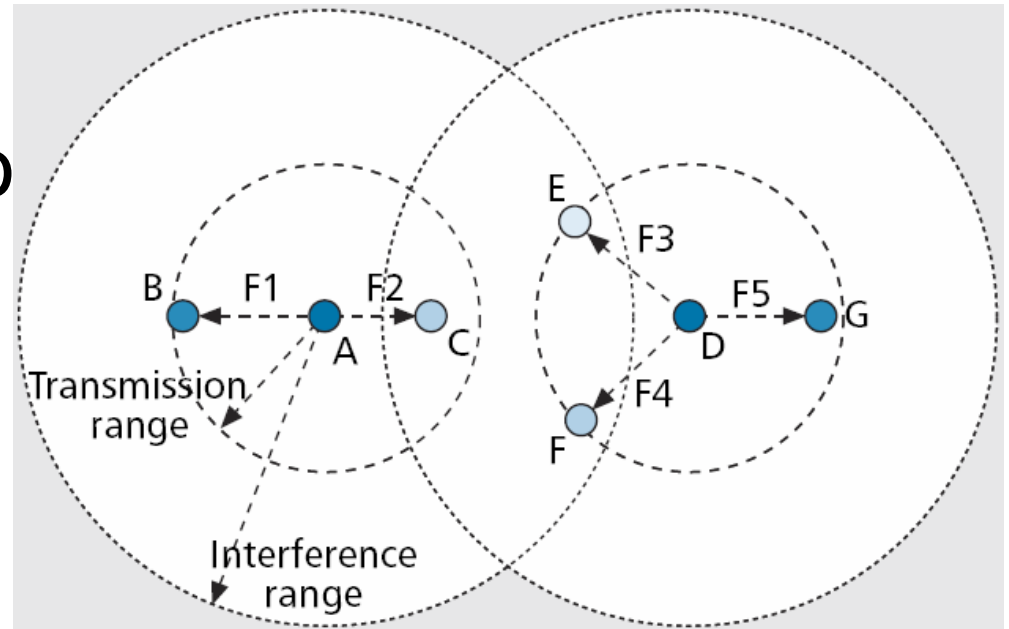
# DCRA Performance

- 1000m \* 1000m / 50 \* 50-node random topologies / 15 traffic requests (rate: 0 ~ max. load)



# Flow Contention

- Link contention
- QoS consideration
- Resource coordination

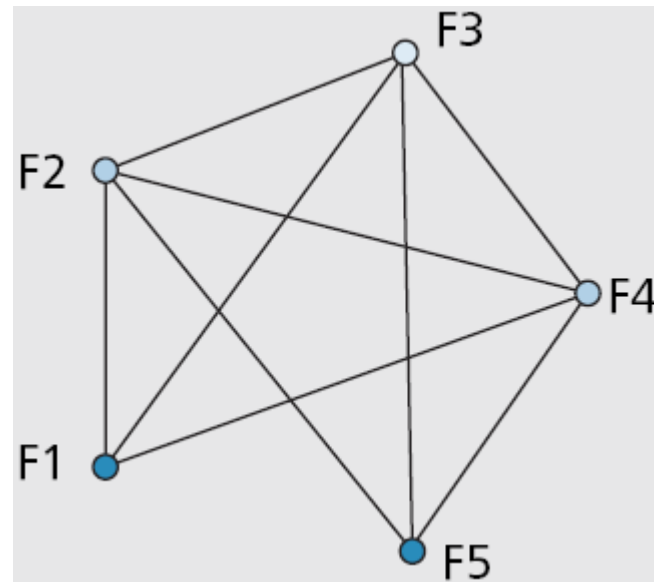


*Spec. of ORiNOCO 802.11b PC Card*

Rates (Mbps)	11.0	5.5	2.0	1.0	CS
Range (m)	399	531	669	796	1783

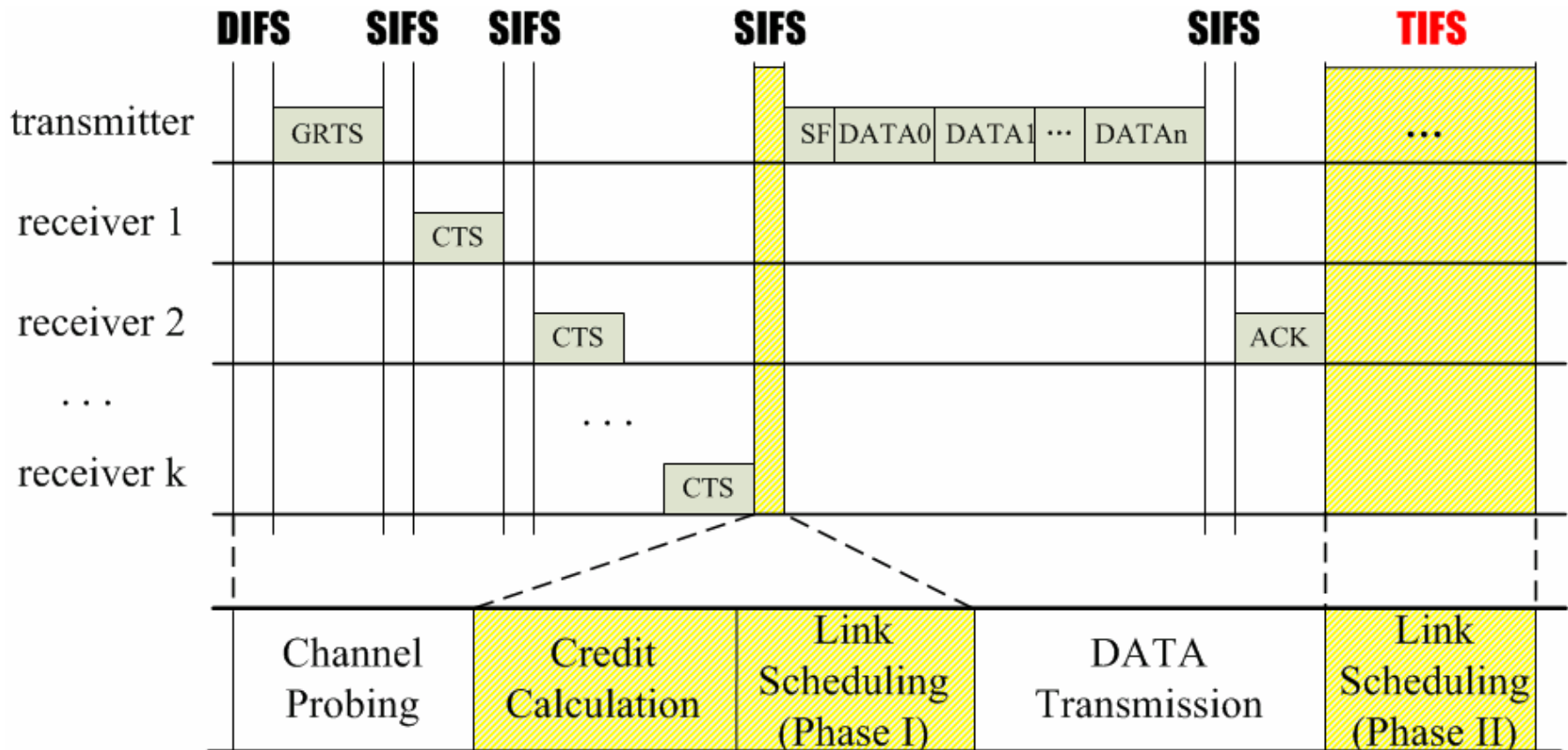
# COS

- Find the globally best set of flows that can transmit simultaneously and maximize overall system performance.
- Channel probing
  - LCG
- Credit calculation
- Data transmission
- Flow scheduling



Local Contention Graph (LCG)

# The 4 Procedures of COS

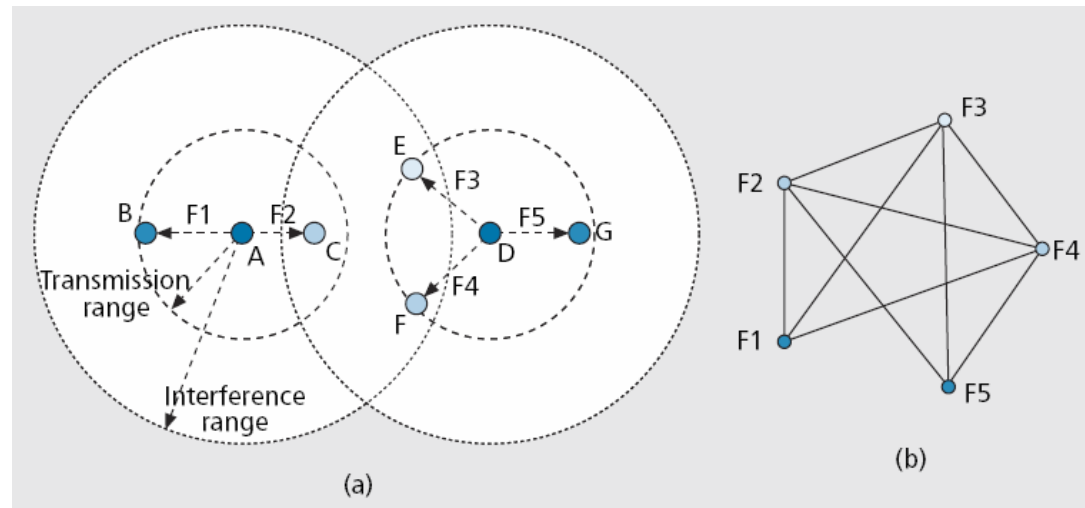


# Credit Calculation

- *The CR* of the MIS (Maximal Independent Subset)
  - $\{S_m\} = \{\{F1, F5\}, \{F2\}, \{F3\}, \{F4\}\}$
  - $CR(S_m) = \sum_{i \in S_m} \mu_i (1 + \lambda_i)$ 
    - $\mu_i$ : the highest rate the  $i$ -th link supports
    - $\lambda_i$ : the QoS factor of the  $i$ -th flow
- *The CR* of the flow: for flow selection
- *The CR* of the transmitter: for TIFS caculation

# An Example

- CR's of MIS's
  - $\{F1, F5\}=7, \{F2\}=6, \{F3\}=5, \{F4\}=4.$
- CR of flows
  - $\{F1, \dots, F5\}=\{7, 6, 5, 4, 7\}.$
- CR of Transmitters
  - $A=7, D=7.$



# Credit Update

- After each transmission, the QoS factor is updated according to:

$$\lambda_i^{k+1} = \begin{cases} \lambda_i^k + a^k (G_i - C_i^k), & \text{if } G_i > C_i^k \\ 0 & \text{otherwise} \end{cases}$$

- $G_i$ : the long-term QoS requirement of the  $i$ -th flow
- $C_i^k$ : the throughput achieved until time slot  $k$
- $a^k = 1/k$



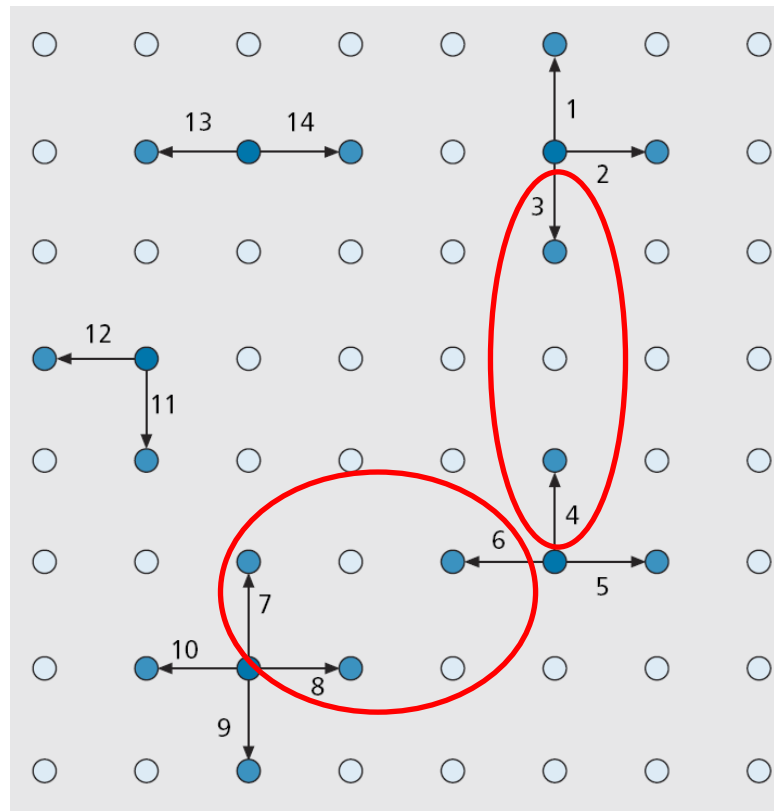
# Flow Scheduling

- Traffic control InterFrame Space
  - Defer non-transmitter's attempt to access for approximate the optimal time scheduling

$$TIFS = \begin{cases} 0, & \text{if } seq = 1 \\ TIFS_{min}, & \text{if } TIFS = 0 \text{ and } seq > 1 \\ \min(TIFS \cdot seq, TIFS_{max}), & \text{otherwise} \end{cases}$$

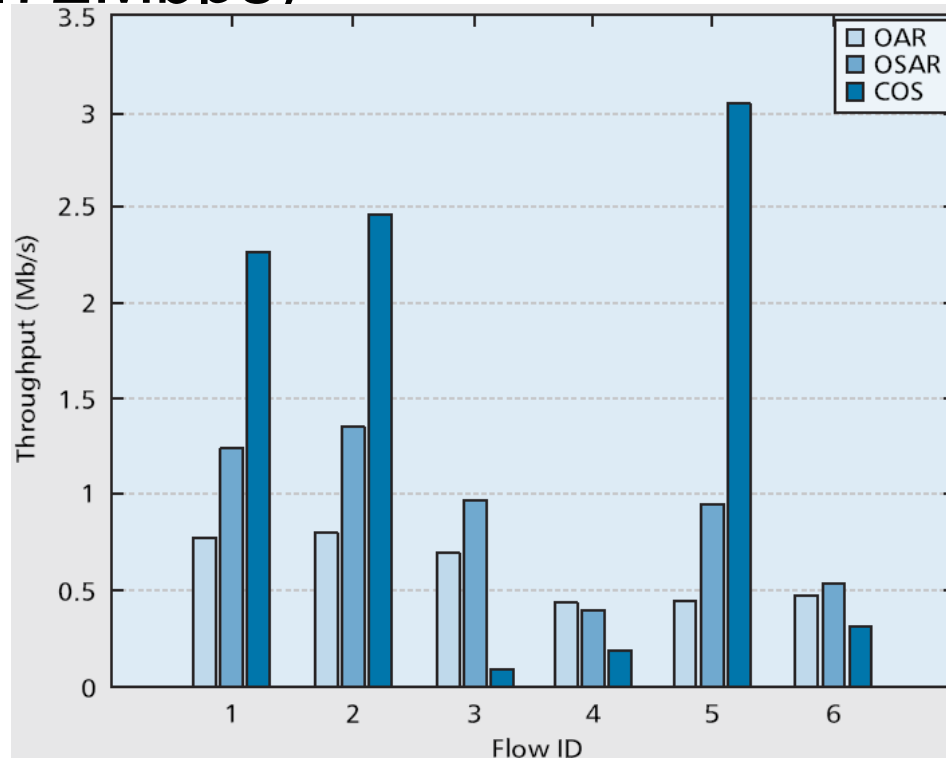
# COS Performance

- 1, 2, 5.5, 11 Mbps / QSAR with flow credit / 8x8-grid=64 nodes / 14 flows



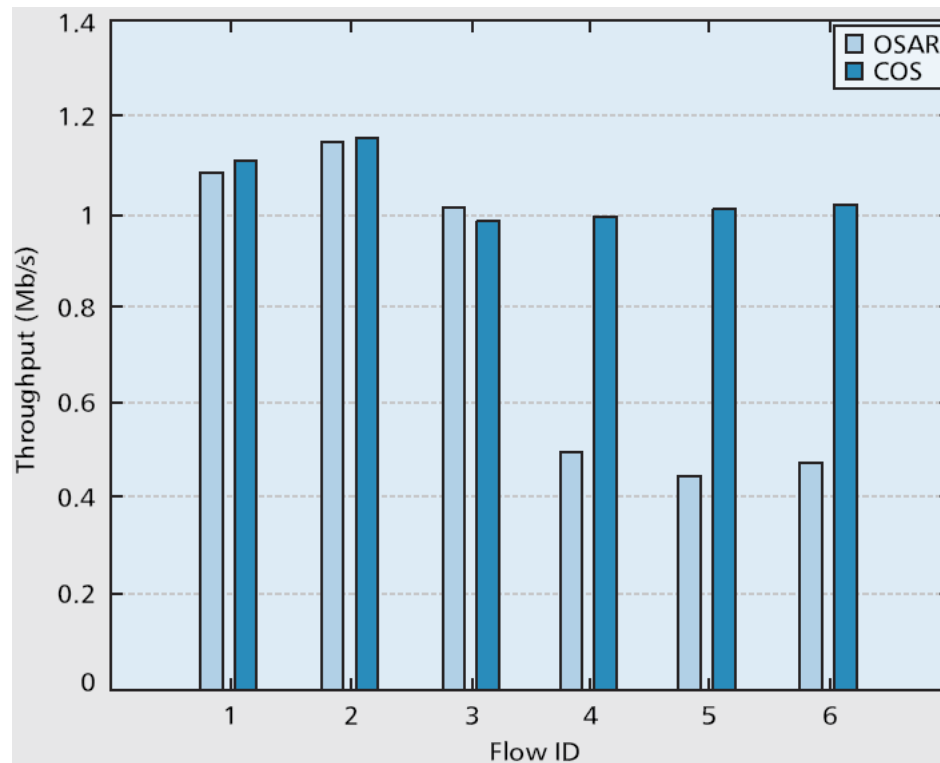
# Throughput w/o QoS

- OAR(8.7Mbps) / OSAR(10.16Mbps) / COS(13.72Mbps)



# Throughput w/ QoS

- $G1=G2=\dots G6=1.0\text{Mbps}$



# Conclusion

- Opportunistically accessing the varying wireless channel opens a new direction for wireless networking related researches.
- Types of diversity that can be considered:
  - Time diversity
  - Multiuser diversity
  - Channel diversity
  - Path diversity
  - Space diversity

# References

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