Combs, Needles, Haystacks: Balancing Push and Pull for Discovery in Large-Scale Sensor Networks

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#### Outline

- □ Introduction
- Comb-Needle Discovery support model
- □ Simulation
- □ Reliability
- □ Adaptive comb-needle Strategy
- Conclusion and Discussion

- Many emerging sensor network applications involve dissemination of observed information to interested clients.
  - Fire emergency system
  - Ecological environment observation
  - Battlefield monitoring

- □ Flooding is a simplest method for obtaining data from a large scale sensor network.
  - Large cost
- Many strategies have been proposed to reduce the cost of discovery and query in large scale sensor network.
  - Improve flooding efficiency
  - Reduce discovery/query cost by taking into account application semantics
  - Reduce search cost by using distributed indexing scheme

(Push-based) Sensor
 nodes push the sensed
 information throughout
 the network.



(Pull-based) The soldier
 broadcast a query for
 the information when it
 is needed.



#### □ Goal

- To investigate efficient and reliable routing mechanisms for supporting queries in large scale ad hoc wireless sensor networks.
- This paper intend to combine the advantages of both push and pull strategies and build an efficient query-support mechanism that adapts to the frequencies of query and events

- Each sensor node pushes its data to a certain neighborhood and the query is disseminated only to a subset of the network.
  - The query process builds a routing structure dynamically that resembles a comb.
  - The sensor nodes push the data duplication structure like a needle.



- □ The principle of the comb-needle structure:
  - Adjust the communication strategy based on the frequency of the query and events.
- Assuming that events and discovery queries occur uniformly in space and time across the sensor network.
  - $f_q$ : the arrival frequency of discovery queries.
  - $f_e$ : the arrival frequency of relevant events.

- □ Compute the minimum push/pull cost
  - Each node sends its data push to 2*l* of its vertical neighbors.
    - $\Box \quad \text{Data push cost: } C_l \sim 2l$
  - The query process first sends the query vertically then fans out horizontally.
    - Query dissemination cost:

$$C_{qd} = n - 1 + (n - 1)(\lfloor \frac{n - 1}{s} \rfloor + 1)$$

- □ Compute the minimum push/pull cost
  - The total cost per query:

$$C = C_{qd} + C_{qr} + f_e * C_l / f_q$$
  

$$\simeq C_{qr} + 2(n-1) + \frac{(n-1)^2}{s} + s * \frac{f_e}{f_q}$$
  

$$\geq C_{qr} + 2(n-1) + 2(n-1) \sqrt{\frac{f_e}{f_q}}$$

This minimum of the cost is around

$$s_{optimal} \sim (n-1) \sqrt{\frac{f_q}{f_e}} \sim \sqrt{N} \sqrt{\frac{f_q}{f_e}}$$

• The query cost is reduced to  $O(\sqrt{N})$ 

- □ The proposed model can cover the whole spectrum of the push and pull strategies.
  - $f_q < f_e$ : the global-pull-local push model
  - $f_q = f_e$  : special push-pull strategy
  - $f_q > f_e$ : reverse combing





Figure 5: Combing

Figure 4: A Special Push-Pull Strategy

Figure 6: Reverse comb

# Simulation

#### □ The simulation environment:

Radio propagation Model

$$P_{rec,ideal}(d) \leftarrow P_{transmit} \frac{1}{1+d^{\gamma}}, \text{ where } 2 \leq \gamma \leq 4$$
$$P_{rec}(i,j) \leftarrow P_{rec,ideal}(d_{i,j})(1+\alpha(i,j))(1+\beta(t))$$

- Topology Model
  - □ Random grid topology
- Routing Protocol
  - □ Constrained Geographical Flooding (CGF)
  - Define CFG width (w), inter-spike distance (s), and the needle push length(l)

## Simulation

 Given query rate and event rate, what is the best spacing for the comb.

• 
$$f_q = 0.1 \text{ p/s} \ f_e = 1 \text{ p/s}$$

• 
$$f_q = 0.1 \text{ p/s}$$
  $f_e = 0.1 \text{ p/s}$ 



Figure 10: Energy consumption:  $f_q = 0.1$ ,  $f_e = 1$ 



 $\label{eq:figure 11: Energy consumption: } f_q = 0.1, \ f_e = 0.1 \\ \mbox{I-Tsung Shen, MNet Lab}$ 

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### Simulation

- Showing the robustness of the protocol with a varying query (CGF) width
  - $f_q = 0.1 \text{ p/s}, f_e = 1 \text{ p/s}$
  - *l*=1, *s* =3



Figure 12: Success rate for different query widths



Figure 13: Energy consumption for different query widths

# Reliability

- Wireless links are unreliable due to the nature of wireless communications.
- This paper study some reliable query and report strategies.
  - Local re-enhancement
    - □ Use interleaving paths
    - Redundancy scheme
  - Spacial redundancy

# Reliability

- Spacial redundancy
  - Constructing an additional vertical query propagation path (level-2 redundancy)
  - Extending the needle length to 2s
    - □ Called 2-level redundancy scheme
    - each report can reach two nodes that are on the query propagation path.

# Reliability



Figure 16: Compare the difference between the calculation and approximation.

#### Figure 15: level-2 redundancy

# Adaptive comb-needle Strategy

- In practice, the query and event frequencies may be time-varying, and thus a good query strategy should adapt to such changes.
- □ Based on the estimate of  $f_q$  and  $f_d$ , the query node calculates *s*.
  - $f_d$ : the probability that a query in generated in a time slot

$$f_d = rac{f_e}{n^2} \qquad \qquad s_{optimal} \sim \sqrt{rac{f}{f}}$$

# Adaptive comb-needle Strategy

Data node estimates the total needle length L by using the query success probability.

$$P_s = P(L \ge s) + E\left(\frac{L}{s}|L < s\right)P(L < s)$$
$$= E\left(\min\left(\frac{L}{s}, 1\right)\right).$$

- □ Since the value of  $f_d$  is small, the value of s obtain form previous queries are important for estimating L.
  - Obtaining s when its data is successfully queried
  - Rotation of the query horizontal duplications

## Adaptive comb-needle Strategy



Figure 17: The query frequency and the event frequency.



Figure 18: The ideal comb width and the estimated comb width.

# Conclusion and Discussion

- This paper proposed the comb-needle model for supporting queries in large scale sensor networks.
- Some enhancement of the reliable schemes have proposed.
- An adaptive comb-needle Strategy have proposed for the query and event frequencies time-varying environment.
- □ The total cost is reduced to  $O(\sqrt{N})$