Error Control Strategies for WiMAX Multi-Hop Relay Networks

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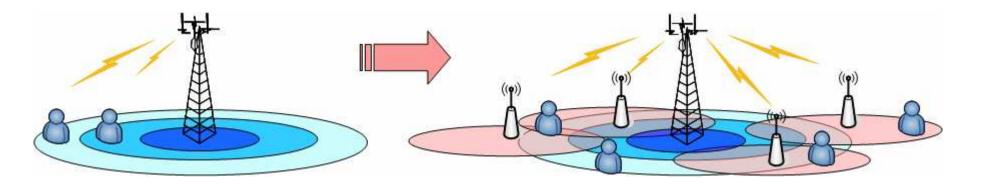
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Outlines

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- Error Control Analysis
- Performance Evaluations
- Conclusions

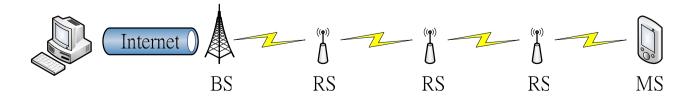
Introduction

- Both advanced WiMAX (802.16m) and LTE (LTE-A) have adapted multi-hop relay architecture for
 - Throughput enhancement
 - Coverage extension
- However, throughput degrades as the number of hops increases
 - Even in an error-free multi-hop environment
- For meet the performance and reliability requirements of 4G, an efficient error control mechanism is desired
 - Relay ARQ and HARQ

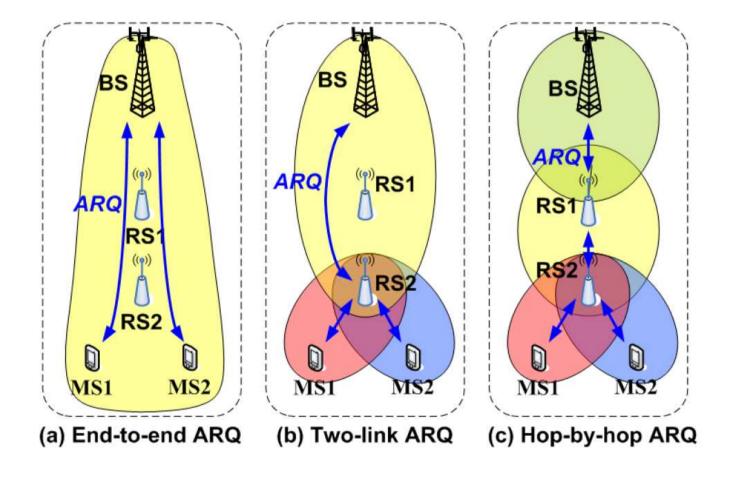


Introduction

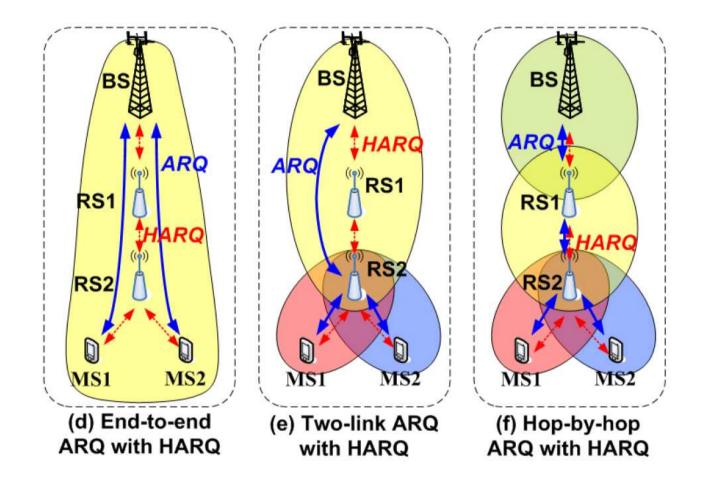
- The factors dominating the performance
 - Buffer size
 - Propagation delay
 - Transmission error probability
 - Acknowledgement feedback
- Different values of the factors lead to completely different conclusions
 - Link-by-link approach is better since it can retransmit packet as soon as possible
 - End-to-end approach is better if feedback delay is considered
- This paper provides
 - The investigation of error control protocols and the combinations
 - A low complexity error control mechanism for multi-hop relay transmission



- Basically, there are three relay ARQs
 - End-to-End ARQ
 - ARQ is performed between BS and MS
 - 2-Link ARQ
 - In first link, ARQ is performed between BS and the last hop RS
 - In second link, ARQ is performed between the last hop RS and MS
 - Hop-by-Hop ARQ
 - ARQ is initiated for each hop



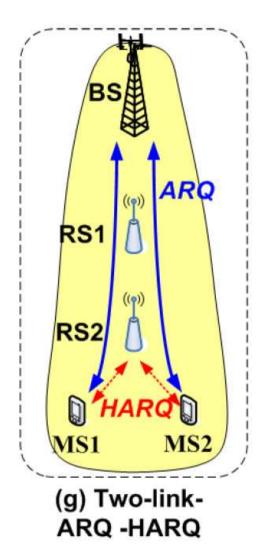
- The combination of ARQ and HARQ
 - End-to-End ARQ with HARQ
 - Per-hop HARQ under End-to-End ARQ
 - 2-Link ARQ with HARQ
 - Per-hop HARQ under 2-Link ARQ
 - Hop-by-Hop ARQ with HARQ
 - ARQ and HARQ are used for every hop
 - Most complicated
 - Every RS has to build state machine and allocate resource for HARQ
 - Large amount of HARQ state machine has to be collected from all RSs during handover



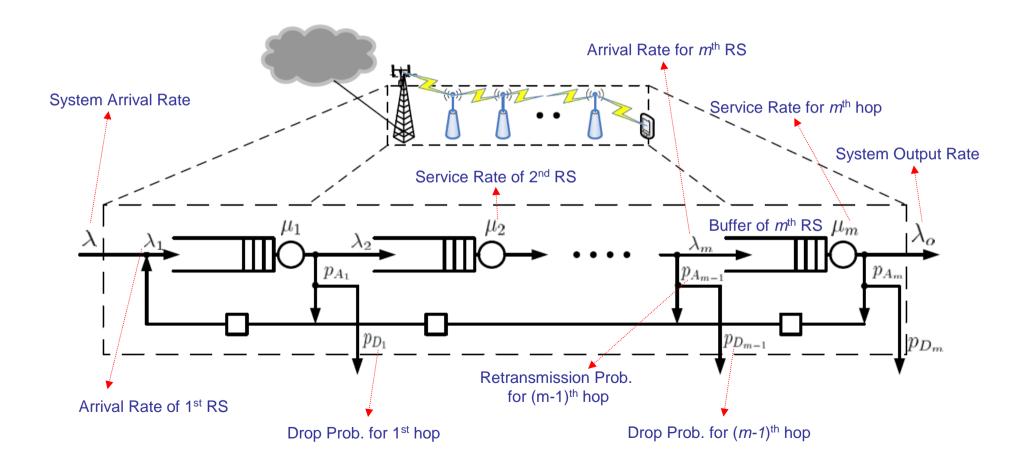
- Observations
 - Relay link (BS⇔RS) has much lower block error rate (BLER) as compared to that of access link (RS⇔MS)
 - RSs are usually stationary and the location is predetermined
 - Per-Hop HARQ is not necessary for relay links
 - MS may have various mobility
 - Higher BLER in access link
 - The usage of HARQ helps in access link

Proposal

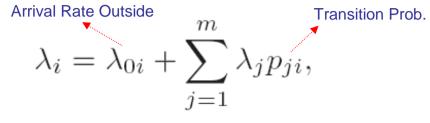
- HARQ only applies in access link
 - Less state machine used and exchanged
- ARQ is between BS and MS for End-to-End error control
 - Relay link can be optimized in implementation
 - BLER is usually low



• Multi-Hop Error Control



• Arrival rate at each queue



- Arrival rate at the first node
 - System arrival + retransmissions

$$\lambda_1 = \lambda + \sum_{j=1}^m \lambda_j p_{A_j} (1 - p_{D_j}).$$

• Arrival rate at *i*th hop node

- Visiting rate at *i*th hop node
 - ARQ burst visiting the *i*th hop node

$$\lambda_{i} = \frac{\lambda P_{l}^{(i-1)}}{1 - \sum_{j=1}^{m} p_{A_{j}} (1 - p_{D_{j}}) P_{l}^{(j-1)}}$$
$$e_{i} = \frac{P_{l}^{(i-1)}}{1 - \sum_{j=1}^{m} p_{A_{j}} (1 - p_{D_{j}}) P_{l}^{(j-1)}}$$

• Mean response time

Service Time

$$\overline{T}_Q(i) = \frac{E[n_i]}{\lambda_i} + T_c(i)$$

• Acknowledge Time (From destination to source)

$$T_{ACK} = \sum_{i=1}^m rac{1}{\mu_i} \stackrel{ ext{Propagation Delay}}{+ m T_p + T_A}_{ ext{Receiving burst and sending the ACK}}$$

• Average delay for ARQ transmission

$$\overline{T}(m, \mathbf{p}_A) = \sum_{i=1}^m e_i \overline{T}_Q(i) + (e_1 - 1) T_{ACK}$$

• Given *h* hops links from BS to MS, the average delay for ARQ transmission is:

$$\overline{D}(h) = \begin{cases} \sum_{i=1}^{h} \overline{T}(1, p_{a_i}), & \text{hop-by-hop} \\ \overline{T}(h - 1, \mathbf{p}_a^{h-1}) + \overline{T}(1, p_{a_h}), & 2\text{-link} \\ \overline{T}(h, \mathbf{p}_a), & \text{end-to-end}, \end{cases}$$

 $P_a = \{P_{a1}, \dots, P_{ah}\}$ as the prob. vector of burst error on total *h* hop links

 The connection delay for various ARQ with HARQ is

$$\overline{D}_{H}(h) = \begin{cases} \sum_{i=1}^{h} \overline{T}_{H}(1, p_{r_{i}}), & \text{hop-by-hop} \\ \frac{\overline{T}_{H}(h-1, \mathbf{p}_{r}^{h-1}) + \overline{T}_{H}(1, p_{r_{h}}), & 2\text{-link} \\ \overline{T}_{H}(h, \mathbf{p}_{r}), & \text{end-to-end.} \end{cases}$$

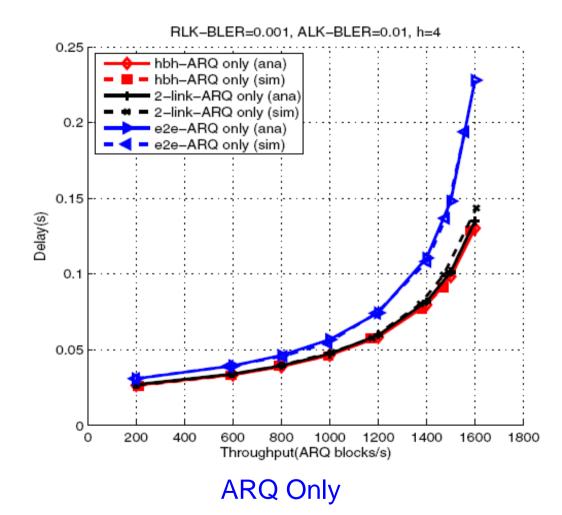
 $\begin{array}{ll} \mbox{Response Time} & \mbox{Acknowledge Time} \\ \overline{T}_Q^H(i) = \frac{E[n_i]}{\lambda_i} + T_c^H(i) & \mbox{} T_{ACK}^H = \sum_{i=1}^m \frac{1}{\mu_i^H} + mT_p + T_H \end{array}$

• The delay for the proposed method is:

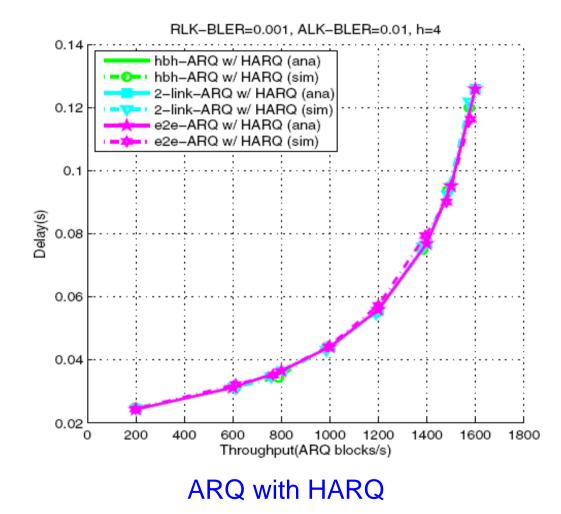
$$\overline{D}_{AH}(h) = \overline{T}_{AH}(h, \{\mathbf{p}_a^{h-1}, p_{r_h}\})$$

$$T_{ACK}^{AH} = \sum_{i=1}^{m-1} \frac{1}{\mu_i} + \frac{1}{\mu_h^H} + mT_p + T_H < T_{ACK}^H$$
Only access link applies HARQ

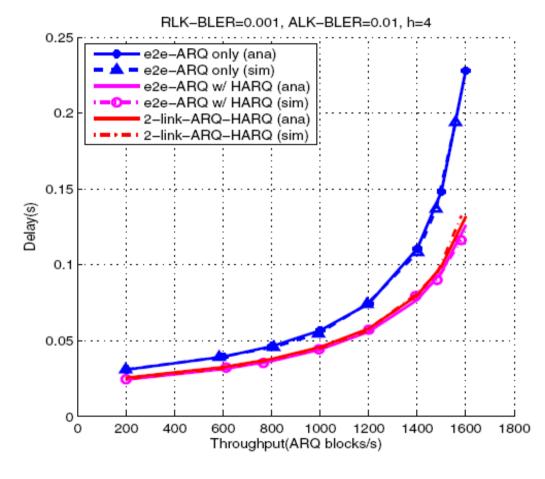
Performance Evaluations



Performance Evaluations



Performance Evaluations



Proposed ARQ vs End-to-End ARQ

Conclusions

- Error controls for multi-hop relay environment are analyzed
 - The analytical model tells that pure ARQ do not perform well in the multi-hop network
 - Retransmission timer has to increase as the number of hop increases
 - Due to longer delay for retransmission
- HARQ can be used in single-hop, but not in multi-hop
 - Error introduced at intermediate RSs will increase since error probability increases with hop-count also
- Compared with End-to-End ARQ with per-hop HARQ, The proposed combination of ARQ and HARQ provide
 - Lower complexity
 - Less information exchange during handover process